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Sadeghimanesh, Amir Hossein; Feliu, Elisenda

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GRÖBNER BASES OF REACTION NETWORKS WITH INTERMEDIATE SPECIES

AMIRHOSEIN SADEGHIMANESH\textsuperscript{1}, ELISENDA FELIU\textsuperscript{1,2}

ABSTRACT. In this work we consider the computation of Gröbner bases of the steady state ideal of reaction networks equipped with mass-action kinetics. Specifically, we focus on the role of intermediate species and the relation between the extended network (with intermediate species) and the core network (without intermediate species).

We show that a Gröbner basis of the steady state ideal of the core network always lifts to a Gröbner basis of the steady state ideal of the extended network by means of linear algebra, with a suitable choice of monomial order. As illustrated with examples, this contributes to a substantial reduction of the computation time, due mainly to the reduction in the number of variables and polynomials. We further show that if the steady state ideal of the core network is binomial, then so is the case for the extended network, as long as an extra condition is fulfilled. For standard networks, this extra condition can be visually explored from the network structure alone.

Keywords: binomial ideals, mass-action kinetics, steady state ideal, invariant, Gröbner basis

INTRODUCTION

Parametric polynomial systems of equations arise in the natural sciences when modeling ecosystems, cell behavior, the spread of an illness, and molecular interactions within the cell, to name a few examples. In these scenarios questions of interest often boil down to describing the solutions to these systems for varying values of the parameters. Although only non-negative solutions are typically meaningful, the standard tool in computational algebraic geometry to study algebraic varieties, namely Gröbner bases, has proven useful. However, due to the parametric coefficients, the computation of a reduced Gröbner basis can be time consuming for realistic examples, which typically involve many variables and parameters. The computation time depends mainly on the degree of the polynomials, the number of variables and coefficients, the choice of the monomial order and the used method \cite{1,2,5,9,29}. These universal considerations target generic polynomial systems, but, in applications, the structure of the particular system of interest might favor one method or one monomial order over another.

\textsuperscript{1}Department of Mathematical Sciences, University of Copenhagen, Universitetsparken 5, 2100 Copenhagen, Denmark.
\textsuperscript{2}Corresponding author: efeliu@math.ku.dk
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We focus on a specific type of polynomial systems that arise when modeling chemical reaction networks with mass-action kinetics \[10, 12\]. Specifically, the evolution of the concentrations of the species of a chemical reaction network in time is described under mass-action by a system of ordinary differential equations in \( \mathbb{R}^n \)

\[
\frac{dx_i}{dt} = f_{\kappa,i}(x), \quad i = 1, \ldots, n
\]

with \( f_{\kappa,i}(x) \) polynomial. The monomials of each \( f_{\kappa,i}(x) \) depend on the reaction network structure alone, and the coefficients depend on the reaction rate constants \( \kappa \), which are often unknown and thus treated as parameters. The steady states, or equilibrium points, of the system are the non-negative points of the variety defined by the steady state ideal \( I_{\kappa} = \langle f_{\kappa,1}(x), \ldots, f_{\kappa,n}(x) \rangle \).

The question of restricting to non-negative steady states remains challenging and no straightforward solutions have been proposed. Despite of this, Gröbner bases have been for example used for model discrimination \[13, 14, 18–20\]. They are also used to decide whether the steady state ideal is binomial, that is, whether any reduced Gröbner basis consists of polynomials with at most two terms. If this is the case, then methods to detect the existence of multiple steady states can be applied \[22, 23\].

In this work we exploit the specific structure of the steady state ideal, which reflects the structure of the reaction network, to guide the selection of good monomial orders and to compute reduced Gröbner bases faster. Specifically, we consider a frequent and nicely-behaved class of species introduced in \[11\] called intermediate species (or intermediates, for short). Intermediates give rise to linear terms in the steady state polynomials, and they can be removed from a reaction network resulting in a smaller core network with only the non-intermediates. A key property is that steady states of the core network can be lifted to steady states of the extended network.

The first main result of this work is Theorem 3.4, where we show how to obtain a Gröbner basis of the extended network from one of the core network using linear algebra. The result implicitly gives good monomial orders, namely, those for which the concentration of the intermediates are larger than for the non-intermediates, and are lexicographic in the variables corresponding to the intermediates. Example 3.5 illustrates the computational advantage of using our approach compared with other methods. Additionally, we conclude that the analysis of the steady state ideal of the core network is sufficient for model discrimination.

The second main result, Theorem 3.10, addresses how to decide whether the steady state ideal is binomial. We show that if the steady state ideal of the core network is binomial, then this is also the case for the steady state ideal of the extended network provided an extra condition is fulfilled. In typical networks, this extra condition can be readily checked from the network structure alone. When the core network has a homogeneous steady state ideal (which happens frequently for realistic reaction networks), then one can employ the linear algebra-based method introduced in \[4\] to detect whether the steady state ideal of the core network is binomial. Then, combined with our result, we obtain a faster method to address whether the steady state ideal of the original network is binomial, which does not rely on the computation of a Gröbner basis.
The key property behind our results is that intermediates define a square linear subsystem of full rank among the steady state polynomials. Its solution and posterior substitution into the remaining polynomials gives rise to a smaller ideal in the non-intermediates. A Gröbner basis of the small ideal can then be lifted to a Gröbner basis of the original ideal. Our approach can be theoretically applied to arbitrary parametric ideals, after detection of linear subsystems among a set of generators. However, technical conditions that are necessary for our results to hold might not be straightforward to check, since we overcome this difficulty by exploiting the network structure.

The structure of the paper is as follows. We start by introducing reaction networks and basic concepts such as the steady state ideal. Intermediates are introduced in Section 2. In Section 3 we address Gröbner bases of networks with intermediates, discuss binomial steady state ideals and relate our work to [4]. In Section 4 a technical condition of algebraic independence of a set of rational functions, which is assumed in the former sections, is discussed. Finally, in the last section, we discuss another class of special species, namely enzymes, that might lead to similar results concerning the computation of Gröbner bases.

1. The steady state ideal of a reaction network

We follow the formalism of [11]. See also [10, 12] for an introduction to reaction networks. Subscripts $\geq 0, > 0$ on $\mathbb{R}$ (resp. $\mathbb{Z}$) refer to the non-negative and positive real numbers (resp. integer numbers).

A reaction network is an ordered triple $N = (S, C, R)$ where $S, C$ and $R$ are three sets called the set of species, complexes and reactions, respectively. Here $S$ is a finite set and $C$ is a finite set of linear combinations of elements of $S$ with coefficients in $\mathbb{Z}_{\geq 0}$. A reaction is an ordered pair of complexes $(c, c')$ in $C^2$, usually denoted as $c \rightarrow c'$. For the reaction $c \rightarrow c'$, the complex $c$ is called the reactant and $c'$ is called the product.

A digraph is associated with a reaction network as follows. The vertex set is $C$ and there is a directed edge from the reactant to the product of every reaction. If both reactions $c \rightarrow c'$ and $c' \rightarrow c$ for two complexes $c$ and $c'$ exist, then the notation $c \rightleftharpoons c'$ is used and the reaction is said to be reversible.

Complexes that are not part of any reaction or species that are not part of any complex do not appear in the digraph. Therefore, the reaction network cannot uniquely be determined from the digraph alone. For simplicity, however, we often introduce a reaction network by its digraph and implicitly assume that the set of complexes equals the set of vertices and the set of species consists of the species that appear in at least one complex.

Write $S = \{X_1, \ldots, X_n\}$, such that the set of species is implicitly ordered. Then a complex $c$ is of the form $c_1X_1 + \cdots + c_nX_n$, which we also write in vector form as $c = (c_1, c_2, \ldots, c_n) \in \mathbb{Z}_n$. With this representation, $c_i$ is called the stoichiometric coefficient of $X_i$ in $c$.

Example 1.1. Let $S = \{X_1, X_2, X_3, X_4\}$, $C = \{X_1 + X_3, X_4, X_2 + X_3\}$, $R = \{X_1 + X_3 \rightarrow X_4, X_4 \rightarrow X_1 + X_3, X_4 \rightarrow X_2 + X_3\}$. The network $N = (S, C, R)$ is
represented with the following digraph

\[ X_1 + X_3 \rightleftharpoons X_4 \rightarrow X_2 + X_3. \]

The complexes \( X_1 + X_3 \) and \( X_4 \) appear both as reactants and products while \( X_2 + X_3 \) appears only as a product.

We next construct a system of Ordinary Differential Equations (ODEs) that models the variation of the concentration of each species in time and introduce the relevant polynomials \( F_i(x) \) that are the focus of this work. We denote the concentration of each species \( X_i \) in lower-case \( x_i \). For each reaction \( c \rightarrow c' \), we introduce a parameter \( k_{c \rightarrow c'} \), and a polynomial \( F_i(x) \) is associated with every species \( X_i \) as follows:

\[ F_i(x) = \sum_{c \rightarrow c' \in \mathcal{R}} (c'_i - c_i)k_{c \rightarrow c'} x^c \in \mathbb{R}(k)[x], \]

where \( x^c = x_1^{c_1} \cdots x_n^{c_n} \). Here \( x = (x_1, \ldots, x_n) \) and \( \mathbb{R}(k) \) is the field of rational functions with variables \( k_{c \rightarrow c'} \) and real coefficients. The symbol \( k \) stands for the parameter vector \( k = (k_{c \rightarrow c'} | c \rightarrow c' \in \mathcal{R}) \).

For a chosen positive value \( k^* \in \mathbb{R}_{>0}^n \) of the parameter vector, we let \( F_{k^*,i}(x) \in \mathbb{R}[x] \) denote the image of \( F_i(x) \) under the evaluation map

\[ \mathbb{R}[k] \to \mathbb{R}, \quad k_{c \rightarrow c'} \mapsto k^*_{c \rightarrow c'}. \]

With this choice of \( k^* \), the ODE system of the reaction network under mass-action kinetics is

\[ \dot{x}_i = F_{k^*,i}(x), \quad i = 1, \ldots, n, \quad x \in \mathbb{R}_{\geq 0}^n. \]

The value \( k^*_{c \rightarrow c'} > 0 \) is called the reaction rate constant of \( c \rightarrow c' \) and is usually depicted as a label of the reaction in the associated digraph. By [27], if the starting condition of \( \dot{x} = F_{k^*,i}(x) \) belongs to \( \mathbb{R}_{>0}^n \) (resp. \( \mathbb{R}_{\geq 0}^n \)), then so does the trajectory for all positive times in the interval of definition.

The steady states of the network are the common zeros of \( F_{k^*,i}(x) \), \( i = 1, \ldots, n \). In applications, only non-negative real solutions have meaning and mostly, positive steady states are interesting, meaning all concentrations are positive. Since the values of the reaction rate constants are in general unknown, they are treated as parameters of the system. Thus we aim at studying the zeros of the system of polynomials \( F_i(x) = 0 \), for \( i = 1, \ldots, n \) in \( \mathbb{R}(k) \) and specially the positive zeros after specifying values for \( k \).

**Definition 1.2.** Let \( \mathcal{N} = (\mathcal{S}, \mathcal{C}, \mathcal{R}) \) be a reaction network with \( \mathcal{S} = \{X_1, \ldots, X_n\} \).

(a) \( F_i(x) \in \mathbb{R}(k)[x] \) is called the steady state polynomial of \( X_i \).

(b) The ideal generated by the steady state polynomials of all the species in the network in the ring \( \mathbb{R}(k)[x] \) is called the steady state ideal of the network:

\[ I_\mathcal{N} = \langle F_i(x) \mid i = 1, \ldots, n \rangle \subseteq \mathbb{R}(k)[x]. \]
The set of steady states for a vector of reaction rate constants $k^*$ is thus the solution set to any basis (set of generators) of $I_N$ specialized to $k^*$.

It follows from (1) and (2) that for all $x \in \mathbb{R}^n$, the vector $F_k(x) = (F_{k,1}(x), \ldots, F_{k,n}(x))$

lies in the vector subspace $S = \text{span}\{(c - c' | c \rightarrow c' \in \mathbb{R})\} \subseteq \mathbb{R}^n$. If $s = \text{dim}(S)$, then $n - s$ of the steady state polynomials can be written as linear combinations of the remaining $s$ polynomials. We conclude that it is always possible to find a basis of $I_N$ with cardinality $\text{dim}(S)$.

Example 1.3. (continued from Example 1.1) The ODE system of the reaction network with digraph

$X_1 + X_3 \xrightarrow{k_1 \ k_2} X_4 \xrightarrow{k_3} X_2 + X_3$

is

$\dot{x}_1 = -k_1 x_1 x_3 + k_2 x_4 \quad \dot{x}_2 = k_3 x_4$

$\dot{x}_3 = -k_1 x_1 x_3 + k_2 x_4 + k_3 x_4 \quad \dot{x}_4 = k_1 x_1 x_3 - k_2 x_4 - k_3 x_4$.

In this case $\text{dim}(S) = 2$, $k = (k_1, k_2, k_3)$ and the steady state ideal is

$I_N = \langle -k_1 x_1 x_3 + k_2 x_4, k_3 x_4 \rangle \subseteq \mathbb{R}(k)[x]$.

2. Intermediates and steady states

In this subsection we introduce a special type of species of interest: intermediates.

Definition 2.1. We say that $Y \subseteq S$ is a subset of intermediates if each $Y \in Y$ fulfills:

- $Y \in \mathcal{C}$ and the stoichiometric coefficient of $Y$ in all other complexes is zero, and
- there exists at least one reaction having $Y$ as reactant and at least one reaction having $Y$ as product.

Each $Y \in Y$ is called an intermediate.

Whenever a set of intermediates $Y$ is given, we partition the set of species into two disjoint subsets $Y = \{Y_1, \ldots, Y_m\}$ and $X = \{X_1, \ldots, X_n\}$ of non-intermediates. We assume further that the set of species is ordered such that the species $Y_1, \ldots, Y_m$ are first. With this convention, we let $(y, x)$ denote the concentration vector of all species: $x$ is the concentration vector of the species in $X$ and $y$ of the species in $Y$.

A complex is either an intermediate in $Y$ or it contains only non-intermediates. In the latter case we say that $c$ is a non-intermediate complex.

Note that given $Y$, we refer to the intermediates of the network as the species in $Y$, even though there might be other species in $X$, regarded as non-intermediates, that fulfill the two items in Definition 2.1.

Example 2.2. The most common mechanism involving intermediates is of the following form:

$X + E \rightarrow Y \rightarrow X' + E$
or variations of it by letting one or both reactions being reversible. *Isomerism mechanisms* among intermediates are also common:

\[ Y \rightarrow Y' \quad Y \rightleftharpoons Y'. \]

Combination of these mechanisms yields to more elaborate networks involving intermediates, as in Examples 2.6 and 3.5 below.

**Definition 2.3.** Let \( \mathcal{Y} \) be a set of intermediates and \( Y \in \mathcal{Y} \).

- A non-intermediate complex \( c \) is called an input for \( Y \) if there is a directed path from \( c \) to \( Y \) in the digraph associated with the network, such that all vertices other than \( c \) belong to \( \mathcal{Y} \).
- \( Y \) is called an \( \ell \)-input intermediate if there are \( \ell \) inputs for \( Y \).

**Example 2.4.** Consider the following network with \( \mathcal{Y} = \{Y_1, Y_2, Y_3\} \):

\[
\begin{align*}
X_1 + X_2 &\rightarrow Y_1 \rightleftharpoons Y_2 \rightarrow Y_3 \rightarrow X_3 + X_4.
\end{align*}
\]

There are two non-intermediate complexes, \( X_1 + X_2 \) and \( X_3 + X_4 \). The species \( Y_1, Y_2, Y_3 \) are all 1-input intermediates. Note that \( Y_2 \) is however the product of two reactions.

Consider now the following network with \( \mathcal{Y} = \{Y\} \):

\[
\begin{align*}
X_1 + X_2 &\rightleftharpoons Y \rightleftharpoons X_3 + X_4.
\end{align*}
\]

The species \( Y \) is a 2-input intermediate and \( X_1 + X_2 \) and \( X_3 + X_4 \) are both inputs for \( Y \).

2.1. **Intermediates and steady states.** Let \( \tilde{\mathcal{N}} \) be a reaction network with a set of intermediates \( \mathcal{Y} = \{Y_1, \ldots, Y_m\} \). Consider the steady state polynomials of the intermediates and denote the parameter vector of reaction rate constants by \( \kappa \) (the reason why will be made clear below). By definition, for every intermediate \( Y_i \), the variable \( y_i \) is only part of the monomial \( y_i \) in (1). Thus, the system with \( m \) equations

\[
F_1(y, x) = \cdots = F_m(y, x) = 0
\]

is linear in \( y_1, \ldots, y_m \). It is shown in [11] that this system has a unique solution for fixed positive values of \( \kappa \) and \( x \), which is further positive. The solution is of the form

\[
y_i = \sum_{c \in C} \mu_{i,c} x^c \quad i = 1, \ldots, m,
\]

where \( \mu_{i,c} \) are rational functions in \( \kappa \) with nonnegative coefficients. The explicit dependence of \( \mu_{i,c} \) on \( \kappa \) is omitted from the notation for simplicity. An explicit description of \( \mu_{i,c} \) can be found using the Matrix-Tree theorem on a suitable labeled digraph, see the supplementary material of [11] or the proof of Lemma 4.7 below.

**Example 2.5.** Consider the following reaction network with \( \mathcal{X} = \{X_1, X_2, X_3\} \) and \( \mathcal{Y} = \{Y_1, Y_2, Y_3\} \):

\[
\begin{align*}
X_1 + X_2 &\xrightarrow{\kappa_1} Y_1 \xrightarrow{\kappa_3} Y_2 \xrightarrow{\kappa_6} 2X_1,
Y_2 &\xrightarrow{\kappa_4} 2X_2,
Y_3 &\xrightarrow{\kappa_5} Y_1 \xrightarrow{\kappa_9} 2X_1.
\end{align*}
\]
The linear system in $y_1, y_2, y_3$ that the steady state polynomials of $Y_1, Y_2, Y_3$ define is:

\[
\begin{align*}
\kappa_1 x_1 x_2 - (\kappa_2 + \kappa_3 + \kappa_5) y_1 &= 0, \\
\kappa_3 y_1 - \kappa_4 y_2 &= 0, \\
\kappa_5 y_1 - (\kappa_6 + \kappa_8) y_3 + \kappa_7 x_1^2 &= 0,
\end{align*}
\]

and its solution is

\[
\begin{align*}
y_1 &= \frac{\kappa_1}{\kappa_2 + \kappa_3 + \kappa_5} x_1 x_2, \\
y_3 &= \frac{\kappa_1 \kappa_5}{(\kappa_6 + \kappa_8)(\kappa_2 + \kappa_3 + \kappa_5)} x_1 x_2 + \frac{\kappa_7}{\kappa_6 + \kappa_8} x_1^2.
\end{align*}
\]

This gives

\[
\begin{align*}
\mu_{1, x_1 + x_2} &= \frac{\kappa_1}{\kappa_2 + \kappa_3 + \kappa_5}, & \mu_{1, x_1} &= 0, & \mu_{1, 2x_1} &= 0, \\
\mu_{2, x_1 + x_2} &= \frac{\kappa_4}{\kappa_4(\kappa_2 + \kappa_3 + \kappa_5)}, & \mu_{2, x_1} &= 0, & \mu_{2, 2x_1} &= 0, \\
\mu_{3, x_1 + x_2} &= \frac{\kappa_1 \kappa_5}{(\kappa_6 + \kappa_8)(\kappa_2 + \kappa_3 + \kappa_5)}, & \mu_{3, x_1} &= \frac{\kappa_7}{\kappa_6 + \kappa_8}, & \mu_{3, 2x_1} &= 0.
\end{align*}
\]

**Example 2.6.** The following digraph corresponds to the Mitogen-Activated Protein Kinase cascade (MAPK) given in [3]:

\[
\begin{align*}
X_0 + E & \xrightarrow[\kappa_1]{} Y_1 \xrightarrow[\kappa_3]{} X_1 + E \xrightarrow[\kappa_4]{} Y_2 \xrightarrow[\kappa_5]{} X_2 + E \\
X_2 + F & \xrightarrow[\kappa_7]{} Y_3 \xrightarrow[\kappa_9]{} Y_4 \xrightarrow[\kappa_{10}]{} X_1 + F \xrightarrow[\kappa_{12}]{} Y_5 \xrightarrow[\kappa_{14}]{} Y_6 \xrightarrow[\kappa_{15}]{} X_0 + F.
\end{align*}
\]

Species $Y_1, \ldots, Y_6$ are intermediates. The non-zero coefficients $\mu_{i,c}$ are:

\[
\begin{align*}
\mu_{1, x_0 + E} &= \frac{\kappa_1}{\kappa_2 + \kappa_3}, & \mu_{2, x_1 + E} &= \frac{\kappa_4}{\kappa_5 + \kappa_6}, & \mu_{3, x_2 + F} &= \frac{\kappa_7}{\kappa_8 + \kappa_9}, \\
\mu_{4, x_2 + F} &= \frac{\kappa_8}{\kappa_8 + \kappa_{10}}, & \mu_{4, x_1 + F} &= \frac{\kappa_{11}}{\kappa_{10}}, & \mu_{5, x_1 + F} &= \frac{\kappa_{12}}{\kappa_{13} + \kappa_{14}}, \\
\mu_{6, x_1 + F} &= \frac{\kappa_{12} \kappa_{14}}{(\kappa_{13} + \kappa_{14}) \kappa_{15}}, & \mu_{6, x_0 + F} &= \frac{\kappa_{16}}{\kappa_{15}}.
\end{align*}
\]

### 2.2. Extended and core networks.

**Definition 2.7.** Let $\mathcal{N} = (\mathcal{S}, \mathcal{C}, \mathcal{R})$ and $\tilde{\mathcal{N}} = (\tilde{\mathcal{S}}, \tilde{\mathcal{C}}, \tilde{\mathcal{R}})$ be two reaction networks. We say that $\tilde{\mathcal{N}}$ is an extension of $\mathcal{N}$ via the addition of intermediates $Y_1, \ldots, Y_m$ if

(i) $\mathcal{Y} = \{Y_1, \ldots, Y_m\}$ is a set of intermediates of $\tilde{\mathcal{N}}$.
(ii) $\mathcal{S} \cup \mathcal{Y} \subseteq \tilde{\mathcal{S}}$ and $\mathcal{C} \cup \mathcal{Y} \subseteq \tilde{\mathcal{C}}$.
(iii) $c \to c' \in \mathcal{R}$ if and only if there is a directed path from $c$ to $c'$ in the digraph associated with $\tilde{\mathcal{N}}$, such that all vertices other than $c$ and $c'$ belong to $\mathcal{Y}$ (there might be none).

In this case $\tilde{\mathcal{N}}$ is called the core network of $\tilde{\mathcal{N}}$.

**Example 2.8.** The core network associated with the network in Example 2.5 is:

\[
\begin{align*}
X_1 + X_2 & \xrightarrow[k_1]{k_1} 2X_2 \\
& \xrightarrow[k_{2,3}]{k_2} 2X_1
\end{align*}
\]
Example 2.9. The core network of the network in Example 2.6 has digraph
\[X_0 + E \xrightarrow{k_1} X_1 + E \xrightarrow{k_2} X_2 + E \xrightarrow{k_3} X_1 + F \xrightarrow{k_4} X_0 + F.\]

Notations \(\kappa, \tilde{I}, \tilde{F}\) are used to address reaction rate constants, steady state ideal and steady state polynomials of the extended network respectively. This notation is fixed from now on whenever we study extensions via the addition of intermediates.

Given \(\tilde{N}\) an extension of \(N\) via the addition of intermediates \(Y_1, \ldots, Y_m\), we define a map
\[\phi: \mathbb{R}[k] \rightarrow \mathbb{R}(\kappa), \quad k_{c \rightarrow c'} \mapsto \phi_{c \rightarrow c'}(\kappa),\]
such that for every reaction \(c \rightarrow c' \in \mathcal{R}\), \(\phi_{c \rightarrow c'}(\kappa)\) is the rational function
\[\phi_{c \rightarrow c'}(\kappa) = \kappa_{c \rightarrow c'} + \sum_{i=1}^{m} \kappa_{Y_i \rightarrow c'} \mu_{i,c},\]
where it is understood that \(\kappa_{c \rightarrow c'} = 0, \kappa_{Y_i \rightarrow c'} = 0\) if respectively \(c \rightarrow c', Y_i \rightarrow c'\) do not belong to \(\mathcal{R}\). Note that \(\phi_{c \rightarrow c'}(\kappa) \neq 0\) for all \(c \rightarrow c'\) by Definition 2.7(iii) and that \(\phi_{c \rightarrow c'}(\kappa)\) is a rational function with positive coefficients.

The map \(\phi\) extends to a map
\[\Phi: \mathbb{R}[k][x] \rightarrow \mathbb{R}(\kappa)[y,x].\]
For example, if \(F_i\) is a steady state polynomial of \(N\), \(\Phi(F_i)\) is the polynomial obtained by replacing \(k_{c \rightarrow c'}\) by the rational function \(\phi_{c \rightarrow c'}(\kappa)\). If the rational functions \(\phi_{c \rightarrow c'}(\kappa)\) are algebraically independent over \(\mathbb{R}\), then \(\phi\) extends to a map of polynomial rings
\[\Phi: \mathbb{R}(k)[x] \rightarrow \mathbb{R}(\kappa)[y,x].\]
We explore in Section 4 ways to check whether the algebraic independence condition holds, and provide types of intermediates for which it holds and no extra check is required.

We introduce the following polynomials
\[H_i(y,x) = y_i - \sum_{c \in \mathcal{C}} \mu_{i,c} x^c \in \mathbb{R}(\kappa)[y,x], \quad i = 1, \ldots, m.\]

Theorem 2.10. ([11, Theorems 3.1 and 3.2]) Let \(\tilde{N}\) be an extension of \(N\) via the addition of intermediates \(Y_1, \ldots, Y_m\).

(i) The coefficient \(\mu_{i,c}\) is nonzero if and only if the non-intermediate complex \(c\) is an input for \(Y_i\) in \(\tilde{N}\).
(ii) The set of steady state polynomials of non-intermediate species and the polynomials \(H_1, \ldots, H_m\) in (4) form a basis of \(\tilde{I}\).
(iii) \(\tilde{F}_i\left(\sum_{c \in \mathcal{C}} \mu_{1,c} x^c, \ldots, \sum_{c \in \mathcal{C}} \mu_{m,c} x^c, x_1, \ldots, x_n\right) = \Phi(F_i(x))\) for \(i = 1, \ldots, n\).

Statements (ii) and (iii) of the previous theorem constitute the proof of the following corollary.
Corollary 2.11. Let $B$ be the set of steady state polynomials of $N$. Then
\[
\tilde{I} = \langle \Phi(B) \cup \{H_1(y,x), \ldots, H_m(y,x)\} \rangle.
\]

We conclude this section with basic properties of $\Phi$.

Lemma 2.12. With the notation above, assume $\phi_{c \to c'}(\kappa)$ for all $c \to c' \in R$ are algebraically independent over $\mathbb{R}$. Let $B = \{f_1, \ldots, f_\ell\}$ and $B' = \{f'_1, \ldots, f'_\ell'\}$ be two sets in $\mathbb{R}(\kappa)[x]$.

(i) If $f \in \langle B \rangle$, then $\Phi(f) \in \langle \Phi(B) \rangle$.

(ii) If $\langle B \rangle = \langle B' \rangle$, then $\langle \Phi(B) \rangle = \langle \Phi(B') \rangle$. Thus $\Phi(\langle B \rangle)$ is well defined.

Proof. (i) Write $f = \sum_{j=1}^\ell \alpha_j f_j$ with $\alpha_j \in \mathbb{R}(\kappa)[x]$. Then
\[
\Phi(f) = \sum_{j=1}^\ell \Phi(\alpha_j) \Phi(f_j) \in \langle \Phi(B) \rangle.
\]

(ii) It is enough to show inclusion $\subseteq$, since the other inclusion is analogous. If $g \in \langle \Phi(B) \rangle$, we have
\[
g = \sum_{i=1}^\ell \lambda_i \Phi(f_i), \quad \lambda_i \in \mathbb{R}(\kappa)[y,x].
\]
Since $f_i \in \langle B' \rangle$, we have by (i) that $\Phi(f_i) \in \langle \Phi(B') \rangle$. In particular, $g$ is an algebraic combination of the polynomials $\Phi(f'_1), \ldots, \Phi(f'_{\ell'})$ with coefficients in $\mathbb{R}(\kappa)[y,x]$. Thus $g \in \langle \Phi(B') \rangle$. \qed

3. Gröbner bases and intermediates

Typically, the values of the reaction rate constants are unknown and reaction networks of interest involve a considerable number of variables. As a consequence, finding a Gröbner basis of the steady state ideal over the field $\mathbb{R}(\kappa)$ can be a demanding task, and sometimes even impossible with standard computers. However, the presence of intermediates, a common feature of reaction networks, can reduce the computation time substantially, by exploiting the structure of the steady state polynomials associated with intermediates given in Theorem 2.10. The main result of this section is Theorem 3.4. Example 3.5 illustrates how the computation time can be reduced by applying our results.

We start with some concepts from computational algebraic geometry.

3.1. Monomial orders and Gröbner bases. We follow the notation on Gröbner bases from [5]. We give here a brief overview of the results required in this text.

Given a monomial order on $R = K[x_1, \ldots, x_n]$, with $K$ a field, let $\text{LM}(f)$ and $\text{LT}(f)$ denote respectively the leading monomial and leading term of $f$. That is, $\text{LT}(f) = \alpha \text{LM}(f)$ if $\alpha$ is the coefficient of the greatest monomial of $f$. Then, for a subset $A \subseteq R$, one defines $\text{LT}(A) = \{ \text{LT}(f) \mid f \in A \}$ and $\text{LM}(A) = \{ \text{LM}(f) \mid f \in A \}$. Clearly,
\[
\langle \text{LT}(A) \rangle = \langle \text{LM}(A) \rangle.
\]
For an ideal $I$, the initial ideal is the ideal generated by the leading terms of the elements of $I$, $\langle \text{LT}(I) \rangle$. A subset $G \subseteq I$ is called a Gröbner basis for $I$ if
\[
\langle \text{LT}(I) \rangle = \langle \text{LT}(G) \rangle, \quad \text{(equiv. } \langle \text{LM}(I) \rangle = \langle \text{LM}(G) \rangle).\]

A Gröbner basis is a basis of $I$ as well. Further, $G$ is a reduced Gröbner basis if additionally for every element $g \in G$ none of its terms can be divided by the leading monomial of an element in $G - \{g\}$, and the coefficient of $\text{LM}(g)$ is 1.

Whether a basis of an ideal is a Gröbner basis depends on the chosen monomial order. Given an ideal and a monomial order, the Gröbner basis is not unique but there is a unique reduced Gröbner basis (see [5]).

We will use the following lemma, which follows from Lemma 2.3.1 and Theorem 2.3.2 of [15]. Alternatively it follows from Proposition 2.9.4 and Buchberger’s criterion (Theorem 2.6.5) in [6].

**Lemma 3.1.** Let $B$ be a basis of $I$. If the leading monomials of every pair $f, g \in B$ are relatively prime, then $B$ is a Gröbner basis.

All monomial orders are defined via a matrix in the following way (though not all matrices $M$ define a monomial order in this way, [5, 26]). For $M \in \mathbb{R}^{n \times n}$ with full rank, the associated order fulfills $x^{c_1} > x^{c_2}$ if the first non-zero entry of the vector $M(c_1 - c_2)$ is positive.

A typical order is the lexicographic monomial order, $\text{lex}$. After choosing a variable order $x^{a_1} > \cdots > x^{a_n}$, $\text{lex}(x^{a_1}, \ldots, x^{a_n})$ is the order defined by the matrix with 1 in positions $(i, a_i)$ for all $i = 1, \ldots, n$ and zero otherwise.

Another monomial order of interest is the graded reverse-lexicographic order, abbreviated $\text{grevlex}$. With this order, $x^{c_1} > x^{c_2}$ if the total degree of the first monomial is larger than the second. If they are equal, then the monomial with the smallest variable with least exponent is the greatest one. Grevlex with order of variables $x^{1} > \cdots > x^{n}$ is defined by the matrix
\[
\begin{pmatrix}
1 & 1 & \cdots & 1 & 1 \\
0 & 0 & \cdots & 0 & -1 \\
0 & 0 & \cdots & -1 & 0 \\
\vdots & \vdots & \ddots & \vdots & \vdots \\
0 & -1 & \cdots & 0 & 0
\end{pmatrix}.
\]

The choice of order plays an important role in the computation time for Gröbner bases, performing lex typically worse than grevlex. However, lex, as any other elimination type order, has a crucial property on elimination of variables. Given a partitioning of the set of variables, $\{x_1, \ldots, x_n\} = \{x_{j_1}, \ldots, x_{j_{n-s}}\} \cup \{x_{i_1}, \ldots, x_{i_s}\}$, a monomial order is of elimination type if $x^{e_{j \ell}}$, for $\ell = 1, \ldots, n-s$, is larger than any monomial in $K[x_{i_1}, \ldots, x_{i_s}]$ [6] §3.1, Exercise 5]. Clearly, $\text{lex}(x_{j_1}, \ldots, x_{j_{n-s}}, x_{i_1}, \ldots, x_{i_s})$ is of elimination type. If $G$ is a Gröbner basis of $I$ with respect to an elimination type order as above, then $G \cap K[x_{i_1}, \ldots, x_{i_s}]$ is a Gröbner basis of $I \cap K[x_{i_1}, \ldots, x_{i_s}]$ with respect to the induced monomial on $K[x_{i_1}, \ldots, x_{i_s}]$, which for lex is $\text{lex}(x_{i_1}, \ldots, x_{i_s})$.

### 3.2. Gröbner bases and intermediates

In this subsection we fix a reaction network $\mathcal{N}$ and an extension $\tilde{\mathcal{N}}$ via the addition of intermediates $Y_1, \ldots, Y_m$. We
show that any Gröbner basis of the steady state ideal of $\tilde{N}$ can be extended to one of $\tilde{N}$ by simply adding the polynomials $H_1, \ldots, H_m$ given in Equation (3). By default, we order the variables $y_1 > \cdots > y_m > x_1 > \cdots > x_n$. We start with some general lemmas.

**Lemma 3.2.** Let $I = \langle f_0, f_1, \ldots, f_s \rangle \subseteq \mathbb{K}[y_1, x_1, \ldots, x_n]$ be an ideal such that $f_i \in \mathbb{K}[x_1, \ldots, x_n]$ for $i = 1, \ldots, s$ and $f_0 = y + f_0'$, with $f_0' \in \mathbb{K}[x_1, \ldots, x_n]$. Consider a monomial order defined by a matrix $M$ whose first row is $(1 \ 0 \ \cdots \ 0)$. Then

$$\langle \text{LT}(I) \rangle = \langle y \rangle + \langle \text{LT}((f_1, \ldots, f_s)) \rangle.$$

Further given $G \subseteq \mathbb{K}[x_1, \ldots, x_n]$, $G$ is a Gröbner basis of $\langle f_1, \ldots, f_s \rangle$ if and only if $\{f_0\} \cup G$ is a Gröbner basis of $I$.

**Proof.** By the choice of monomial order, the monomial $y$ is larger than any monomial not involving $y$. Consider a reduced Gröbner basis $G'$ of $\langle f_1, \ldots, f_s \rangle$. Then the leading terms of the elements in $G'$ are relatively prime with each other and with the leading term of $f_0$. Since $\{f_0\} \cup G'$ is a basis of $I$, then by Lemma 3.1 $\{f_0\} \cup G'$ is a Gröbner basis of $I$. Now, the initial ideal of $I$ is generated by the leading terms of $\{f_0\} \cup G'$. So:

$$\langle \text{LT}(I) \rangle = \langle \text{LT}(\{f_0\} \cup G') \rangle = \langle \{y\} \cup \text{LT}(G') \rangle = \langle y \rangle + \langle \text{LT}(G') \rangle = \langle y \rangle + \langle \text{LT}((f_1, \ldots, f_s)) \rangle.$$

This proves the first part of the lemma.

For the second part, note that

$$\langle y \rangle + \langle \text{LT}(G) \rangle = \langle \{\text{LT}(f_0)\} \cup \text{LT}(G) \rangle = \langle \text{LT}(\{f_0\} \cup G) \rangle.$$

Using this equality and the first part of the lemma, we have $\{f_0\} \cup G$ is a Gröbner basis of $I$ if and only if $\langle y \rangle + \langle \text{LT}(G) \rangle = \langle y \rangle + \langle \text{LT}(\{f_1, \ldots, f_s\}) \rangle$. Since $y$ is not part of any polynomial in $G$, this equality holds if and only if $\langle \text{LT}(G) \rangle = \langle \text{LT}(\{f_1, \ldots, f_s\}) \rangle$, i.e. $G$ is a Gröbner basis of $\langle f_1, \ldots, f_s \rangle$. □

Recall that we write $I \subseteq \mathbb{R}(k)[x_1, \ldots, x_n]$ and $\tilde{I} \subseteq \mathbb{K}(\kappa)[y_1, \ldots, y_m, x_1, \ldots, x_n]$ for the steady state ideals of $N$ and $\tilde{N}$ respectively. For the rest of the section, we assume that the rational functions $\phi_{\kappa \to \kappa'}(\kappa)$ are algebraically independent over $\mathbb{R}$, such that $\Phi(A)$ is defined for all subsets $A$ of $\mathbb{R}(k)[x]$.

For an arbitrary basis $B$ of $I$, define

$$B = \Phi(B) \cup \{H_1(y, x), \ldots, H_m(y, x)\} \subseteq \mathbb{K}(\kappa)[y, x].$$

**Lemma 3.3.** If $B$ is a basis of $I \subseteq \mathbb{R}(k)[x]$, then $B$ is a basis of $\tilde{I} \subseteq \mathbb{K}(\kappa)[y, x]$.

**Proof.** Let $B'$ be the set of steady state polynomials of $N$. By Corollary 2.11

$$\tilde{I} = \langle \Phi(B') \cup \{H_1(y, x), \ldots, H_m(y, x)\} \rangle.$$

Let now $B$ be an arbitrary basis of $I$. Then $\langle B \rangle = \langle B' \rangle$ and thus by Lemma 2.12 (ii), $\langle \Phi(B) \rangle = \langle \Phi(B') \rangle$. Therefore

$$\langle \Phi(B) \cup \{H_1(y, x), \ldots, H_m(y, x)\} \rangle = \langle \Phi(B') \cup \{H_1(y, x), \ldots, H_m(y, x)\} \rangle = \tilde{I}.$$

This completes the proof. □
Let \( \text{rem}(p, B) \) be the remainder of the division of the polynomial \( p \) by a set of polynomials \( B \).

**Theorem 3.4.** Fix a monomial order on \( \mathbb{R}(k)[x] \) associated with an \( n \times n \) matrix \( Q \), and let \( G \) be a Gröbner basis of \( I \) with this order. Then, \( \widetilde{G} \) is a Gröbner basis of \( \widetilde{I} \) with the monomial order on \( \mathbb{R}(\kappa)[y, x] \) associated with the matrix

\[
\widetilde{Q} = \begin{pmatrix} \text{Id}_m & 0 \\ 0 & Q \end{pmatrix},
\]

where \( \text{Id}_m \) is the identity matrix of size \( m \).

If \( G \) is reduced, then \( \Phi(G) \cup \{ y_i - \text{rem} \left( \sum_{c \in C} \mu_{i,c} x^c, \Phi(G) \right) \} \) is the reduced Gröbner basis of \( \widetilde{I} \).

**Proof.** First note that by the monomial order given by \( \widetilde{Q} \), we have \( y_1 > \cdots > y_m > x_i \) for all \( i = 1, \ldots, n \). Also, the polynomial \( H_i \) has degree one in \( y_i \) and none of the elements of \( \Phi(G) \cup \{ H_j \mid j \neq i \} \) involves \( y_i \).

Let us assume we have shown that \( \Phi(G) \) is a Gröbner basis of \( \langle \Phi(G) \rangle \) with the given order, that is

\[
\langle \text{LT}(\langle \Phi(G) \rangle) \rangle = \langle \text{LT}(\Phi(G)) \rangle.
\]

Then by Lemmas 3.2 and 3.3, \( \Phi(G) \cup \{ H_1(y, x), \ldots, H_m(y, x) \} \) is a Gröbner basis of \( \widetilde{I} \). Therefore the first part of the statement holds provided (8) holds.

Let us show (8). We start by noting that for a subset \( J \) in \( \mathbb{R}(k)[x] \), the set \( \text{LM}(J) \) consists only of monomials in \( x_1, \ldots, x_n \), and thus is naturally included in \( \mathbb{R}(\kappa)[y, x] \) as well. Further

\[
\text{LM}(J) = \text{LM}(\Phi(J)).
\]

Let \( G' \) be a reduced Gröbner basis of \( I \). Since \( G' \) is reduced, pairs of monomials in \( \text{LM}(G') = \text{LM}(\Phi(G')) \) are relatively prime. Since \( \Phi(G') \) is a basis of \( \langle \Phi(G') \rangle \), then by Lemma 3.1 and Equation (5), it is actually a Gröbner basis and (8) holds for \( G' \). Now, consider an arbitrary Gröbner basis \( G \) of \( I \). In \( \mathbb{R}(k)[x] \) it holds

\[
\langle \text{LM}(G) \rangle = \langle \text{LM}(G') \rangle.
\]

This means that every monomial in \( \langle \text{LM}(G') \rangle \) is divisible by a monomial in \( \langle \text{LM}(G) \rangle \) and vice versa \( \mathbb{R} \) §2.4, Lemma 2. Since this fact holds also in \( \mathbb{R}(\kappa)[y, x] \), (10) holds also in \( \mathbb{R}(\kappa)[y, x] \). Combined with (9) this gives

\[
\langle \text{LM}(\Phi(G)) \rangle = \langle \text{LM}(\Phi(G')) \rangle.
\]

By Lemma 2.12(ii), \( \langle G \rangle = \langle G' \rangle \) in \( \mathbb{R}(k)[x] \) implies \( \langle \Phi(G) \rangle = \langle \Phi(G') \rangle \). Thus in \( \mathbb{R}(\kappa)[y, x] \) we have

\[
\langle \text{LM}(\Phi(G)) \rangle = \langle \text{LM}(\Phi(G')) \rangle = \langle \text{LM}(\langle \Phi(G') \rangle) \rangle = \langle \text{LM}(\langle \Phi(G) \rangle) \rangle.
\]

This shows that (8) holds.

The second part of the lemma is clear from the definition of a reduced Gröbner basis and using that \( \Phi(G) \cup \{ y_i - \text{rem} \left( \sum_{c \in C} \mu_{i,c} x^c, \Phi(G) \right) \} \) is also a Gröbner basis.

\( \square \)
From the computational point of view, Theorem 3.4 is very useful. Instead of computing a Gr"obner basis of $\bar{I}$ directly, one can first compute a Gr"obner basis $G$ for the core network, with a smaller number of variables and polynomials, then add the polynomials $y_i - \sum_{c \in C} \mu_{i,c} x^c$, and, finally, simplify them using polynomial division by $\Phi(G)$. The second step involves only linear algebra. A possible issue here is to verify that the rational functions $\phi_{c \rightarrow c'}$ are algebraically independent. We provide in Section 4 a list of network structures involving intermediates for which the condition is fulfilled.

**Example 3.5.** An interesting example to show the advantage of using Theorem 3.4 is Example 4.4 of [4]. We consider the reaction network $\tilde{N}$ with associated digraph given in Figure 1.

This reaction network has 29 species and 46 reactions. Therefore the steady state ideal is generated by 29 polynomials in 29 variables and 46 parameters. Using Singular [7] and monomial order grevlex with $x_1 > \cdots > x_{29}$ (the same monomial order that is used in [4]), it took between 110 and 115 seconds\(^1\) to compute the reduced Gröbner basis. This basis consists of 169 polynomials.

\(^1\)Information about the processor: Intel(R) Core(TM) i5-3570 CPU 3.3GHz 3.4GHz with 8GB RAM. We report the interval of obtained times after several runs of Singular, computed in milliseconds.
Now we consider the monomial order introduced in Theorem 3.4 for the removal of the 15 intermediates:

\[ X_3, X_6, X_8, X_{10}, X_{12}, X_{13}, X_{15}, X_{17}, X_{20}, X_{21}, X_{23}, X_{24}, X_{25}, X_{27}, X_{29}. \]

The original network \( \widetilde{N} \) is an extension of the following core network \( N \) with 14 species and 16 reactions:

\[
\begin{align*}
X_1 + X_2 & \xrightarrow{k_1} X_2 + X_4 \\
X_4 + X_{18} & \xrightarrow{k_2} X_{18} + X_{26} \\
X_{18} + X_{22} & \xrightarrow{k_{13}} X_{16} + X_{22} \\
X_5 + X_{16} & \xrightarrow{k_{14}} X_5 + X_{14} \\
X_{18} + X_{28} & \xrightarrow{k_{15}} X_{18} + X_{19} \\
X_{19} & \xrightarrow{k_{16}} X_{28}.
\end{align*}
\]

The functions \( \phi_{c\to c'} \) are algebraically independent over \( \mathbb{R} \) by Corollary 4.6 in Section 4. We consider grevlex with \( x_1 > \cdots > x_{28} \) for the monomials corresponding to \( N \). The monomial order in Theorem 3.4 is then associated with the following matrix

\[
\tilde{Q} = \begin{pmatrix}
1 & 0 & \cdots & 0 \\
0 & 1 & \ddots & \vdots \\
\vdots & \ddots & \ddots & \cdots \\
0 & \cdots & 0 & 1
\end{pmatrix},
\]

and order of variables

\[
x_3 > x_6 > x_8 > x_{10} > x_{12} > x_{13} > x_{15} > x_{17} > x_{20} > x_{21} > x_{23} > x_{24} \\
> x_{25} > x_{27} > x_{29} > x_{21} > x_2 > x_4 > x_5 > x_7 > x_9 > x_{11} > x_{14} > x_{16} \\
> x_{18} > x_{19} > x_{22} > x_{26} > x_{28}.
\]

The reduced Gröbner basis of \( \tilde{I} \) with this monomial order has 33 polynomials and it takes about 96 seconds to compute it directly with Singular. Alternatively the strategy outlined in Theorem 3.4 can be applied. The steady state ideal \( I \) of \( N \) is generated by 11 polynomials in 14 variables and 16 parameters. Using Singular, the reduced Gröbner basis of \( I \) has 18 polynomials and its computation takes less than a millisecond. The computation time for the polynomials \( H_i(y,x) \) is negligible, since they are found by solving 15 independent linear equations. Therefore the reduced Gröbner basis of the ideal of the original system has 18+15=33 polynomials and can be computed in less than a millisecond.

We conclude that in general, regarding computational time, the monomial order introduced in Theorem 3.4 is a good choice for networks with intermediates, and further, by applying the strategy of Theorem 3.4 we reduce the computation time considerably, compared with direct computation of the reduced Gröbner basis.
Remark 3.6. Theorem 3.4 holds regardless the choice of method to compute a Gröbner basis. Since the computation of the polynomials $H_i$ is simple linear algebra, even for the fastest available methods for the computation of Gröbner bases, decomposing the computation as in Theorem 3.4 should be faster than direct computation of the basis of the steady state ideal of $\tilde{\mathcal{N}}$.

Remark 3.7. For polynomials with integer coefficients, it is usually faster to compute a Gröbner basis using the so-called modular approach, see e.g. [23, 29]. These methods first choose a so-called lucky prime and compute a Gröbner basis of the ideal in $\mathbb{Z}_p[x]$. Then the coefficients of this Gröbner basis are lifted to a Gröbner basis in $\mathbb{Q}[x]$. For the sake of comparison, we also computed how long it takes to find a Gröbner basis using $p$-modular approaches on the extended network in Example 3.5 with grevlex and $x_1 > \cdots > x_{29}$. Using the largest prime number in Singular, $p = 32003$, it takes 127 seconds to compute the Gröbner basis over $\mathbb{Z}_{32003}$. Since coefficients in the starting basis are 1 or $-1$, one may think that $p = 2$ is a lucky prime. It took 97 seconds to compute the Gröbner basis over $\mathbb{Z}_2$. These times are larger than the times reported in Example 3.5 (and these Gröbner bases still need to be lifted to $\mathbb{Q}(\kappa)[y,x]$).

An important consequence of Theorem 3.4 concerns parameter-free model discrimination. In this setting one seeks elements of the steady state ideal $\tilde{I}$ involving only the concentrations of species that are experimentally measurable. These elements are called invariants. Each invariant implies that there is a set of monomials that lie on a hyperplane, and the hypothesis of coplanary is then tested using experimental data [13, 14, 18–20]. This approach is attractive because it does not require knowing the values of the reaction rate constants.

Experimentally measurable species do not typically involve intermediates. In this case, Theorem 3.4 tells us that invariants on the non-intermediate species can be computed directly from the core network, using elimination ideals.

Corollary 3.8. Let $\mathcal{N}$ be a reaction network and $\tilde{\mathcal{N}}$ an extension of it via the addition of $m$ intermediates $Y_1, \ldots, Y_m$. Let $X_{i_1}, \ldots, X_{i_p}$ be non-intermediates. Then

$$\tilde{I} \cap \mathbb{R}(\kappa)[x_{i_1}, \ldots, x_{i_p}] = \Phi(I \cap \mathbb{R}(k)[x_{i_1}, \ldots, x_{i_p}]).$$

Proof. For simplicity, assume $\{i_1, \ldots, i_p\} = \{n-p+1, \ldots, n\}$ and let $\bar{x} = (x_{n-p+1}, \ldots, x_n)$. Consider the monomial order $\text{lex}(y_1, \ldots, y_m, x_1, \ldots, x_n)$ on $\mathbb{R}(\kappa)[y,x]$, and $\text{lex}(x_1, \ldots, x_n)$ on $\mathbb{R}(k)[x]$. Let $G$ be a Gröbner basis of $I$. By Theorem 3.4 $\tilde{G}$ is a Gröbner basis of $\tilde{I}$. By the properties of lex and Lemma 2.12(ii) we have

$$\tilde{I} \cap \mathbb{R}(\kappa)[\bar{x}] = \langle \tilde{G} \cap \mathbb{R}(\kappa)[\bar{x}] \rangle = \langle \Phi(G \cap \mathbb{R}(k)[\bar{x}]) \rangle = \Phi(I \cap \mathbb{R}(k)[\bar{x}]).$$

This concludes the proof.

Note that the monomial order on $\mathbb{R}(\kappa)[y,x]$ given in Theorem 3.4 is of elimination type with respect to the partition $\{y_1, \ldots, y_m\} \cup \{x_1, \ldots, x_n\}$.

Example 3.9. Consider the network in Example 2.6 and its core network in Example 2.9. In order to find invariants of the extended network involving the concentration of the non-intermediate species $E, X_0, X_1, X_2$, we consider the ideal $I \cap
\( \mathbb{R}(k)[e, x_0, x_1, x_2] \), which is generated by the polynomial

\[ e(k_1 k_3 x_0 x_2 - k_2 k_4 x_1^2) \]

We have

\[ \phi(k_1, k_2, k_3, k_4) = \begin{pmatrix} x_1 x_3, & k_2 k_4, & k_5 + k_6, & k_7 + k_8, & k_9 + k_10, & k_11 + k_12 \end{pmatrix} \]

The functions \( \phi_{c \to c'} \) are algebraically independent over \( \mathbb{R} \) by Corollary 4.6. By Corollary 3.8 the ideal \( \tilde{I} \cap \mathbb{R}(\kappa)[e, x_0, x_1, x_2] \) is generated by the polynomial

\[ e \left( \frac{k_1 k_3}{k_2 + k_3} \frac{k_5 k_6}{k_5 + k_6} \frac{k_7 k_8}{k_7 + k_8} \frac{k_9 k_{10}}{k_9 + k_{10}} \frac{k_11 k_{12}}{k_11 + k_{12}} x_0 x_2 - \frac{k_4 k_6}{k_4 + k_6} \frac{k_5 k_{13}}{k_5 + k_{13}} \frac{k_6 k_{14}}{k_6 + k_{14}} x_1^2 \right) \]

3.3. Detecting binomial steady state ideals. A binomial is a polynomial having at most two terms. An ideal is said to be binomial if it admits a set of generators consisting of binomials only. By [8, Corollary 1.2], an ideal is binomial if and only if any reduced Gröbner basis (with respect to any monomial order) consists of binomials.

It is of biological relevance in the study of reaction networks to determine whether there exists a choice of reaction rate constants \( k \) for which there are multiple positive steady states in some coset \( x_0 + S \) defined by the vector subspace \( S \) that contains the image of \( F_k \) (see Section 1). This property is termed multistationarity. If the steady state ideal is binomial, then there exist efficient ways to determine whether the network admits multistationarity. This leads to the problem of determining whether an ideal is binomial, and in case it is, of finding a binomial basis of it. As noted, both questions can be addressed by finding a Gröbner basis of the steady state ideal of the network. Thus, for networks with intermediates, our results can be applied also to detect binomial steady state ideals.

Recall that we are assuming that the rational functions \( \phi_{c \to c'}(\kappa) \) are algebraically independent over \( \mathbb{R} \).

Theorem 3.10. Let \( \mathcal{N} \) be a reaction network and \( \tilde{\mathcal{N}} \) an extension of it via the addition of \( m \) intermediates \( Y_1, \ldots, Y_m \).

The steady state ideal \( \tilde{I} \) is binomial if and only if

\( \bullet \) \( \tilde{I} \) is binomial, and,

\( \bullet \) for any reduced Gröbner basis \( G \) of \( I \) and for every \( i = 1, \ldots, m \), the remainder of the division of \( \sum_{c \in C} \mu_i c x^c \) by \( \Phi(G) \) has at most one term.

Proof. Fix any monomial order on \( \mathbb{R}(k)[x_1, \ldots, x_n] \) associated with an \( n \times n \) matrix \( Q \) and consider the monomial order with matrix \( \tilde{Q} \) from Theorem 3.4. Let \( G \) be the reduced Gröbner basis of \( I \) and

\[ \tilde{G}' = \Phi(G) \cup \left\{ y_i - \text{rem} \left( \sum_{c \in C} \mu_{i,c} x^c, \Phi(G) \right) \right\} \]

the reduced Gröbner basis of \( \tilde{I} \) (cf. Theorem 3.4). Using that an ideal is binomial if and only if any reduced Gröbner basis consists of binomials, the theorem is a consequence of the following two facts:

\( \bullet \) By definition, \( \tilde{G}' \) consists of binomials if and only if \( \Phi(G) \) is a set of binomials and the remainder of the division of \( \sum_{c \in C} \mu_{i,c} x^c \) by \( \Phi(G) \) has at most one term.
By the algebraic independence of $\phi_{c \rightarrow c'}$, $\Phi(G)$ consists of binomials if and only if $G$ does.

Since the polynomial $\sum_{c \in C} \mu_i,c x^c$ has exactly one term for 1-input intermediates, we readily obtain the following corollary.

**Corollary 3.11.** Let $N$ be a reaction network and $\tilde{N}$ an extension of it via the addition of $m$ 1-input intermediates $Y_1, \ldots, Y_m$. Then $\tilde{I}$ is binomial if and only if $I$ is binomial.

Since 1-input intermediates are the most abundant form of intermediates found in realistic networks, this corollary implies that in order to check whether a steady state ideal is binomial, we can often remove intermediates and check whether the steady state ideal of the core network is binomial.

**Example 3.12.** Consider the network in Example 2.5 and its core network given in Example 2.8. The functions $\phi_{c \rightarrow c'}$ are algebraically independent over $\mathbb{R}$ by Example 4.1. Since the steady state ideal of $N$ is $\langle -(k_1 - k_2)x_1x_2 - 2k_3x_1^2, (k_1 - k_2)x_1x_2 + 2k_3x_1^2 \rangle$, the core network has a binomial steady state ideal. The reduced Gröbner basis $G$ for this ideal with monomial order $\text{lex}(x_1, x_2, x_3)$ is

$$G = \left\{ x_1^2 - \frac{(k_1 - k_2)}{2k_3} x_1x_2 \right\}.$$

We apply Theorem 3.10 to conclude that the steady state ideal of the extended network is also binomial. The intermediates $Y_1, Y_2$ are 1-input intermediates and hence the remainder condition of the theorem is automatically fulfilled. For the intermediate $Y_3$, $\text{rem} \left( \mu_4, x_1 + x_2, x_2 + x_3, \Phi(G) \right)$ has a single term with monomial $x_1x_2$. Therefore we conclude that the extended network also has a binomial steady state ideal.

The following example shows that extended networks with multi-input intermediates might not have binomial steady state ideals, even though their core networks have.

**Example 3.13.** Consider the network given in Example 2.6 and its core network given in Example 2.9. The steady state ideal of the core network is binomial with basis $B = \{k_1x_0e - k_4x_1f, k_2x_1e - k_3x_2f\}$. The intermediates $Y_4$ and $Y_6$ are 2-input intermediates. The remainder of the division of $\mu_4, x_2 + x_2f + \mu_4, x_1 + f x_1f$ by $\Phi(G)$ for $G$ the reduced Gröbner basis of $I$ with the monomial order $\text{lex}(x_2, x_1, x_0, f, e)$ is

$$\frac{\kappa_4}{\kappa_{10}} x_1f + \frac{\kappa_7 \kappa_9}{\kappa_8 \kappa_{10} + \kappa_9 \kappa_{10}} x_2f,$$

which has two terms. Therefore by Theorem 3.10 the steady state ideal of the network in Example 2.6 is not binomial.

**Remark 3.14.** In [4], a method for determining whether a homogeneous ideal is binomial is introduced. The method avoids the computation of Gröbner bases and
4. Algebraic independence

In this section we discuss how to check whether the functions $\phi_{c \to c'}$ are algebraically independent over $\mathbb{R}$ and provide classes of intermediates for which this property holds. Consider a set of rational functions $A = \{ \frac{f_1}{y_1}, \ldots, \frac{f_m}{y_m} \} \subseteq \mathbb{R}(x_1, \ldots, x_n)$. By §III.7, Theorem III, in [16], the set $A$ is algebraically independent over $\mathbb{R}$ if and only if the rank of the associated Jacobian matrix $\left( \frac{\partial (f_i/g_j)}{\partial x_k} \right)_{i,j,k}$ over $\mathbb{R}(x)$ is $m$.

Another way to check algebraic independence that requires the computation of a Gröbner basis is as follows. Define $U$ to be the open set given by $\mathbb{R}^n$ minus the zero locus of the product $g_1 \cdots g_m$, and let $\varphi$ be the following function on $U$

$$x = (x_1, \ldots, x_n) \mapsto \left( \frac{f_1(x)}{g_1(x)}, \ldots, \frac{f_m(x)}{g_m(x)} \right).$$

By §3.3, Theorem 2, in [6], the closure of $\text{Im}(\varphi)$ is the variety associated with the ideal

$$J := \langle g_1 T_1 - f_1, \ldots, g_m T_m - f_m, 1 - y g_1 \cdots g_m \rangle \cap \mathbb{R}[T_1, \ldots, T_m].$$

Since the sets of polynomials vanishing on a set and on its closure agree (see [6] after Definition 2 in §4.4), $A$ is algebraically independent over $\mathbb{R}$ if and only if $J = \{0\}$.

**Example 4.1.** The functions $\phi_{c \to c'}$ of Examples 2.5 and 2.8 are

$$\phi_{X_1 + X_2 \to 2X_2} = \kappa_4 \mu_2, x_1 + x_2 + \kappa_8 \mu_3, x_1 + x_2 = \frac{\kappa_1 \kappa_3}{\kappa_2 + \kappa_3 + \kappa_5} + \frac{\kappa_1 \kappa_3 \kappa_8}{(\kappa_6 + \kappa_8)(\kappa_2 + \kappa_3 + \kappa_5)},$$
$$\phi_{X_1 + X_2 \to 2X_1} = \kappa_6 \mu_3, x_1 + x_2 = \frac{\kappa_1 \kappa_5 \kappa_8}{(\kappa_6 + \kappa_8)(\kappa_2 + \kappa_3 + \kappa_5)},$$
$$\phi_{2X_1 \to 2X_2} = \kappa_9 + \kappa_8 \mu_3, 2X_1 = \kappa_9 + \frac{\kappa_8 \kappa_9}{\kappa_6 + \kappa_8}.$$

We find that $J = \{0\}$. Hence the algebraic independence condition holds for the network in Example 2.8. Alternatively, one easily checks that the associated Jacobian matrix has rank 3.

The computations above can be simplified by taking into account what parameters occur in each of the rational functions.

**Definition 4.2.** Let $\tilde{\mathcal{N}}$ be an extension of $\mathcal{N}$ via the addition of the intermediates $\{Y_1, \ldots, Y_m\}$. Consider the digraph associated with $\tilde{\mathcal{N}}$. Let $\tilde{\gamma}_1, \ldots, \tilde{\gamma}_p$ denote the vertex sets of the connected components of the subgraph induced by the subset of vertices $\{Y_1, \ldots, Y_m\}$.

Let $\tilde{R}' \subseteq \tilde{R}$ be the subset of reactions of the core network that are not in $\tilde{R}$. These reactions arise necessarily from paths through intermediates. We say that

is regarded as a fast method. If the steady state ideal of the core network is homogeneous, then Theorem 3.10 or Corollary 3.11 in combination with this method provide a fast procedure to detect binomial steady state ideals.

Interestingly, steady state polynomials of core networks are often homogeneous of degree two, since it is common that non-intermediate species appear in complexes of the form $X_i + X_j$, yielding quadratic terms in the steady state polynomials. This is for example the case for so-called Post-Translational Modification Networks [28].
two reactions $r_1: c_1 \rightarrow c_1', r_2: c_2 \rightarrow c_2' \in \mathcal{R}'$ overlap if there exist paths through intermediates

$$c_1 \rightarrow Y_{i_1} \rightarrow \cdots \rightarrow Y_{i_p} \rightarrow c_1', \quad c_2 \rightarrow Y_{j_1} \rightarrow \cdots \rightarrow Y_{j_q} \rightarrow c_2'$$

with all intermediates belonging to the same set $\mathcal{Y}_i$.

Consider the equivalence relation on $\mathcal{R}'$ generated by the overlap relation: $r \sim r'$ if and only if there exist $r_0 = r, r_1, \ldots, r_p = r'$ such that $r_i, r_{i+1}$ overlap for all $i = 0, \ldots, p-1$. Let $\mathcal{R}'_1, \ldots, \mathcal{R}'_t$ be the equivalence classes of this equivalence relation.

**Example 4.3.** Consider the network in Example 2.8. The set $\mathcal{R}'$ consists of two reactions $X_1 + X_2 \rightarrow 2X_2$ and $X_1 + X_2 \rightarrow 2X_1$. The subgraph of the digraph associated with $\tilde{\mathcal{N}}$ induced by the set of intermediates is connected. Thus the two reactions of $\mathcal{R}'$ are equivalent and there is one equivalence class.

**Lemma 4.4.** The set $\{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R}' \}$ is algebraically independent over $\mathbb{R}$ if and only if the set $\{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R}'_i \}$ is algebraically independent over $\mathbb{R}$ for all $i = 1, \ldots, t$.

**Proof.** Since $\mathcal{R}'_i \subseteq \mathcal{R}$ for all $i = 1, \ldots, t$, the forward implication is clear.

To prove the reverse implication, assume that the sets $T_i = \{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R}'_i \}$ are algebraically independent over $\mathbb{R}$ for all $i = 1, \ldots, t$. By construction, the sets of parameters appearing in the rational functions $\phi_{c \rightarrow c'}(\kappa)$ are disjoint for two reactions in different equivalence classes. Therefore the union of the sets $T_1, \ldots, T_t$ is algebraically independent over $\mathbb{R}$. Furthermore if $c \rightarrow c' \in \mathcal{R} \setminus \mathcal{R}'$, then the parameter $\kappa_{c \rightarrow c'}$ appears only in $\phi_{c \rightarrow c'}(\kappa)$. As a consequence the set

$$\bigcup_{i=1}^{t} T_i \cup \{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R} \setminus \mathcal{R}' \} = \{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R} \}$$

is algebraically independent over $\mathbb{R}$.

**Example 4.5.** Consider the network in Example 4.3. The algebraic independence of the functions $\phi_{c \rightarrow c'}(\kappa)$ for all reactions $c \rightarrow c'$ in $\mathcal{R}$ follows in this case from the algebraic independence of the functions $\phi_{c \rightarrow c'}(\kappa)$ for the reactions $X_1 + X_2 \rightarrow 2X_2$ and $X_1 + X_2 \rightarrow 2X_1$.

**Corollary 4.6.** If $\mathcal{R}' = \emptyset$ or each of the equivalence classes $\mathcal{R}'_1, \ldots, \mathcal{R}'_t$ consist of one reaction, then the rational functions $\phi_{c \rightarrow c'}(\kappa)$ are algebraically independent over $\mathbb{R}$.

For the networks in Example 2.6 and Example 3.5, each of the equivalence classes consist of one reaction. Therefore, by Corollary 4.6, the algebraic independence condition holds.

We next show that the algebraic independence condition holds for specific classes of intermediates without the need of doing any extra computation.

**Lemma 4.7.** For the following extension networks, with intermediates $Y_1, \ldots, Y_m$, the set $\{ \phi_{c \rightarrow c'}(\kappa) \mid c \rightarrow c' \in \mathcal{R} \}$ is algebraically independent over $\mathbb{R}$.

(i) $c \leftrightarrow Y_1 \leftrightarrow Y_2 \leftrightarrow \ldots \leftrightarrow Y_m \leftrightarrow c'$, provided $\{Y_1, \ldots, Y_m\}$ is a set of intermediates and where $\leftrightarrow$ means the reaction can be irreversible or reversible.
with an arbitrary digraph structure among the complexes \( c_0, Y_1, \ldots, Y_m \) such that there exists a directed path from \( c_0 \) to \( Y_m \).

(iii) \[
c_0 \xrightarrow{\kappa_1} Y_1 \xrightarrow{\kappa_3} Y_2 \xrightarrow{\kappa_5} \cdots \xrightarrow{\kappa_{2m-1}} Y_m
\]

where some of the reactions with label \( \kappa_{2i} \) might not exist, and \( t_i \geq 0 \) for each \( 1 \leq i \leq m \).

Proof. We start by recalling how to find \( \mu_{i,c} \) using a labeled digraph (see proof of Theorem 2 of the electronic supplementary material of [11]). For each non-intermediate complex \( c \), consider the labeled digraph \( \hat{G}_c \) with vertex set \( \{ Y_1, \ldots, Y_m, \star \} \) and labeled edges \( Y_i \xrightarrow{\kappa_{Y_i \rightarrow Y_j}} Y_j \) if \( Y_i \rightarrow Y_j \in \tilde{R} \), \( \star \xrightarrow{\kappa_{\star \rightarrow Y_i}} Y_i \) if \( c \rightarrow Y_i \in \tilde{R} \) and \( Y_i \xrightarrow{\beta_i} \) \( \star \) with \( \beta_i = \sum_{Y_i \rightarrow c'} \kappa_{Y_i \rightarrow c'} \) if \( \beta_i \neq 0 \).

For every vertex \( v \) of \( \hat{G}_c \), define \( \theta(v) \) as the set of all spanning trees rooted at \( v \) \( ^2 \)

Given such a tree \( \tau \), let \( \pi(\tau) \) be the product of the labels of the edges of \( \tau \). Then

\[
\mu_{i,c} = \frac{\sum_{\tau \in \theta(Y_i)} \pi(\tau)}{\sum_{\tau \in \theta(\star)} \pi(\tau)}.
\]

(i) If one of the reactions is irreversible, then the core network consists of exactly one reaction, either \( c \rightarrow c' \) or \( c' \rightarrow c \), and the set \( \{ \phi_{c \rightarrow c'}(\kappa) | c \rightarrow c' \in \mathcal{R} \} \) is algebraically independent over \( \mathbb{R} \).

If all reactions are reversible, we write

\[
c \xrightarrow{\kappa_1} Y_1 \xrightarrow{\kappa_3} Y_2 \xrightarrow{\kappa_5} \cdots \xrightarrow{\kappa_{2m-1}} Y_m \xrightarrow{\kappa_{2m+1}} c',
\]

and we have \( \phi_{c \rightarrow c'}(\kappa) = \kappa_2 \mu_{1,c'} \), \( \phi_{c \rightarrow c'}(\kappa) = \kappa_{2m+1} \mu_{m,c} \). By the expressions for \( \mu_{1,c'} \) and \( \mu_{m,c} \) in (11), both rational functions have the same denominator and \( \kappa_{2m+1} \) is not part of their numerator. Therefore, algebraic independence of \( \kappa_2 \mu_{1,c'} \) and \( \kappa_{2m+1} \mu_{m,c} \)

\(^2\)a spanning tree is rooted at \( v \) if \( v \) is the only vertex with no outgoing edges
follows from the algebraic independence of the numerators of these two rational functions. Since \( \kappa_{2m+1} \) is a factor of \( \phi_{c \rightarrow c'}(\kappa) \) and is not part of the numerator of \( \phi_{c' \rightarrow c}(\kappa) \), the two functions \( \phi_{c \rightarrow c'}(\kappa) \) and \( \phi_{c' \rightarrow c}(\kappa) \) are algebraically independent over \( \mathbb{R} \).

(ii) We have \( \phi_{c_0 \rightarrow c_i}(\kappa) = \ell_i \mu_{m,c_0} \) for \( i = 1, \ldots, p \). Thus the set \( \{ \phi_{c_0 \rightarrow c_i} \mid 1 \leq i \leq p \} \) is algebraically independent over \( \mathbb{R} \) if and only if \( \{ \ell_i \mid 1 \leq i \leq p \} \) is, which clearly holds.

(iii) The reactions of the core network are of the form \( c_0 \rightarrow c_{i,j} \). We consider the graph \( \hat{G}_{c_0} \) (removing the edges for which there is no reaction):

\[
\kappa_1 \xrightarrow{\sum_{j=1}^{i_1} \ell_{1,j}} Y_1 \xrightarrow{\sum_{j=1}^{i_2} \ell_{2,j}} \ldots \xrightarrow{\sum_{j=1}^{i_m} \ell_{m,j}} Y_m
\]

We have \( \phi_{c_0 \rightarrow c_{i,j}}(\kappa) = \ell_{i,j} \mu_{i,c_0} \). The denominators of the rational functions \( \mu_{i,c_0} \) as given in (11) agree. Therefore it is enough to check that the polynomials \( \rho_{i,j} := \ell_{i,j} \sum_{\tau \in \Theta(Y_j)} \pi(\tau) \) for all \( i, j \) are algebraically independent over \( \mathbb{R} \).

For each \( 1 \leq i \leq m \), there exists a spanning tree rooted at \( Y_i \) involving an edge of the form \( Y_j \rightarrow * \) only for \( j \geq i \). Now consider the smallest index \( i \) such that there exists a complex \( c_{i,j} \). The parameter \( \ell_{i,j} \) appears in a polynomial \( \rho_{1,i_2} \) only for \( i_1 = i \). Hence the polynomials \( \rho_{1,i_2} \) are algebraically independent if and only if they are for \( i_1 > i \). We proceed in the same way now considering the smallest index \( k > i \) such that there exists a complex \( c_{k,j} \). This process terminates in at most \( m \) steps.

\[\square\]

Corollary 4.6 and Lemma 4.7(i) show that typical rational functions arising from realistic networks, such as those built from the mechanism in Example 2.2, fulfill the algebraic independence condition.

5. Another class of species: enzymes

In this final section we consider another class of species for which reduction mechanisms have also been defined, namely enzymes, and study how Gröbner bases of extended and reduced networks relate.

5.1. Enzymes. A species \( E \in S \) is an enzyme if for every reaction the stoichiometric coefficient of \( E \) in the reactant and the product agree [21]. This automatically gives that the steady state polynomial of \( E \) is identically zero, and implies that the concentration of \( E \) is constant in time and only depends on the initial amount \( e_0 \) of \( E \). For example, \( E \) and \( F \) are enzymes in the network of Example 2.2.

The core network obtained by removal of \( E \) consists of simply removing \( E \) from each side of the reaction (this is an example of an embedded network, see [17]). For
example, a reaction

\[(12) \quad X_1 + E \xrightleftharpoons{k_1} X_2 + E \quad \text{becomes} \quad X_1 \xrightarrow{k_1} X_2.\]

After fixing the initial amount of enzyme $e_0$, the steady states of the extended network satisfying that the concentration of $E$ is $e_0$ agree with the steady states of the core network with $k_1 = e_0k_1$.

This might lead one to think that enzymes are redundant and that similar properties as those that hold for intermediates also hold for enzymes. For example, one might think there is an easy way to obtain a Gröbner basis of the steady state ideal of the extended network from one of the core network, or that a binomial steady state ideal remains binomial upon removal of intermediates. But this is not the case, as the following examples illustrate.

**Example 5.1.** Let $\mathcal{N}$ be the network

\[
\begin{array}{c}
2X \xrightarrow{k_1} 3X \xrightarrow{k_2} X \\
\end{array}
\]

A binomial basis of the steady state ideal is $\{-2k_2x^3 + (k_1 - k_3)x^2\}$. Now consider the following network by adding one enzyme $E$:

\[
\begin{array}{c}
2X \xrightarrow{\kappa_1} 3X \xrightarrow{\kappa_2} X \\
2X + E \xrightarrow{\kappa_3} X + E.
\end{array}
\]

A reduced Gröbner basis of its steady state ideal is $\{x^3 - \frac{\kappa_1}{2k_2}x^2 + \frac{\kappa_3}{2k_2}x^2e\}$, and hence this ideal is not binomial.

The previous example suggests the following: Consider a reaction as in \((12)\). One might obtain a Gröbner basis of the steady state ideal of the extended network by considering a Gröbner basis of the steady state ideal of the core network and substituting the parameter $\kappa_1$ by $k_1e$. The following example gives a negative answer to this question.

**Example 5.2.** Let $\mathcal{N}$ be the following network

\[
\begin{array}{c}
X_1 \xrightarrow{k_1} 0 \quad 3X_1 \xrightarrow{k_3} X_2. \\
2X_1 \xrightarrow{k_2} \quad 
\end{array}
\]

The set of steady state polynomials is

\[
\{-k_1x_1 - 2k_2x_1^2 - 3k_3x_1^3, \ k_3x_2^3\}.
\]

With every arbitrary monomial order on $\mathbb{R}(k)[x]$, the reduced Gröbner basis of the steady state ideal is $\{x_1\}$.

Let now $\mathcal{N}'$ be the extension of $\mathcal{N}$ via the enzyme $E$:

\[
\begin{array}{c}
X_1 \xrightarrow{k_1} 0 \\
2X_1 \xrightarrow{k_2} \\
3X_1 + E \xrightarrow{k_3} X_2 + E.
\end{array}
\]
The set of steady state polynomials of $\mathcal{N}'$ is
$$\{-\kappa_1 x_1 - 2\kappa_2 x_1^2 - 3\kappa_3 x_1 e, \kappa_3 x_1^3 e\}.$$  

The steady state ideal is different from $\langle x_1 \rangle$. Thus, there is not a monomial order on $\mathbb{R}(\kappa)[x, e]$ for which the reduced Gröbner basis can be obtained from the set $\{x_1\}$ by making the substitution $k_3 = \kappa_3 e$.

**Example 5.3.** When a binomial basis of the steady state ideal is obtained from linear combinations of the steady state polynomials (see [4]), then the steady state ideal of the core network is binomial if and only if that of the extended network is.

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**References**


