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MARKOV EQUIVALENCE OF MARGINALIZED LOCAL INDEPENDENCE GRAPHS

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Symmetric independence relations are often studied using graphical representations. Ancestral graphs or acyclic directed mixed graphs with $m$-separation provide classes of symmetric graphical independence models that are closed under marginalization. Asymmetric independence relations appear naturally for multivariate stochastic processes, for instance, in terms of local independence. However, no class of graphs representing such asymmetric independence relations, which is also closed under marginalization, has been developed. We develop the theory of directed mixed graphs with $\mu$-separation and show that this provides a graphical independence model class which is closed under marginalization and which generalizes previously considered graphical representations of local independence.

Several graphs may encode the same set of independence relations and this means that in many cases only an equivalence class of graphs can be identified from observational data. For statistical applications, it is therefore pivotal to characterize graphs that induce the same independence relations. Our main result is that for directed mixed graphs with $\mu$-separation each equivalence class contains a maximal element which can be constructed from the independence relations alone. Moreover, we introduce the directed mixed equivalence graph as the maximal graph with dashed and solid edges. This graph encodes all information about the edges that is identifiable from the independence relations, and furthermore it can be computed efficiently from the maximal graph.

1. Introduction. Graphs have long been used as a formal tool for reasoning with independence models. Most work has been concerned with symmetric independence models arising from standard probabilistic independence for discrete or real-valued random variables. However, when working with dynamical processes it is useful to have a notion of independence that can distinguish explicitly between the present and the past, and this is a key motivation for considering local independence.

The notion of local independence was introduced for composable Markov processes by Schweder [37] who also gave examples of graphs describing local independence structures. Aalen [1] discussed how one could extend the definition of local independence in the broad class of semimartingales using the Doob–Meyer decomposition. Several authors have since then used graphs to represent local independence structures of multivariate stochastic process models, in particular for point process models; see, for example, [4, 11–13, 35]. Local independence takes a dynamical point of view in the sense that it evaluates the dependence of the present on the past. This provides a natural link to statistical causality as cause must necessarily precede effect [1, 2, 28, 37]. Furthermore, recent work argues that for some applications it can be important to consider continuous-time models, rather than only cross-sectional models, when trying to infer causal effects [3].

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Local independence for point processes has been applied for data analysis (see, e.g., [2, 23, 44]), but in applications a direct causal interpretation may be invalid if only certain dynamical processes are observed while other processes of the system under study are unobserved. Allowing for such latent processes is important for valid causal inference, and this motivates our study of representations of marginalized local independence graphs.

Graphical representations of independence models have also been studied for time series [14–17]. In the time series context—using the notion of Granger causality—Eichler [15] gave an algorithm for learning a graphical representation of local independence. However, the equivalence class of graphs that yield the same local independences was not identified, and thus the learned graph does not have any clear causal interpretation. Related research has been concerned with inferring the graph structure from subsampled time series, but under the assumption of no latent processes; see, for example, [9, 22].

In this paper, we give a formal, graphical framework for handling the presence of unobserved processes and extend the work on graphical representations of local independence models by formalizing marginalization and giving results on the equivalence classes of such graphical representations. The graphical framework that we propose is a generalization of that of Didelez [11–13]. This development is analogous to work on marginalizations of graphical models using directed acyclic graphs, DAGs. Starting from a DAG, one can find graphs (e.g., maximal ancestral graphs or acyclic directed mixed graphs) that encode marginal independence models [8, 18, 19, 25, 33, 34, 36, 39]. One can then characterize the equivalence class of graphs that yield the same independence model [5, 45]—the so-called Markov equivalent graphs—and construct learning algorithms to find such an equivalence class from data. The purpose of this paper is to develop the necessary theoretical foundation for learning local independence graphs by developing a precise characterization of the learnable object: the class of Markov equivalent graphs.

The paper is structured as follows: in Section 2, we discuss abstract independence models, relevant graph-theoretical concepts and the notion of local independence and local independence graphs. In Section 3, we introduce $\mu$-separation for directed mixed graphs, which will be used to represent marginalized local independence graphs, and we describe an algorithm to marginalize a given local independence graph. In Sections 4 and 5, we develop the theory of $\mu$-separation for directed mixed graphs further, and we discuss, in particular, Markov equivalence of such graphs. All proofs of the main paper are given in the Supplementary Material [29]. Sections A to F are in the Supplementary Material.

### 2. Independence models and graph theory

Graphical separation criteria as well as probabilistic models give rise to abstract conditional independence statements. Graphical modeling is essentially about relating graphical separation to probabilistic independence. We will consider both as instances of abstract independence models.

Consider some set $\mathcal{S}$. An independence model, $\mathcal{I}$, on $\mathcal{S}$ is a set of triples $(A, B, C)$ where $A, B, C \in \mathcal{S}$, that is, $\mathcal{I} \subseteq \mathcal{S} \times \mathcal{S} \times \mathcal{S}$. Mathematically, an independence model is a ternary relation. In this paper, we will consider independence models over a finite set $V$ which means that $\mathcal{S} = \mathcal{P}(V)$, the power set of $V$. In this case, an independence model $\mathcal{I}$ is a subset of $\mathcal{P}(V) \times \mathcal{P}(V) \times \mathcal{P}(V)$. We will call an element $s \in \mathcal{P}(V) \times \mathcal{P}(V) \times \mathcal{P}(V)$ an independence statement and write $s$ as $(A, B \mid C)$ for $A, B, C \subseteq V$. This notation emphasizes that $s$ is thought of as a statement about $A$ and $B$ conditionally on $C$.

Graphical and probabilistic independence models have been studied in very general settings, though mostly under the assumption of symmetry of the independence model, that is,

$$(A, B \mid C) \in \mathcal{I} \Rightarrow (B, A \mid C) \in \mathcal{I};$$

see, for example, [7, 10, 26] and references therein. These works take an abstract axiomatic approach by describing and working with a number of properties that hold in, for example,
models of conditional independence. In this paper, we consider independence models that do not satisfy the symmetry property as will become evident when we introduce the notion of local independence.

2.1. Local independence. We consider a real-valued, multivariate stochastic process

\[ X_t = (X^1_t, X^2_t, \ldots, X^n_t), \quad t \in [0, T] \]

defined on a probability space \((\Omega, \mathcal{F}, P)\). In this section, the process is a continuous-time process indexed by a compact time interval. The case of a discrete time index, corresponding to \( X = (X_t) \) being a time series, is treated in Section C in the Supplementary Material. We will later identify the coordinate processes of \( X \) with the nodes of a graph; hence, both are indexed by \( V = \{1, 2, \ldots, n\} \). As illustrated in Example 2.3 below, the index set may be chosen in a more meaningful way for a specific application. In that example, \( X^i_t \geq 0 \) is a price process, \( X^L_t \in \mathbb{N}_0 \) is a counting process of events, and the remaining four processes take values in \( \{0, 1\} \) indicating if an individual at a given time is a regular user of a given substance. Figure 1 shows examples of sample paths for three individuals.

To avoid technical difficulties, irrelevant for the present paper, we restrict attention to right-continuous processes with coordinates of finite and integrable variation on the interval \([0, T]\). This includes most nonexplosive multivariate counting processes as an important special case, but also other interesting processes such as piecewise-deterministic Markov processes.

To define local independence below, we need a mathematical description of how the stochastic evolution of one coordinate process depends infinitesimally on its own past and the past of the other processes. To this end, let \( \mathcal{F}^{C,0}_t \) denote the \( \sigma \)-algebra generated by \( \{X^\alpha_s : s \leq t, \alpha \in C\} \) for \( C \subseteq V \). For technical reasons, we need to enlarge this \( \sigma \)-algebra, and we define \( \mathcal{F}^C_t \) to be the completion of \( \bigcap_{s > t} \mathcal{F}^{C,0}_s \) w.r.t. \( P \). Thus \( (\mathcal{F}^C_t) \) is a right-continuous and complete filtration which represents the history of the processes indexed by \( C \subseteq V \) until time \( t \). Figure 2 illustrates, in the context of Example 2.3, the filtrations \( \mathcal{F}^V_t \), \( \mathcal{F}^{\{L,M,H\}}_t \) and \( \mathcal{F}^{\{T,A,M,H\}}_t \).

For \( \beta \in V \) and \( C \subseteq V \), let \( \Lambda^{C,\beta}_t \) denote an \( \mathcal{F}^C_t \)-predictable process of finite and integrable variation such that

\[ E(X^\beta_t | \mathcal{F}^C_t) = \Lambda^{C,\beta}_t \]
is an $\mathcal{F}_t^C$ martingale. Such a process exists (see Section E for the technical details), and is usually called the compensator or the dual predictable projection of $E(X^{\beta} | \mathcal{F}_t^C)$. It is in general unique up to evanescence.

**Definition 2.1 (Local independence).** Let $A, B, C \subseteq V$. We say that $X^B$ is locally independent of $X^A$ given $X^C$ if there exists an $\mathcal{F}_t^C$-predictable version of $\Lambda_{A \cup C, \beta}$ for all $\beta \in B$. We use $A \not\rightarrow B \mid C$ to denote that $X^B$ is locally independent of $X^A$ given $X^C$.

In words, the process $X^B$ is locally independent of $X^A$ given $X^C$ if, for each time point, the past up until time $t$ of $X^C$ gives us the same predictable information about $E(X^{\beta}_t | \mathcal{F}_t^{A \cup C})$ as the past of $X^{A \cup C}$ until time $t$. Note that when $\beta \in C$, $E(X^{\beta}_t | \mathcal{F}_t^C) = X^\beta_t$.

Local independence was introduced by Schweder [37] for composable Markov processes and extended by Aalen [1]. Local independence and graphical representations thereof were later considered by Didelez [11–13] and by Aalen et al. [4]. Didelez [12] also discussed local independence models of composable finite Markov processes under some specific types of marginalization. Commenges and Gégout-Petit [6, 21] discussed definitions of local independence in classes of semimartingales. Note that Definition 2.1 allows a process to be separated from itself by some conditioning set $C$, generalizing the definition used, for example, by Didelez [13].

Local independence defines the independence model
$$\mathcal{I} = \{ (A, B \mid C) \mid X^B \text{ is locally independent of } X^A \text{ given } X^C \}$$
such that the local independence statement $A \not\rightarrow B \mid C$ is equivalent to $(A, B \mid C) \in \mathcal{I}$ in the abstract notation. We note that the local independence model is generally not symmetric. Using Definition 2.1, we introduce below an associated directed graph in which there is no directed edge from a node $\alpha$ to a node $\beta$ if and only if $\beta$ is locally independent of $\alpha$ given $V \setminus \{\alpha\}$.

**Definition 2.2 (Local independence graph).** For the local independence model determined by $X$, we define the local independence graph to be the directed graph, $\mathcal{D}$, with nodes $V$ such that for $\alpha, \beta \in V$,
$$\alpha \not\rightarrow_{\mathcal{D}} \beta \iff \alpha \not\rightarrow \beta \mid V \setminus \{\alpha\}$$
where $\alpha \not\rightarrow_{\mathcal{D}} \beta$ denotes that there is no directed edge from $\alpha$ to $\beta$ in the graph $\mathcal{D}$.
Didelez [11] gives almost the same definition of a local independence graph, however, in essence always assumes that there is a dependence of each process on its own past. See also Sections A and B.

The local independence graph induces an independence model by \( \mu \)-separation as defined below. The main goal of the present paper is to provide a graphical representation of the induced independence model for a subset of coordinate processes corresponding to the case where some processes are unobserved. This is achieved by establishing a correspondence, which is preserved under marginalization, between directed mixed graphs and independence models induced via \( \mu \)-separation. We emphasize that the correspondence only relates local independence to graphs when the local independence model satisfies the global Markov property with respect to a graph.

The local independence model satisfies the global Markov property with respect to the local independence graph if every \( \mu \)-separation in the graph implies a local independence. This has been shown for point processes under some mild regularity conditions [13] using the slightly different notion of \( \delta \)-separation. Section A discusses how \( \delta \)-separation is related to \( \mu \)-separation, and Section B shows how to translate the global Markov property of [13] into our framework. Moreover, general sufficient conditions for the global Markov property were given in [30] covering point processes as well as certain diffusion processes. Section C provides, in addition, a discussion of Markov properties in the context of time series.

To help develop a better understanding of local independence and its relevance for applications, we discuss an example of drug abuse progression.

**Example 2.3 (Gateway drugs).** The theory of gateway drugs has been discussed for many years in the literature on substance abuse [24, 40]. In short, the theory posits that the use of “soft” and often licit drugs precedes (and possibly leads to) later use of “hard” drugs. Alcohol, tobacco and marijuana have all been discussed as candidate gateway drugs to “harder” drugs such as heroin.

We propose a hypothetical, dynamical model of transitions into abuse via a gateway drug, and more generally, a model of substance abuse progression. Substance abuse is known to be associated with social factors, genetics and other individual and environmental factors [43]. Substance abuse can evolve over time when an individual starts or stops using some drug. In this example, we consider substance processes Alcohol (A), Tobacco (T), Marijuana (M) and Hard drugs (H) modeled as zero-one processes, that is, stochastic processes that are piecewise constantly equal to zero (no substance use) or one (substance use). We also include \( L \), a process describing life events, and a process \( I \), which can be thought of as an exogenous process that influences the tobacco consumption of the individual, for example, the price of tobacco which may change due to changes in tobacco taxation. Let \( V = \{A, T, M, H, L, I\} \).

We will visualize each process as a node in a graph and draw an arrow from one process to another if the first has a direct influence on the second. We will not go into a full discussion of how to formalize “influence” in terms of a continuous-time causal dynamical model as this would lead us astray; see instead [13, 27, 38]. The upshot is that for a (faithful) causal model, there is no direct influence if and only if \( \alpha \not \rightarrow \beta \mid V \setminus \{\alpha\} \), which identifies the “influence” graph with the local independence graph.

Several formalizations of the gateway drug question are possible. We will focus on the questions “is the use of hard drugs locally independent of use of alcohol for some conditioning set?” and “is the use of hard drugs locally independent of the use of tobacco for some conditioning set?” Using the dynamical nature of local independence, we are asking if, for example, the past alcohol usage changes the hard drug usage propensity when accounting for the past of all other processes in the model. This is one possible formalization of the gateway drug question as a negative answer would mean that there exist some gateway processes through which any influence of alcohol usage on hard drug usage is mediated. If the
visualizaion in Figure 3 is indeed a local independence graph in the above sense we see that conditioning on all other processes, H is indeed locally independent of A and locally independent of T. In this hypothetical scenario, we could interpret this as marijuana in fact acting as a gateway drug to hard drugs. If the global Markov property holds, we can furthermore use $\mu$-separation to obtain further local independences from the graph. We return to this example in Section 5.5 to illustrate how the main results of the paper can be applied. In particular, we are interested in what conclusions we can make when we do not observe all the processes but only a subset.

2.2. Marginalization and separability.

**Definition 2.4 (Marginalization).** Given an independence model $\mathcal{I}$ over $V$, the marginal independence model over $O \subseteq V$ is defined as

$$\mathcal{I}^O = \{ \langle A, B \mid C \rangle \mid \langle A, B \mid C \rangle \in \mathcal{I}; A, B, C \subseteq O \}.$$

Marginalization is defined abstractly above, though we are primarily interested in the marginalization of the independence model encoded by a local independence graph via $\mu$-separation. The main objective is to obtain a graphical representation of such a marginalized independence model involving only the nodes $O$. To this end, we consider the notion of separability in an independence model.

**Definition 2.5 (Separability).** Let $\mathcal{I}$ be an independence model over $V$. Let $\alpha, \beta \in V$. We say that $\beta$ is separable from $\alpha$ if there exists $C \subseteq V \setminus \{\alpha\}$ such that $\langle \alpha, \beta \mid C \rangle \in \mathcal{I}$, and otherwise we say that $\beta$ is inseparable from $\alpha$. We define

$$s(\beta, \mathcal{I}) = \{ \gamma \in V \mid \beta \text{ is separable from } \gamma \}. $$

We also define $u(\beta, \mathcal{I}) = V \setminus s(\beta, \mathcal{I})$.

We show in Proposition 3.6 that if $\mathcal{I}$ is the independence model induced by a directed graph via $\mu$-separation, then $\alpha \in u(\beta, \mathcal{I})$ if and only if there is a directed edge from $\alpha$ to $\beta$. In this case, the graph is thus directly identifiable from separability properties of $\mathcal{I}$. That is, however, not true in general for a marginalization of $\mathcal{I}$, and this is the motivation for developing a theory of directed mixed graphs with $\mu$-separation.

2.3. Graph theory. A graph, $G = (V, E)$, is an ordered pair where $V$ is a finite set of vertices (also called nodes) and $E$ is a finite set of edges. Furthermore, there is a map that to each edge assigns a pair of nodes (not necessarily distinct). We say that the edge is between these two nodes. We consider graphs with two types of edges: directed ($\rightarrow$) and bidirected.
We can think of the edge set as a disjoint union, $E = E_d \cup E_b$, where $E_d$ is a set of ordered pairs of nodes $(\alpha, \beta)$ corresponding to directed edges, and $E_b$ is a set of unordered pairs of nodes $\{\alpha, \beta\}$ corresponding to bidirected edges. This implies that the edge $\alpha \leftrightarrow \beta$ is identical to the edge $\beta \leftrightarrow \alpha$, but the edge $\alpha \rightarrow \beta$ is different from the edge $\beta \rightarrow \alpha$. It also implies that the graphs we consider can have multiple edges between a pair of nodes $\alpha$ and $\beta$, but they will always be a subset of the edges $\{\alpha \rightarrow \beta, \alpha \leftarrow \beta, \alpha \leftrightarrow \beta\}$.

**Definition 2.6 (DMG).** A directed mixed graph (DMG), $G = (V, E)$, is a graph with node set $V$ and edge set $E$ consisting of directed and bidirected edges as described above.

Throughout the paper, $G$ will denote a DMG with node set $V$ and edge set $E$. Occasionally, we will also use $D$ and $M$ to denote DMGs. We use $D$ only when the DMG is also a directed graph, that is, has no bidirected edges. We use $M$ to stress that some DMG is obtained as a marginalization of a DMG on a larger node set. We will use notation such as $\leftrightarrow_G$ or $\rightarrow_D$ to denote the specific graph that an edge belongs to.

If $\alpha \rightarrow \beta$, we say that the edge has a *tail* at $\alpha$ and a *head* at $\beta$. Jointly tails and heads are called (edge) *marks*. An edge $e \in E$ between nodes $\alpha$ and $\beta$ is a *loop* if $\alpha = \beta$. We also say that the edge is *incident* with the node $\alpha$ and with the node $\beta$ and that $\alpha$ and $\beta$ are *adjacent*.

For $\alpha, \beta \in V$, we use the notation $\alpha \sim \beta$ to denote a generic edge of any type between $\alpha$ and $\beta$. We use the notation $\alpha \leftarrow\rightarrow \beta$ to indicate an edge that has a head at $\beta$ and may or may not have a head at $\alpha$. Note that the presence of one edge, $\alpha \rightarrow \beta$, say, does not in general preclude the presence of other edges between these two nodes. Finally, $\alpha \leftarrow\rightarrow\rightarrow G \beta$ means that there is no edge in $G$ between $\alpha$ and $\beta$ that has a head at $\beta$ and $\alpha \leftarrow\rightarrow G \beta$ means that there is no directed edge from $\alpha$ to $\beta$. Note that $\alpha \leftarrow\rightarrow G \beta$ is a statement about the absence of an edge in the graph $G$ and to avoid confusion with local independence, $\alpha \leftarrow\rightarrow \beta \mid C$, we always include the conditioning set when writing local independence statements, even if $C = \emptyset$ (see also Definition 2.2).

We say that $\alpha$ is a *parent* of $\beta$ in the graph $G$ if $\alpha \rightarrow \beta$ is present in $G$ and that $\beta$ is a *child* of $\alpha$. We say that $\alpha$ is a *sibling* of $\beta$ (and that $\beta$ is a sibling of $\alpha$) if $\alpha \leftrightarrow \beta$ is present in the graph. The motivation of the term sibling will be explained in Section 3. We use $pa(\alpha)$ to denote the set of parents of $\alpha$.

A *walk* is an ordered, alternating sequence of vertices, $\gamma_1$, and edges, $e_j$, denoted $\omega = \langle \gamma_1, e_1, \ldots, e_n, \gamma_n \rangle$, such that each $e_i$ is between $\gamma_i$ and $\gamma_{i+1}$, along with an orientation of each directed loop along the walk (if $e_i$ is a loop then we also know if $e_i$ points in the direction of $\gamma_i$ or in the direction of $\gamma_{i+1}$). Without the orientation, for instance, the walks $\alpha \rightarrow \beta \rightarrow \gamma$ and $\alpha \rightarrow \beta \leftarrow \beta \rightarrow \gamma$ would be indistinguishable. See Figure 4 for examples. We will often present the walk $\omega$ using the notation

$$\gamma_1 \sim e_1 \sim \cdots \sim e_n \gamma_{n+1},$$

where the loop orientation is explicit. We will omit the edge superscripts when they are not needed.

![Fig. 4](image_url)  

A directed mixed graph with node set $\{\alpha, \beta, \gamma, \delta\}$. Consider first the walk $\alpha \rightarrow \beta$. This is different from the walk $\beta \leftarrow \alpha$ as walks are ordered. Consider instead the two walks $\beta \leftrightarrow \gamma \leftarrow \delta$ and $\beta \leftrightarrow \gamma \rightarrow \delta$. These two walks have the same (ordered) sets of nodes and edges but are not equal as the loop at $\gamma$ has different orientations between the two walks. Furthermore, one can note that for the first of the two walks, $\gamma$ is a collider in the first instance, but not in the second. The walks $\alpha \rightarrow \beta \rightarrow \alpha$ and $\alpha \rightarrow \beta \leftarrow \alpha$ are both cycles, and the second is an example of the fact that the same edge can occur twice in a cycle.
We say that the walk \( \omega \) contains nodes \( \gamma_i \) and edges \( e_j \). The length of the walk is \( n \), the number of edges that it contains. We define a trivial walk to be a walk with no edges and, therefore, only a single node. Equivalently, a trivial walk can be defined as a walk of length zero. A subwalk of \( \omega \) is either itself a walk of the form \( \langle \gamma_k, e_k, \ldots, e_m, \gamma_m \rangle \) where \( 1 \leq k < m \leq n + 1 \) or a trivial walk \( \langle \gamma_k \rangle, 1 \leq k \leq n + 1 \). A (nontrivial) walk is uniquely identified by its edges, and the ordering and orientation of these edges, hence the vertices can be omitted when describing the walk. At times, we will omit the edges to simplify notation, however, we will always have a specific, uniquely identified walk in mind even when the edges and/or their orientation is omitted. The first and last nodes of a walk are called endpoint nodes (these could be equal) or just endpoints, and we say that a walk is between its endpoints, or alternatively from its first node to its last node. We call the walk \( \omega^{-1} = \langle \gamma_{n+1}, e_n, \ldots, e_1, \gamma_1 \rangle \) the inverse walk of \( \omega \). Note that the orientation of directed loops is also reversed in the inverse walk such that they point toward \( \gamma_1 \) in the inverse if and only if they point toward \( \gamma_1 \) in the original walk. A path is a walk on which no node is repeated.

Consider a walk \( \omega \) and a subwalk thereof, \( \langle \alpha, e_1, \gamma, e_2, \beta \rangle \), where \( \alpha, \beta, \gamma \in V \) and \( e_1, e_2 \in E \). If \( e_1 \) and \( e_2 \) both have heads at \( \gamma \), then \( \gamma \) is a collider on \( \omega \). If this is not the case, then \( \gamma \) is a noncollider. Note that an endpoint of a walk is neither a collider, nor a noncollider. We stress that the property of being a collider/noncollider is relative to a walk (see also Figure 4).

Let \( \omega_1 = \langle \alpha, e_1^1, \gamma_1^1, \ldots, \gamma_{n-1}^1, e_n^1, \beta \rangle \) and \( \omega_2 = \langle \alpha, e_1^2, \gamma_1^2, \ldots, \gamma_{n-1}^2, e_n^2, \beta \rangle \) be two (nontrivial) walks. We say that they are endpoint-identical if \( e_1^1 \) and \( e_1^2 \) have the same mark at \( \alpha \) and \( e_n^1 \) and \( e_n^2 \) have the same mark at \( \beta \). Note that this may depend on the orientation of directed edges in the two walks. Assume that some edge \( e \) is between \( \alpha \) and \( \beta \). We say that the (nontrivial) walk \( \omega_1 \) is endpoint-identical to \( e \) if it is endpoint-identical to the walk \( \langle \alpha, e, \beta \rangle \). If \( \alpha = \beta \) and \( e \) is directed, this should hold for just one of the possible orientations of \( e \).

Let \( \omega_1 \) be a walk between \( \alpha \) and \( \gamma \), and \( \omega_2 \) a walk between \( \gamma \) and \( \beta \). The composition of \( \omega_1 \) with \( \omega_2 \) is the walk that starts at \( \alpha \), traverses every node and edge of \( \omega_1 \), and afterwards every node and edge of \( \omega_2 \), ending in \( \beta \). We say that we compose \( \omega_1 \) with \( \omega_2 \).

A directed path from \( \alpha \) to \( \beta \) is a path between \( \alpha \) and \( \beta \) consisting of edges of type \( \rightarrow \) only (possibly of length zero) such that they all point in the direction of \( \beta \). A cycle is either a loop, or a (nontrivial) path from \( \alpha \) to \( \beta \) composed with \( \beta \sim \alpha \). This means that in a cycle of length 2, an edge can be repeated. A directed cycle is either a loop, \( \alpha \rightarrow \alpha \), or a (nontrivial) directed path from \( \alpha \) to \( \beta \) composed with \( \beta \rightarrow \alpha \). For \( \alpha \in V \), we let An(\( \alpha \)) denote the set of ancestors, that is,

\[
\text{An}(\alpha) = \{ \gamma \in V \mid \text{there is a directed path from } \gamma \text{ to } \alpha \}.
\]

This is generalized to nonsingleton sets \( C \subseteq V \),

\[
\text{An}(C) = \bigcup_{\alpha \in C} \text{An}(\alpha).
\]

We stress that \( C \subseteq \text{An}(C) \) as we allow for trivial directed paths in the definition of an ancestor. We use the notation An\( C \) if we wish to emphasize in which graph the ancestry is read, but omit the subscript when no ambiguity arises.

Let \( G = (V, E) \) be a graph, and let \( O \subseteq V \). Define the subgraph induced by \( O \) to be the graph \( G_O = (O, E_O) \) where \( E_O \subseteq E \) is the set of edges that are between nodes in \( O \). If \( G_1 = (V, E_1) \) and \( G_2 = (V, E_2) \), we will write \( G_1 \subseteq G_2 \) to denote \( E_1 \subseteq E_2 \) and say that \( G_2 \) is a supergraph of \( G_1 \).

A directed graph (DG), \( D = (V, E) \), is a graph with only directed edges. Note that this also allows directed loops. Within a class of graphs, we define the complete graph to be the graph which is the supergraph of all graphs in the class when such a graph exists. For the class of DGs on node set \( V \), the complete graph is the graph with edge set \( E = \{(\alpha, \beta) \mid \alpha, \beta \in V \} \).

A directed acyclic graph (DAG) is a DG with no loops and no directed cycles. An acyclic directed mixed graph (ADMG) is a DMG with no loops and no directed cycles.
3. Directed mixed graphs and separation. In this section, we introduce $\mu$-separation for DMGs which are then shown to be closed under marginalization. In particular, we obtain a DMG representing the independence model arising from a local independence graph via marginalization.

The class of DMGs contains as a subclass the ADMGs that have no directed cycles [19, 32]. ADMGs have been used to represent marginalized DAG models, analogously to how we will use DMGs to represent marginalized DGs. ADMGs come with the $m$-separation criterion which can be extended to DMGs, but this criterion differs in important ways from the $\mu$-separation criterion introduced below. These differences also mean that our main result on Markov equivalence does not apply to, for example, DMGs with $m$-separation, and thus our theory of Markov equivalence hinges on the fact that we are considering DMGs using the asymmetric notion of $\mu$-separation.

3.1. $\mu$-separation. We define $\mu$-separation as a generalization of $\delta$-separation introduced by Didelez [11], analogously to how $m$-separation is a generalization of $d$-separation; see, for example, [33]. In Section A, we make the connection to Didelez’s $\delta$-separation exact and elaborate further on this in Section B.

**Definition 3.1 ($\mu$-connecting walk).** A nontrivial walk

$$\langle \alpha, e_1, \gamma_1, \ldots, \gamma_{n-1}, e_n, \beta \rangle$$

in $\mathcal{G}$ is said to be $\mu$-connecting (or simply open) from $\alpha$ to $\beta$ given $C$ if $\alpha \notin C$, every collider is in $\text{An}(C)$, no noncollider is in $C$, and $e_n$ has a head at $\beta$.

When a walk is not $\mu$-connecting given $C$, we say that it is closed or blocked by $C$. One should note that if $\omega$ is a $\mu$-connecting walk from $\alpha$ to $\beta$ given $C$, the inverse walk, $\omega^{-1}$, is not in general $\mu$-connecting from $\beta$ to $\alpha$ given $C$. The requirement that a $\mu$-connecting walk be nontrivial, that is, of strictly positive length, leads to the possibility of a node being separated from itself by some set $C$ when applying the following graph separation criterion to the class of DMGs.

**Definition 3.2 ($\mu$-separation).** Let $A, B, C \subseteq V$. We say that $B$ is $\mu$-separated from $A$ given $C$ if there is no $\mu$-connecting walk from any $\alpha \in A$ to any $\beta \in B$ given $C$ and write $A \perp_{\mu} B \mid C$, or write $A \perp_{\mu} B \mid C [\mathcal{G}]$ if we want to stress to what graph the separation statement applies.

The above notion of separation is given in terms of walks of which there are infinitely many in any DMG with a nonempty edge set. However, we will see that it is sufficient to consider a finite subset of walks from $A$ to $B$ (Proposition 3.5).

Given a DMG, $\mathcal{G} = (V, E)$, we define an independence model over $V$ using $\mu$-separation,

$$\mathcal{I}(\mathcal{G}) = \{ (A, B \mid C) \mid (A \perp_{\mu} B \mid C) \}.$$

Definition 3.1 implies $A \perp_{\mu} B \mid C$ whenever $A \subseteq C$ and, therefore, $\mathcal{I}(\mathcal{G}) \neq \emptyset$.

Below we state two propositions that essentially both give equivalent ways of defining $\mu$-separation. The propositions are useful when proving results on $\mu$-separation models.

**Proposition 3.3.** Let $\alpha, \beta \in V$, $C \subseteq V$. If there is a $\mu$-connecting walk from $\alpha$ to $\beta$ given $C$, then there is a $\mu$-connecting walk from $\alpha$ to $\beta$ that furthermore satisfies that every collider is in $C$. 

DEFINITION 3.4. A route from $\alpha$ to $\beta$ is a walk from $\alpha$ to $\beta$ such that no node different from $\beta$ occurs more than once, and $\beta$ occurs at most twice.

A route is always a path, a cycle or a composition of a path and a cycle that share no edge and only share the vertex $\beta$.

PROPOSITION 3.5. Let $\alpha, \beta \in V$, $C \subseteq V$. If $\omega$ is a $\mu$-connecting walk from $\alpha$ to $\beta$ given $C$, then there is a $\mu$-connecting route from $\alpha$ to $\beta$ given $C$ consisting of edges in $\omega$.

If there is a $\mu$-connecting walk from $A$ to $B$ given $C$, it does not in general follow that we can also find a $\mu$-connecting path or cycle from $A$ to $B$ given $C$. As an example of this, consider the following DMG on nodes $\{\alpha, \beta, \gamma\}$: $\alpha \leftarrow \beta \leftarrow \gamma$. There is a $\mu$-connecting walk from $\alpha$ to $\beta$ given $\emptyset$, and a $\mu$-connecting route, but no $\mu$-connecting path from $\alpha$ to $\beta$ given $\emptyset$.

3.2. Marginalization of DMGs. Given a DG or a DMG, $G$, we are interested in finding a graph that represents the marginal independence model over a node set $O \subseteq V$, that is, finding a graph $M$ such that

$$I(M) = (I(G))^O.$$  

It is well known that the class of DAGs with $d$-separation is not closed under marginalization, that is, for a DAG, $D = (V, E)$, and $O \subseteq V$, it is not in general possible to find a DAG with node set $O$ that encodes the same independence model among the variables in $O$ as did the original graph. Richardson and Spirtes [33] gave a concrete counterexample and in Example 3.7 we give a similar example to make the analogous point: DGs read with $\mu$-separation are not closed under marginalization. In this example, we use the following proposition which gives a simple characterization of separability in DGs.

PROPOSITION 3.6. Consider a DG, $D = (V, E)$, and let $\alpha, \beta \in V$. Then $\beta$ is $\mu$-separable (see Definition 2.5) from $\alpha$ in $D$ if and only if $\alpha \not\rightarrow_D \beta$.

EXAMPLE 3.7. Consider the directed graph, $G$, in Figure 5. We wish to show that it is not possible to encode the $\mu$-separations among nodes in $O = \{\alpha, \beta, \gamma, \delta\}$ using a DG on these nodes only. To obtain a contradiction, assume $D = (O, E)$ is a DG such that

$$A \perp_B C \mid D \iff A \perp_B C \mid G$$

for $A, B, C \subseteq O$. There is no $C \subseteq O \setminus \{\alpha\}$ such that $\alpha \perp_B \mid C \mid D$ and no $C \subseteq O \setminus \{\beta\}$ such that $\beta \perp_C \gamma \mid D$. If $D$ has the property (3.2), then it follows from Proposition 3.6 that $\alpha \rightarrow_D \beta$ and $\beta \rightarrow_D \gamma$. However, then $\gamma$ is not $\mu$-separated from $\alpha$ given $\emptyset$ in $D$. This shows that there exists no DG, $D$, that satisfies (3.2).

We note that marginalization of a probability model does not only impose conditional independence constraints on the observed variables but also so-called equality and inequality constraints; see, for example, [18] and references therein. In this paper, we will only be

\[
\alpha \quad \beta \quad \gamma \quad \delta
\]

FIG. 5. The directed graph of Example 3.7 which exemplifies that DGs are not closed under marginalization.
concerned with the graphical representation of local independence constraints, and not with representing analogous equality or inequality constraints.

In the remainder of this section, we first introduce the latent projection of a graph (see also [41] and [34]), and then show that it provides a marginalized DMG in the sense of (3.1). At the end of the section, we give an algorithm for computing the latent projection of a DMG. This algorithm is an adapted version of one described by Sadeghi [36] for a different class of graphs. Koster [25] described a similar algorithm for ADMGs.

**DEFINITION 3.8 (Latent projection).** Let $G = (V, E)$ be a DMG, $V = M \cup O$. We define the latent projection of $G$ on $O$ to be the DMG $(O, D)$ such that $\alpha \sim \beta \in D$ if and only if there exists an endpoint-identical (and nontrivial) walk between $\alpha$ and $\beta$ in $G$ with no colliders and such that every nonendpoint node is in $M$. Let $m(G, O)$ denote the latent projection of $G$ on $O$.

The definition of latent projection motivates the graphical term sibling for DMGs, as one way to obtain an edge $\alpha \leftrightarrow \beta$ is through a latent projection of a larger graph in which $\alpha$ and $\beta$ share a parent.

To characterize the class of graphs obtainable from a DG via a latent projection, we introduce the canonical DG of the DMG $G$, $C(G)$, as follows: for each (unordered) pair of nodes $\{\alpha, \beta\} \subseteq V$ such that $\alpha \leftrightarrow G \beta$, add a distinct auxiliary node, $m_{\{\alpha, \beta\}}$, add edges $m_{\{\alpha, \beta\}} \rightarrow \alpha$, $m_{\{\alpha, \beta\}} \rightarrow \beta$ to $E$ and then remove all bidirected edges from $E$. If $D$ is a DG, then $M = m(C, V)$ will satisfy

$$\alpha \leftrightarrow M \beta \iff \alpha \leftrightarrow M \alpha \quad \text{for all } \alpha, \beta \in O$$

for all subsets of vertices $O$. Conversely, if $G = (V, E)$ is a DMG that satisfies (3.3), then $G$ is the latent projection of its canonical DG; $m(C(G), V) = G$. The class of DMGs that satisfy (3.3) is closed under marginalization (Proposition 3.9) and has certain regularity properties (see, e.g., Proposition 3.10). These result provide the means for graphically representing marginals of local independence graphs. However, the theory that leads to our main results on Markov equivalence does not require the property (3.3) and, therefore, we develop it for general DMGs.

**PROPOSITION 3.9.** Let $O \subseteq V$. The graph $M = m(G, O)$ is a DMG. If $G$ satisfies (3.3), then $M$ does as well.

**PROPOSITION 3.10.** Assume that $G$ satisfies (3.3) and let $\alpha \in V$. Then $\alpha$ has no loops if and only if $\alpha \perp \mu_\alpha \mid V \setminus \{\alpha\}$.

We also observe directly from the definition that the latent projection operation preserves ancestry and nonancestry in the following sense.

**PROPOSITION 3.11.** Let $O \subseteq V$, $M = m(G, O)$ and $\alpha, \beta \in O$. Then $\alpha \in An_G(\beta)$ if and only if $\alpha \in An_M(\beta)$.

The main result of this section is the following theorem, which states that the marginalization defined by the latent projection operation preserves the marginal independence model encoded by a DMG.

**THEOREM 3.12.** Let $O \subseteq V$, $M = m(G, O)$. Assume $A, B, C \subseteq O$. Then

$$A \perp_\mu B \mid C \ [G] \iff A \perp_\mu B \mid C \ [M].$$
input: a DMG, \( G = (V, E) \) a subset \( M \subseteq V \) over which to marginalize

output: a graph \( M = (O, \bar{E}) \), \( O = V \setminus M \)

Initialize \( E_0 = E \), \( M_0 = (V, E_0) \), \( k = 0 \);

while \( \Omega_M(M_k) \neq \emptyset \) do

Choose \( \theta = \theta(\alpha, m, \beta) \) \( \in \Omega_M(M_k) \);

Set \( e_{k+1} \) to be the edge \( \alpha \sim \beta \) which is endpoint-identical to \( \theta \);

Set \( E_{k+1} = E_k \cup \{e_{k+1}\} \);

Set \( M_{k+1} = (V, E_{k+1}) \);

Update \( k = k + 1 \)

end

return \((M_k)_O\)

Algorithm 1: Computing the latent projection of a DMG

3.3. A marginalization algorithm. We describe an algorithm to compute the latent projection of a graph on some subset of nodes. For this purpose, we define a triroute, \( \theta \), to be a walk of length 2, \( \langle \alpha, e_1, \gamma, e_2, \beta \rangle \), such that \( \gamma \neq \alpha, \beta \). We suppress \( e_1 \) and \( e_2 \) from the notation and use \( \theta(\alpha, \gamma, \beta) \) to denote the triroute. We say that a triroute is colliding if \( \gamma \) is a collider on \( \theta \), and otherwise we say that it is noncolliding. This is analogous to the concept of a tripath (see, e.g., [26]), but allows for \( \alpha = \beta \).

Define \( \Omega_M(G) \) to be the set of noncolliding triroutes \( \theta(\alpha, m, \beta) \) such that \( m \in M \) and such that an endpoint-identical edge \( \alpha \sim \beta \) is not present in \( G \).

**Proposition 3.13.** Algorithm 1 outputs the latent projection of a DMG.

4. Properties of DMGs.

**Definition 4.1 (Markov equivalence).** Let \( G_1 = (V, E_1) \) and \( G_2 = (V, E_2) \) be DMGs. We say that \( G_1 \) and \( G_2 \) are Markov equivalent if \( \mathcal{I}(G_1) = \mathcal{I}(G_2) \). This defines an equivalence relation and we let \([G_1]\) denote the (Markov) equivalence class of \( G_1 \).

**Example 4.2 (Markov equivalence in DGs).** Let \( D = (V, E) \) be a DG. There is a directed edge from \( \alpha \) to \( \beta \) if and only if \( \beta \) cannot be separated from \( \alpha \) by any set \( C \subseteq V \setminus \{\alpha\} \) (Proposition 3.6). This implies that two DGs are Markov equivalent if and only if they are equal. Thus, in the restricted class of DGs, every Markov equivalence class is a singleton and in this sense identifiable from its induced independence model. However, when considering Markov equivalence in the more general class of DMGs not every equivalence class of a DG is a singleton as the DG might be Markov equivalent to a DMG. As an example of this, consider the complete DG on a node set \( V \) which is Markov equivalent to the complete DMG on \( V \).

**Definition 4.3 (Maximality of a DMG).** We say that \( G \) is maximal if it is complete, or if any added edge changes the induced independence model \( \mathcal{I}(G) \).

4.1. Inducing paths. Separability of nodes can be studied using the concept of an inducing path which has also been used in other classes of graphs [33, 41]. In the context of DMGs and \( \mu \)-separation, it is natural to define several types of inducing paths due to the asymmetry of \( \mu \)-separation and the possibility of directed cycles in DMGs.
DEFINITION 4.4 (Inducing path). An inducing path from \( \alpha \) to \( \beta \) is a nontrivial path or cycle, \( \pi = (\alpha, \ldots, \beta) \), which has a head at \( \beta \) and such that there are no noncolliders on \( \pi \) and every node is an ancestor of \( \alpha \) or \( \beta \). The inducing path \( \pi \) is bidirected if every edge on \( \pi \) is bidirected. If \( \pi \) is not bidirected, it has one of the forms \( \alpha \to \beta \) or \( \alpha \to \gamma_1 \leftrightarrow \cdots \leftrightarrow \gamma_n \leftrightarrow \beta \). and we say that it is unidirected. If, furthermore, \( \gamma_i \in \text{An}(\beta) \) for all \( i = 1, \ldots, n \) (or it is on the form \( \alpha \to \beta \)) then we say that it is directed.

Note that an inducing path is by definition either a path or a cycle. An inducing path is either bidirected or unidirected. Some unidirected inducing paths are also directed; see Figure 6 for examples. Propositions 4.7 and 4.8 show how bidirected and directed inducing paths in a certain sense correspond to bidirected and directed edges, respectively.

PROPOSITION 4.5. Let \( \nu \) be an inducing path from \( \alpha \) to \( \beta \). The following holds for any \( C \subseteq V \setminus \{\alpha\} \). If \( \alpha \neq \beta \), then there exists a \( \mu \)-connecting path from \( \alpha \) to \( \beta \) given \( C \). If \( \alpha = \beta \), then there exists a \( \mu \)-connecting cycle from \( \alpha \) to \( \beta \) given \( C \). We call such a path or cycle a \( \nu \)-induced open path or cycle, respectively, or simply a \( \nu \)-induced open walk to cover both the case \( \alpha = \beta \) and the case \( \alpha \neq \beta \). If the inducing path is bidirected or directed, then the \( \nu \)-induced open walk is endpoint-identical to the inducing path.

The following corollary is a direct consequence of Proposition 4.5, showing that \( \beta \) is inseparable from \( \alpha \) if there is an inducing path from \( \alpha \) to \( \beta \) irrespectively of whether the nodes are adjacent.

COROLLARY 4.6. Let \( \alpha, \beta \in V \). If there exists an inducing path from \( \alpha \) to \( \beta \) in \( G \), then \( \beta \) is not \( \mu \)-separated from \( \alpha \) given \( C \) for any \( C \subseteq V \setminus \{\alpha\} \), that is, \( \alpha \in u(\beta, \mathcal{I}(G)) \).

The following two propositions show that for two of the three types of inducing paths there is a Markov equivalent supergraph in which the nodes are adjacent. This illustrates how one can easily find Markov equivalent DMGs that do not have the same adjacencies. Example 4.12 shows that for a unidirected inducing path it may not be possible to add an edge without changing the independence model.

PROPOSITION 4.7. If there exists a bidirected inducing path from \( \alpha \) to \( \beta \) in \( G \), then adding \( \alpha \leftrightarrow \beta \) in \( G \) does not change the independence model.

PROPOSITION 4.8. If there exists a directed inducing path from \( \alpha \) to \( \beta \) in \( G \), then adding \( \alpha \to \beta \) in \( G \) does not change the independence model.

We say that nodes \( \alpha \) and \( \beta \) are collider-connected if there exists a nontrivial walk between \( \alpha \) and \( \beta \) such that every nonendpoint node is a collider on the walk. We say that \( \alpha \) is directly collider-connected to \( \beta \) if \( \alpha \) and \( \beta \) are collider-connected by a walk with a head at \( \beta \).
FIG. 7. A maximal DMG in which $\delta$ is inseparable from $\beta$, though no edge is between the two. See Example 4.12. We will in general omit the bidirected loops from the visual presentations of DMGs; see also the discussion in Section 5.4.

**DEFINITION 4.9.** Let $\alpha, \beta \in V$. We define the set

$$D(\alpha, \beta) = \{ \gamma \in \text{An}(\alpha, \beta) \mid \gamma \text{ is directedly collider-connected to } \beta \} \setminus \{ \alpha \}.$$  

Note that if $\alpha \not\rightarrow_{\mathcal{G}} \beta$, then $\text{pa}(\beta) \subseteq D(\alpha, \beta)$, and if the graph is furthermore a directed graph then $\text{pa}(\beta) = D(\alpha, \beta)$.

**PROPOSITION 4.10.** If there is no inducing path from $\alpha$ to $\beta$ in $\mathcal{G}$, then $\beta$ is separated from $\alpha$ by $D(\alpha, \beta)$.

**EXAMPLE 4.11 (Inducing paths).** Consider the DMG on nodes $\{\alpha, \gamma\}$ and with a single edge $\gamma \rightarrow \alpha$. In this case, there is no inducing path from $\alpha$ to $\alpha$ and $\alpha$ is $\mu$-separated from $\alpha$ by $D(\alpha, \alpha) = \{\gamma\}$. Now add the edge $\alpha \leftrightarrow \gamma$. In this new DMG, there is an inducing path from $\alpha$ to $\alpha$ and therefore $\alpha$ is inseparable from itself.

**EXAMPLE 4.12 (Nonadjacency of inseparable nodes in a maximal DMG).** Consider the DMG in Figure 7. One can show that this DMG is maximal (Definition 4.3). There is an inducing path from $\beta$ to $\delta$ making $\delta$ inseparable from $\beta$, yet no arrow can be added between $\beta$ and $\delta$ without changing the independence model. This example illustrates that maximal DMGs do not have the property that inseparable nodes are adjacent. This is contrary to MAGs which form a subclass of ancestral graphs and have this exact property [33].

5. Markov equivalence of DMGs. The main result of this section is that each Markov equivalence class of DMGs has a greatest element, that is, an element which is a supergraph of all other elements. This fact is helpful for understanding and graphically representing such equivalence classes, and potentially also for constructing learning algorithms. We will prove this result by arguing that the independence model of a DMG, $\mathcal{G} = (V, E)$, defines for each node $\alpha \in V$ a set of potential parents and a set of potential siblings. We then construct the greatest element of $[\mathcal{G}]$ by simply using these sets, and argue that this is in fact a Markov equivalent supergraph. As we only use the independence model to define the sets of potential parents and siblings, the supergraph is identical for all members of $[\mathcal{G}]$, and thus a greatest element. Within the equivalence class, the greatest element is also the only maximal element, and we will refer to it as the maximal element of the equivalence class.

5.1. Potential siblings.

**DEFINITION 5.1.** Let $\mathcal{I}$ be an independence model over $V$ and let $\alpha, \beta \in V$. We say that $\alpha$ and $\beta$ are potential siblings in $\mathcal{I}$ if (s1)-(s3) hold:

(s1) $\beta \in u(\alpha, \mathcal{I})$ and $\alpha \in u(\beta, \mathcal{I})$,
(s2) for all $\gamma \in V, C \subseteq V$ such that $\beta \in C$,

$$\langle \gamma, \alpha \mid C \rangle \in \mathcal{I} \Rightarrow \langle \gamma, \beta \mid C \rangle \in \mathcal{I}.$$
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(s3) for all \( \gamma \in V, C \subseteq V \) such that \( \alpha \in C \),
\[
\langle \gamma, \beta | C \rangle \in \mathcal{I} \implies \langle \gamma, \alpha | C \rangle \in \mathcal{I}.
\]

Potential siblings are defined abstractly above in terms of the independence model only. The following proposition gives a useful characterization for graphical independence models by simply contraposing (s2) and (s3).

**Proposition 5.2.** Let \( \mathcal{I}(G) \) be the independence model induced by \( G \). Then \( \alpha, \beta \in V \) are potential siblings if and only if (gs1)–(gs3) hold:

- (gs1) \( \beta \in u(\alpha, \mathcal{I}(G)) \) and \( \alpha \in u(\beta, \mathcal{I}(G)) \),
- (gs2) for all \( \gamma \in V, C \subseteq V \) such that \( \beta \in C \): if there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \beta \) given \( C \), then there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \alpha \) given \( C \),
- (gs3) for all \( \gamma \in V, C \subseteq V \) such that \( \alpha \in C \): if there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \alpha \) given \( C \), then there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \beta \) given \( C \).

**Proposition 5.3.** Assume that \( \alpha \leftrightarrow \beta \) is in \( G \). Then \( \alpha \) and \( \beta \) are potential siblings in \( \mathcal{I}(G) \).

**Lemma 5.4.** Assume that \( \alpha \) and \( \beta \) are potential siblings in \( \mathcal{I}(G) \). Let \( G^+ \) denote the DMG obtained from \( G \) by adding \( \alpha \leftrightarrow \beta \). Then \( \mathcal{I}(G) = \mathcal{I}(G^+) \).

The above shows that if \( \alpha \) and \( \beta \) are potential siblings in \( \mathcal{I}(G) \) then there exists a supergraph, \( G^+ \), which is Markov equivalent with \( G \), such that \( \alpha \) and \( \beta \) are siblings in \( G^+ \). This motivates the term potential siblings.

5.2. Potential parents. In this section, we will argue that also a set of potential parents are determined by the independence model. This case is slightly more involved for two reasons. First, the relation is asymmetric, as for each potential parent edge there is a parent node and a child node. Second, adding directed edges potentially changes the ancestry of the graph.

**Definition 5.5.** Let \( \mathcal{I} \) be an independence model over \( V \) and let \( \alpha, \beta \in V \). We say that \( \alpha \) is a potential parent of \( \beta \) in \( \mathcal{I} \) if (p1)–(p4) hold:

- (p1) \( \alpha \in u(\beta, \mathcal{I}) \),
- (p2) for all \( \gamma \in V, C \subseteq V \) such that \( \alpha \notin C \),
\[
\langle \gamma, \beta | C \rangle \in \mathcal{I} \implies \langle \gamma, \alpha | C \rangle \in \mathcal{I},
\]
- (p3) for all \( \gamma, \delta \in V, C \subseteq V \) such that \( \alpha \notin C, \beta \in C \),
\[
\langle \gamma, \delta | C \rangle \in \mathcal{I} \implies \langle \gamma, \beta | C \rangle \in \mathcal{I} \lor \langle \alpha, \delta | C \rangle \in \mathcal{I},
\]
- (p4) for all \( \gamma \in V, C \subseteq V \), such that \( \alpha \notin C \),
\[
\langle \beta, \gamma | C \rangle \in \mathcal{I} \implies \langle \beta, \gamma | C \cup \{\alpha\} \rangle \in \mathcal{I}.
\]

**Proposition 5.6.** Let \( \mathcal{I}(G) \) be the independence model induced by \( G \). Then \( \alpha \in V \) is a potential parent of \( \beta \in V \) if and only if (gp1)–(gp4) hold:

- (gp1) \( \alpha \in u(\beta, \mathcal{I}(G)) \),
- (gp2) for all \( \gamma \in V, C \subseteq V \) such that \( \alpha \notin C \): if there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \alpha \) given \( C \), then there exists a \( \mu \)-connecting walk from \( \gamma \) to \( \beta \) given \( C \),
for all $\gamma, \delta \in V$, $C \subseteq V$ such that $\alpha \notin C$, $\beta \in C$: if there exists a $\mu$-connecting walk from $\gamma$ to $\beta$ given $C$ and a $\mu$-connecting walk from $\alpha$ to $\delta$ given $C$, then there exists a $\mu$-connecting walk from $\gamma$ to $\delta$ given $C$,

for all $\gamma \in V$, $C \subseteq V$, such that $\alpha \notin C$: if there exists a $\mu$-connecting walk from $\beta$ to $\gamma$ given $C \cup \{\alpha\}$, then there exists a $\mu$-connecting walk from $\beta$ to $\gamma$ given $C$.

**Proposition 5.7.** Assume that $\alpha \rightarrow \beta$ is in $\mathcal{G}$. Then $\alpha$ is a potential parent of $\beta$ in $\mathcal{I}(\mathcal{G})$.

**Lemma 5.8.** Assume that $\alpha$ is a potential parent of $\beta$ in $\mathcal{I}(\mathcal{G})$. Let $\mathcal{G}^+$ denote the DMG obtained from $\mathcal{G}$ by adding $\alpha \rightarrow \beta$. Then $\mathcal{I}(\mathcal{G}) = \mathcal{I}(\mathcal{G}^+)$.

**5.3. A Markov equivalent supergraph.** Let $\mathcal{G} = (V, E)$ be a DMG. Define $\mathcal{N}(\mathcal{I}(\mathcal{G})) = (V, E^m)$ to be the DMG with edge set $E^m = E^d \cup E^b$ where $E^d$ is a set of directed edges and $E^b$ a set of bidirected edges such that the directed edge from $\alpha$ to $\beta$ is in $E^d$ if and only if $\alpha$ is a potential parent of $\beta$ in $\mathcal{I}(\mathcal{G})$ and the bidirected edge between $\alpha$ and $\beta$ is in $E^b$ if and only if $\alpha$ and $\beta$ are potential siblings in $\mathcal{I}(\mathcal{G})$.

**Theorem 5.9.** Let $\mathcal{N} = \mathcal{N}(\mathcal{I}(\mathcal{G}))$. Then $\mathcal{N} \in [\mathcal{G}]$ and $\mathcal{N}$ is a supergraph of all elements of $[\mathcal{G}]$. Furthermore, if we have a finite sequence of DMGs $\mathcal{G}_0, \mathcal{G}_1, \ldots, \mathcal{G}_m, \mathcal{G}_i = (V, E_i)$, such that $\mathcal{G}_0 = \mathcal{G}$, $\mathcal{G}_m = \mathcal{N}$, and $E_i \subseteq E_{i+1}$ for all $i = 0, \ldots, m - 1$, then $\mathcal{G}_i$ is Markov equivalent with $\mathcal{N}$ for all $i = 0, \ldots, m - 1$.

The graph $\mathcal{N}$ in the above theorem is a supergraph of every Markov equivalent DMG and, therefore, maximal. On the other hand, every maximal DMG is a representative of its equivalence class, and also a supergraph of all Markov equivalent DMGs. This means that we can use the class of maximal DMGs to obtain a unique representative for each DMG equivalence class.

Lemmas 5.4 and 5.8 show that conditions (gs1)–(gs3) and (gp1)–(gp4) are sufficient to Markov equivalently add a bidirected or a directed edge, respectively. The conditions are also necessary in the sense that for each condition one can find example graphs where only a single condition is violated and where the larger graph is not Markov equivalent to the smaller graph.

We can note that $\alpha$ is a potential parent and a potential sibling of $\alpha$ if and only if $\alpha \in \mu(\alpha, \mathcal{I}(\mathcal{G}))$. This means that in $\mathcal{N}(\mathcal{I}(\mathcal{G}))$ for each node either both a directed and a bidirected loop is present or no loop at all.

**5.4. Directed mixed equivalence graphs.** Theorem 5.9 suggests that one can represent an equivalence class of DMGs by displaying the maximal element and then simply indicate which edges are not present for all members of the equivalence class.

**Definition 5.10 (DMEG).** Let $\mathcal{N} = (V, F)$ be a maximal DMG. Define $\tilde{F} \subseteq F$ such that for $e \in F$ we let $e \in \tilde{F}$ if and only if there exists a DMG $\mathcal{G} = (V, \tilde{F})$ such that $\mathcal{G} \in [\mathcal{N}]$ and $e \notin \tilde{F}$. We call $\mathcal{N}' = (V, F, \tilde{F})$ a directed mixed equivalence graph (DMEG). When visualizing $\mathcal{N}'$, we draw $\mathcal{N}$, but use dashed edges for the set $\tilde{F}$; see Figure 8.

Let $\mathcal{N}' = (V, F, \tilde{F})$ be a DMEG. The DMG $(V, F)$ is in the equivalence class represented by $\mathcal{N}'$. However, one cannot necessarily remove any subset of $\tilde{F}$ and obtain a member of the Markov equivalence class (see Figure 8). Moreover, an equivalence class does not in general contain a least element, that is, an element which is a subgraph of all Markov equivalent graphs.
We will throughout this section let $\mathcal{N} = (V, F)$ be a maximal DMG. For $e \in F$, we will use $\mathcal{N} - e$ to denote the graph $(V, F \setminus \{e\})$. Assume that we have a maximal DMG from which we wish to derive the DMEG. Consider some edge $e \in F$. If $\mathcal{N} - e \in [\mathcal{N}]$, then $e \in \tilde{F}$ as there exists a Markov equivalent subgraph of $\mathcal{N}$ in which $e$ is not present. On the other hand, if $\mathcal{N} - e \notin [\mathcal{N}]$ then we note that $\mathcal{N} - e$ is the largest subgraph of $\mathcal{N}$ that does not contain $e$. Let $\mathcal{K}$ be a subgraph of $\mathcal{N}$ that does not contain $e$. Then $\mathcal{I}(\mathcal{N}) \subseteq \mathcal{I}(\mathcal{N} - e) \subseteq \mathcal{I}(\mathcal{K})$. Using Theorem 5.9, we know that all $\mathcal{N}$-Markov equivalent DMGs are in fact subgraphs of $\mathcal{N}$, and using that $\mathcal{K}$ is not Markov equivalent to $\mathcal{N}$ we see that all graphs in $[\mathcal{N}]$ must contain $e$. This means that when $\mathcal{N} - e \notin [\mathcal{N}]$ then $e \notin \tilde{F}$ as $e$ must be present in all Markov equivalent DMGs.

Any loop should in principle be dashed when drawing a DMEG as for each node in a maximal DMG either both the directed and the bidirected loop are present or neither of them. However, we choose to not present them as dashed as if they are present in the maximal DMG, then at least one of them will be present in any Markov equivalent DMG satisfying (3.3), that is, for any DMG which is a marginalization of a DG. In addition, we only draw the directed loop to not overload the visualizations.

5.5. Constructing a directed mixed equivalence graph. When constructing a DMEG from $\mathcal{N}$, it suffices to consider the graphs $\mathcal{N} - e$ for each $e \in E$ and determine if they are Markov equivalent to $\mathcal{N}$ or not. A brute-force approach to doing so is to simply check all separation statements in both graphs. However, one can make a considerably more efficient algorithm.

PROPOSITION 5.11. Assume $\alpha \xleftrightarrow{\mathcal{N}} \beta$. It holds that $\mathcal{N} - e \in [\mathcal{N}]$ if and only if $\alpha \in u(\beta, \mathcal{I}(\mathcal{N} - e))$.

PROPOSITION 5.12. Assume $\alpha \xleftrightarrow{\mathcal{N}} \beta$. Then $\mathcal{N} - e \in [\mathcal{N}]$ if and only if $\alpha \in u(\beta, \mathcal{I}(\mathcal{N} - e))$ and $\beta \in u(\alpha, \mathcal{I}(\mathcal{N} - e))$.

Fig. 8. The DMG 1 is maximal (the bidirected loops at $\alpha$, $\beta$ and $\delta$ have been omitted from the visual presentation). The DMGs 1–6 are the six elements of its Markov equivalence class (when ignoring Markov equivalent removal of loops). The graph 7 is the corresponding DMEG. In a DMEG, every solid edge is in every graph in the equivalence class, every absent edge is not in any graph, and every dashed edge is in some, but not in others. Note that every DMG in the above equivalence class contains the edge $\gamma \rightarrow \beta$ or the edge $\delta \rightarrow \beta$ even though both are dashed in the DMEG. This example shows that not every equivalence class contains a least element.
We can now outline a two-step algorithm for constructing the DMEG from an arbitrary DMG, \( \mathcal{G} \). We first construct the maximal Markov equivalent graph, \( \mathcal{N} \). We know from Theorem 5.9 that one can simply check if each pair of nodes are potential siblings/parents in the independence model induced by \( \mathcal{G} \) and construct the maximal Markov equivalent graph directly. This may, however, not be computationally efficient.

The above propositions show that given the maximal DMG, one can efficiently construct the DMEG by evaluating separability once for each directed edge and twice for each bidirected edge. Using Proposition 4.10, one can determine separability by testing a single separation statement, and this means that starting from \( \mathcal{N} \), one can construct the corresponding DMEG in a way such that the number of separation statements to test scales linearly in the number of edges in \( \mathcal{N} \).

EXAMPLE 5.13 (Gateway drugs, continued). We return to the model in Example 2.3 to consider what happens when it is only partially observed and to give an interpretation of the corresponding local independence model. The local independence graph is assumed to be as depicted on Figure 9, left.

Consider first the situation where \( L \) and \( I \) are unobserved. In this case, under the faithfulness assumption of the full model (Definition C.5) we can construct the DMEG, which is shown in the center panel of Figure 9, from the local independence model. The DMEG represents the Markov equivalence class which we can infer from the marginal local independence model (\( L \) and \( I \) are unobserved). Theoretically, the inference requires an oracle to provide us with local independence statements, which will in practice have to be approximated by statistical tests. What is noteworthy is that the DMEG can be inferred from the distribution of the observed variables only, and we do not need to know the local independences of the full model.

If we ignore which edges are dashed and which are not, the graph simply represents the local independence model of the marginal system as the maximal element in the Markov equivalence class. The dashed edges give us additional—and in some sense local—information. As an example, the directed edge from \( A \) to \( H \) is dashed and we cannot know if there exists a conditioning set that would render \( H \) locally independent of \( A \) in the full system. On the other hand, the directed edge from \( T \) to \( H \) is absent, and we can conclude that tobacco use is not directly affecting hard drug use.

Consider instead the situation where \( I \) is also observed. \( I \) serves as an analogue to an instrumental variable (see, e.g., [31] for an introduction to instrumental variables). The inclusion of this variable identifies some of the structure by removing some dashed edges and making others nondashed.

6. Discussion and conclusion. In this paper, we introduced a class of graphs to represent local independence structures of partially observed multivariate stochastic processes.
Previous work based on directed graphs, that allows for cycles and use the asymmetric $\delta$-separation criterion, was extended to mixed directed graphs to account for latent processes and we introduced $\mu$-separation in mixed directed graphs.

An important task is the characterization of equivalence classes of graphs and this has been studied, for example, in MAGs [5, 45]. In the case of MAGs, a key result is that every element in a Markov equivalence class has the same skeleton, that is, the same adjacencies [5]. As shown by Propositions 4.7 and 4.8, this is not the case for DMGs, and Example 4.12 shows that one cannot necessarily within a Markov equivalence class find an element such that two nodes are inseparable if and only if they are adjacent.

We proved instead a central maximality property which allowed us to propose the use of DMEGs to represent a Markov equivalence class of DMGs in a concise way. Given a maximal DMG, we furthermore argued that one can efficiently find the DMEG. Similar results are known for chain graphs, as one can also in a certain sense find a unique, largest graph representing a Markov equivalence class [20], though this graph is not a supergraph of all Markov equivalent graphs as in the case of DMGs. Volf and Studený [42] suggested to use this largest graph as a unique representative of the Markov equivalence class, and they provided an algorithm to construct it.

We emphasize that the characterization given of the maximal element of a Markov equivalence class of DMGs is constructive in the sense that it straightforwardly defines an algorithm for learning a maximal DMG from a local independence oracle. This learning algorithm may not be computationally efficient or even feasible for large graphs, and it is ongoing research to develop efficient learning algorithms and to develop the practical implementations of the tools needed for replacing the oracle by statistical tests.

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SUPPLEMENTARY MATERIAL

Additional results and proofs (DOI: 10.1214/19-AOS1821SUPP; .pdf). The supplementary material consists of Sections A to F. In Sections A and B, we relate $\mu$-separation to Didelez’s $\delta$-separation, and also relate our slightly different definitions of local independence. Section C describes how one can unroll a local independence graph and obtain a DAG. We use this to discuss Markov properties and faithfulness in the time series case. In Section D, we provide an augmentation criterion to determine $\mu$-separation using an auxiliary undirected graph. In Section E, we discuss conditions for existence of compensators and elaborate on the definition of local independence. Section F contains the proofs of the main paper.

REFERENCES


