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Best Laid Plans of Lions and Men

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Abstract
We answer the following question dating back to J.E. Littlewood (1885–1977): Can two lions catch a man in a bounded area with rectifiable lakes? The lions and the man are all assumed to be points moving with at most unit speed. That the lakes are rectifiable means that their boundaries are finitely long. This requirement is to avoid pathological examples where the man survives forever because any path to the lions is infinitely long. We show that the answer to the question is not always “yes” by giving an example of a region $R$ in the plane where the man has a strategy to survive forever. $R$ is a polygonal region with holes and the exterior and interior boundaries are pairwise disjoint, simple polygons. Our construction is the first truly two-dimensional example where the man can survive.

Next, we consider the following game played on the entire plane instead of a bounded area: There is any finite number of unit speed lions and one fast man who can run with speed $1 + \varepsilon$ for some value $\varepsilon > 0$. Can the man always survive? We answer the question in the affirmative for any constant $\varepsilon > 0$.

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1 Introduction

‘A lion and a man in a closed circular arena have equal maximum speeds. What tactics should the lion employ to be sure of his meal?’1 These words (including the footnote) introduce the now famous lion and man problem, invented by R. Rado in the late thirties, in Littlewood’s Miscellany [15]. It was for a long time believed that in order to avoid the lion, it was optimal for the man to run on the boundary of the arena. A simple argument then shows that the

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1 The curve of pursuit ($L$ running always straight at $M$) takes infinite time, so the wording has its point.

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lion could always catch the man by staying on the radius $OM$ defined by the man while approaching him as much as possible. However, A.S. Besicovitch proved in 1952 that the man has a very simple strategy (following which he will approach but not reach the boundary) that enables him to avoid capture forever no matter what the lion does. See [15] for details.

Throughout this paper, all men, lions, and other animals are assumed to be points. One can prove that two lions are enough to catch the man in a circular arena, and Croft [8] proves that in general a necessary and sufficient number of birds to catch a fly inside an $n$-dimensional spherical cage is just $n$ (again, we assume that the fly and the birds have equal maximum speeds).

A well-known related discrete game is the cop and robber game: Let $G$ be a finite connected undirected graph. Two players called cop $C$ and robber $R$ play a game on $G$ according to the following rules: First $C$ and then $R$ occupy some vertex of $G$. After that they move alternately along edges of $G$. The cop $C$ wins if at some point in time $C$ and $R$ are on the same vertex. If the robber $R$ can prevent this situation forever, then $R$ wins. The robber has a winning strategy on many graphs including all cycles of length at least 4. Therefore, the cop player $C$ can be given a better chance by allowing him, say, $k$ cops $C_1, \ldots, C_k$. At every turn $C$ moves any non-empty subset of $\{C_1, \ldots, C_k\}$. Now, the cop-number of $G$ is the minimal number of cops needed for $C$ to win. Aigner and Fromme [2] observes that the cop-number of the dodecahedron graph is at least 3, since if there are only 2 cops, the robber can always move to a vertex not occupied by a cop and not in the neighbourhood of any. Furthermore, they prove that the cop-number of any planar graph is at most 3. Thus, the cop-number of the dodecahedron is exactly 3.

Returning to the lion and man game, Bollobás [6] writes that the following open problem was already mentioned by J.E. Littlewood (1885–1977): Can two lions catch a man in a bounded (planar) area with rectifiable lakes? An informal definition of a rectifiable curve is that it has finite length. We require that the boundaries of the lakes and the exterior boundary are all rectifiable curves to avoid pathological examples where the man survives forever because any path to the lions is infinite. Bollobás mentions the same problem in a comment in his edition of Littlewood’s Miscellany [15] and in [7]. The problem is also stated by Fokkink et al. [11]. Berarducci and Intrigila [4] prove that the man can survive forever (for some initial positions of the man and lions) if the area is a planar embedding of the dodecahedron graph where each edge is a curve with the same length, say length 1. The proof is essentially the same as the proof by Aigner and Fromme [2] that the cop-number of the dodecahedron is at least 3: When the man is standing at a vertex, there will always be a neighbouring vertex with distance more than 1 to the nearest lion. It is thus safe for the man to run to that vertex. This, however, is a one-dimensional example. Berarducci and Intrigila raise the question whether it is possible to replace the one-dimensional edges by two-dimensional thin lines.

We present a truly two-dimensional region $R$ in the plane where two lions are not enough to ever catch the man. We say that $R$ is truly two-dimensional since $R$ is a polygonal region with holes and the exterior and interior boundaries are all pairwise disjoint, simple polygons – in particular, they are clearly rectifiable. We were likewise inspired by the dodecahedron in the construction of our example. We explain the construction in Section 2.

Rado and Rado [16] and Janković [13] consider the problem where there are many lions and one man, but where the game is played in the entire unbounded plane. They prove that the lions can catch the man if and only if the man starts in the interior of the convex hull of the lions. Inspired by that problem, we ask the following question: What if the lions have maximum speed 1 and the man has maximum speed $1 + \varepsilon$ for some $\varepsilon > 0$? We prove
that for any constant $\varepsilon$ and any finite number of lions, such a fast man can survive forever provided that he does not start at the same point as one of the lions. We explain a strategy in Section 3.

Other variants of the game with a faster man have been studied previously. Flynn [9, 10] and Lewin [14] study the problem where there is one lion and one fast man in a circular arena. The lion tries to get as close to the man as possible and the man tries to keep the distance as large as possible. Variants of the cop and robber game where the robber is faster than the cops have also been studied. See for instance [3, 12].

### 1.1 Definitions

We follow the conventions of Bollobás et al. [5]. Let $R \subseteq \mathbb{R}^2$ be a region in the plane on which the lion and man game is to be played, and assume that the lion starts at point $l_0$ and the man at point $m_0$. We define a man path as a function $m: [0, \infty) \to R$ satisfying $m(0) = m_0$ and the Lipschitz condition $\|m(s) - m(t)\| \leq V \cdot |s - t|$, where $V$ is the speed of the man. In our case, we either have $V = 1$ or, in the case of a fast man, $V = 1 + \varepsilon$ for some small constant $\varepsilon > 0$. Note that it follows from the Lipschitz condition that any man path is continuous. A lion path $l$ is defined similarly, but the lions we consider always run with at most unit speed. Let $\mathcal{L}$ be the set of all lion paths and $\mathcal{M}$ be the set of all man paths. Then a strategy for the man is a function $M: \mathcal{L} \to \mathcal{M}$ such that if $l, l' \in \mathcal{L}$ agree on $[0, t]$, then $M(l)$ and $M(l')$ also agree on $[0, t]$. This last condition is a formal way to describe that the man’s position $M(l)(t)$, when he follows strategy $M$, depends only on the position of the lion at points in time before and including time $t$, i.e., he is not allowed to act based on the lion’s future movements. (By the continuity of any man path, the man’s position at time $t$ is in fact determined by the lion’s position at all times strictly before time $t$.) A strategy $M$ for the man is winning if for any $l \in \mathcal{L}$ and any $t \in [0, \infty)$, it holds that $M(l)(t) \neq l(t)$. Similarly, a strategy for the lion $L: \mathcal{M} \to \mathcal{L}$ is winning if for any $m \in \mathcal{M}$, it holds that $L(m)(t) = m(t)$ for some $t \in [0, \infty)$. These definitions are extended to games with more than one lion in the natural way.

It might seem unfair that the lion is not allowed to react on the man’s movements when we evaluate whether a strategy $M$ for the man is winning. However, we can give the lion full information about $M$ and allow it to choose its path $l$ depending on $M$ prior to the start of the game. If $M$ is a winning strategy, the man can also survive the lion running along $l$.

We call a man strategy $M$ locally finite if it satisfies the following property: if $l$ and $l'$ are any two lion paths that agree on $[0, t]$ for some $t$ then the corresponding man paths $M(l)$ and $M(l')$ agree on $[0, t + \delta]$ for some $\delta > 0$ (we allow that $\delta$ depends on $l_{|[0,t]}$). Thus, informally, the man commits to doing something for some positive amount of time dependent only on the situation so far. Bollobás et al. [5] prove that if the man has a locally finite winning strategy, then the lion does not have any winning strategy. The argument easily extends to games with multiple lions. At first sight, it might sound absurd to even consider the possibility that the lion has a winning strategy when the man also does. However, it does not follow from the definition that the existence of a winning strategy for the man implies that the lion does not also have a winning strategy. See the paper by Bollobás et al. [5] for a detailed discussion of this (including descriptions of natural variants of the lion and man game where both players have winning strategies). In each of the problems we describe, the winning strategy of the man is locally finite, so it follows that the lions do not have winning strategies. In fact, the strategies we describe satisfy the much stronger condition that they are equitemporal, i.e., there is a constant $\Delta > 0$ such that the man at any point in time $i \cdot \Delta$, for $i = 0, 1, \ldots$, decides where he wants to run until time $(i + 1) \cdot \Delta$. 
In this section, we present a polygonal region $R$ in the plane with 11 lakes. See [1] for an illustration of such a region. The exterior and interior boundaries of $R$ are all pairwise disjoint simple polygons, and a man can survive forever in $R$ against two lions provided that the lions are initially at a sufficient distance.

Consider a planar embedding $D$ of the dodecahedron where each edge is a polygonal curve. We can obtain that all edges have the same length by prolonging some edges using a zig-zag pattern. This embedding corresponds to an area with 11 lakes and infinitely thin paths between the lakes, and as observed by Berarducci and Intrigila [4], the man can survive forever against two lions on such an embedding by deciding at each vertex which neighbouring vertex to visit next. First, we explain why it is not straightforward to obtain the region $R$ from $D$, or, at least, why some natural initial attempts will not work.

We want to “thicken” each edge of $D$ such that the boundaries of the lakes become disjoint, thus obtaining a truly two-dimensional region $D'$ containing $D$ as a subset. However, doing so, the point in $D'$ corresponding to a vertex of $D$ does not necessarily lie on the shortest path between its neighbours. We thus cannot simply employ the strategy from $D$, roughly speaking, because the man must plan in advance which turn to take in the upcoming vertex. Thus, before he reaches the region $R_v$ corresponding to a given vertex, he should already know which neighbouring vertex he will visit afterwards. Then, he can choose a path through $R_v$ that makes the concatenated path shortest possible.

In order to carry out this idea, we first need to describe a winning strategy of the man on the dodecahedron graph with the special property that he does not make his decisions at the vertices. Let $G$ be a planar embedding of the dodecahedron where all edges have length 4. The distance between two points in $G$ is the length of a shortest path between the points. Let the quarters denote the points on the edges of $G$ at distance 1 to the closest vertex. Consider a quarter $x$ on the edge $ab$ of $G$. For a point $p \in G$, $p \neq x$, let $d_a(x, p)$ be the length of a shortest simple path in $G$ from $x$ to $p$ that initially follows the edge $\{a, b\}$ in direction towards $a$. Let $d_b(x, p)$ be defined similarly.

When the man is at a quarter $x$ with distance 1 to the vertex $a$ and 3 to the vertex $b$, we let $d_{\text{near}}$ denote the distance from $x$ to the closest lion with respect to $d_a$, and let $d_{\text{far}}$ denote the distance from $x$ to the closest lion with respect to $d_b$. To avoid confusion, we write them as $d_{\text{near}}(t)$ and $d_{\text{far}}(t)$ when $x$ is the position of the man at the time $t$.

We will now show that if the lions are sufficiently far away in the initial situation, there exists a winning strategy for the man where he only takes stock of the situations in the quarters. That is, when he reaches a quarter, he must plan for the next 2 units of time where to run to, and then he has reached a quarter again, and so on.

**Invariant 1.** In the scenario described above:
1. The man is standing on a quarter.
2. $\min\{d_{\text{near}}, d_{\text{far}}\} \geq 1$.
3. At least one of the two following statements is true:
   - $d_{\text{near}} \geq 3$
   - $d_{\text{far}} \geq 7$

**Lemma 2.** If Invariant 1 is satisfied initially, the man has a winning strategy by which he runs from quarter to quarter at unit speed so that Invariant 1 is true at any quarter. The strategy maintains Invariant 1 Point 2 at all times, that is, that the closest lion is always at least at distance 1.
Proof. Let \( x \) denote the position of the man at the time \( t \), and assume the invariant holds. We prove that he can run to another quarter \( x' \) without getting caught such that the invariant again holds when he reaches \( x' \).

The proof goes by inspecting cases. Let \( ab \) be the edge containing \( x \) and suppose \( a \) be the nearest vertex to \( x \) and \( b \) the furthest.

**Case 1:** \( d_{\text{far}}(t) \geq 7 \). Let \( y \) denote the other quarter on the same edge as \( x \). We claim that the man can run to \( y \) without violating the invariant. We must thus argue that the invariants are satisfied at time \( t+2 \) for a man situated at \( y \). First, note that he will not encounter any lion while running towards \( y \) because \( d_{\text{far}}(t) > 4 \). Note also that \( d_{\text{far}}(t + 2) \geq 1 \), since \( d_{\text{near}}(t) \geq 1 \) and the worst case is that the lion follows the man. Furthermore, \( d_{\text{near}}(t + 2) \geq 7 - 4 = 3 \), since \( d_{\text{far}}(t) \geq 7 \) and the worst case is that the man and lion have run towards each other. Thus, the invariant holds at the time \( t + 2 \).

**Case 2:** \( d_{\text{far}} < 7 \), and thus \( d_{\text{near}}(t) \geq 3 \). In this case, we exploit the fact that \( d_{\text{far}} \) is so small that we can bound \( d_{\text{far}}(t + 2) \) from below. Let \( l_{\text{far}} \) denote the lion at distance \( d_{\text{far}} \) from \( x \), and let \( l_{\text{near}} \) denote the other lion. Consider the two other quarters at distance 1 from \( a \), call them \( q_1 \) and \( q_2 \). Assume without loss of generality that \( q_1 \) is furthest from \( l_{\text{near}} \). The situation is sketched in Figure 1. We now argue that the man can choose to run towards \( q_1 \) without getting eaten, and while maintaining the invariant. Let \( b' \) denote the vertex at distance 3 to \( q_1 \). Note that \( d_{b'}(q_1, l_{\text{near}}(t)) \geq 11 \) and thus, \( d_{b'}(q_1, l_{\text{near}}(t + 2)) \geq 9 \).

In Figure 1, the points that are both \( \geq 3 \) from \( x \), and (weakly) closer to \( q_2 \) than to \( q_1 \), are marked with red, and hence by our choice of \( q_1 \), \( l_{\text{near}} \) must be in the subset marked with red at time \( t \). As is easily seen by inspection, \( d_{b'}(q_1, l_{\text{near}}(t)) \geq 9 \), and thus \( d_{b'}(q_1, l_{\text{near}}(t + 2)) \geq 7 \). But then, \( d_{\text{far}}(t + 2) \geq \min\{9, 7\} = 7 \), and Invariant 1.1 and 3 are maintained.

To see that Invariant 1.2 is still maintained, note that \( d_{a}(q_1, l_{\text{near}}(t)) \geq 3 \) and therefore \( d_{a}(q_1, l_{\text{near}}(t + 2)) \geq 1 \). Similarly, since \( d_{b}(x, l_{\text{far}}(t)) \geq 1 \), we have \( d_{a}(q_1, l_{\text{far}}(t)) \geq 3 \) so that \( d_{a}(q_1, l_{\text{far}}(t + 2)) \geq 1 \). Thus, \( l_{\text{near}}(t + 2) \geq 1 \), and we are done.

\[
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\]

Our first goal is to find an embedding \( G \) of the dodecahedron in the plane with the properties described below, which will make it easier for us to construct the region \( R \).
Regardless of angles between $a, b, c$, we can introduce bends to make the three edges meet at $v$ in angles of size $\frac{3\pi}{2}$ and at the same time extend the lengths suitably.

The shortest paths in the circle $D_v$ between any two of $a, b, c$, that avoid crossing the polygonal curves $P_{vf}, P_{vg}, P_{vh}$ all have length $\frac{1}{8}$.

Lemma 3. There exists a planar embedding $G$ of the dodecahedron such that

- all edges have length 4,
- all edges consist of line segments with lengths being multiples of $\frac{1}{5}$,
- any pair of line segments from different edges that meet at a vertex each have length $\frac{1}{4}$ and form an angle of size $\frac{2\pi}{3}$, and
- for any vertex $v$, the circle $D_v$ centered at $v$ with radius $\frac{1}{16}$ only intersects the three edges incident to $v$.

After proving this lemma, we derive from $G$ a truly two-dimensional area $R$ in the plane where the man can survive against two lions. Lemma 2 gives a winning strategy for the man in $G$ where he runs from quarter to quarter. The paths along which he runs in $R$ will be exactly the same as in $G$ except for inside the circles $D_v$.

We first need the following elementary geometric observations:

Observation 4. There exists a planar embedding $D$ of the dodecahedron such that all edges have length 1 or 3. $D$ furthermore has the property that the circle of radius $\frac{1}{4}$ centered at any vertex $v$ only intersects the three edges incident to $v$. (See Figure 2.)

Lemma 5. For any three points $a, b, c$ on a circle $C$, there exist an equilateral triangle with corners $a', b', c'$ on $C$ where $\{a, b, c\}$ and $\{a', b', c'\}$ are disjoint and such that, when considering the points $a, b, c, a', b', c'$ all together, $a$ is a neighbour of $a'$, and $b$ is a neighbour of $b'$, and $c$ is a neighbour of $c'$.

Proof. See Figure 3. The points $a, b, c$ divide $C$ into three arcs. Clearly, we can choose an equilateral triangle with corners on $C$ disjoint from $\{a, b, c\}$ so that not all three corners of the triangle are on the same arc. It is now easy to label the corners of the triangle with $a', b', c'$ to satisfy the lemma.

We are now ready to prove that a planar embedding $G$ of the dodecahedron exists as stated in Lemma 3.

Proof of Lemma 3. Start with the embedding $D$ shown in Figure 2, where all edges have length 1 or 3. Consider a vertex $v$ and the circle $C_v$ of radius $r = \frac{1}{4}$ centered at $v$. Assume the three edges incident to $v$ enter $C_v$ in the points $a, b, c$, and let $u_a, u_b, u_c$ be the neighbouring
vertices of $v$ such that $a$ is a point on the edge $\{u_a, v\}$, $b$ is a point on $\{u_b, v\}$, and $c$ is a point on $\{u_c, v\}$. We now delete the segments $va$, $vb$, and $vc$, and therefore need to reconnect $a$, $b$, and $c$ to $v$. We explain how to reconnect $a$ to $v$; $b$ and $c$ are handled analogously. We find points $a', b', c'$ on $C_v$ as described in Lemma 5. See Figure 3. We first connect $a'$ to $v$. We now need to connect $a$ to $a'$ using some bends. A bend is two segments $xy$ and $yz$, each of length $r/2 = 1/8$, such that $x$ and $z$ are on $C_v$ and $y$ is in the interior of $C_v$. If the edge $\{u_a, v\}$ had length 3 in $D$, we make two bends that together connect $a$ and $a'$. We thus increase the length of the edge $\{u_a, v\}$ by $1/2$ in each end and the resulting edge has length 4. If the edge $\{u_a, v\}$ had length 1 in $D$, we connect $a$ and $a'$ by 6 bends, corresponding to extending the length of the edge by 3. The result is a planar embedding $\mathcal{G}$ of the dodecahedron with the properties stated in the lemma.

We now describe how to make the region $R$. We want each quarter of $\mathcal{G}$ to be a point in $R$ and we want all pairs of quarters to have the same distances in $\mathcal{G}$ and $R$. It will then follow from Lemma 2 that the man has a winning strategy by running from quarter to quarter in $R$. We make one lake $L_f$ corresponding to each face $f$ of $\mathcal{G}$. Here, we also consider the outer boundary of $R$ to be the boundary of an unbounded lake corresponding to the exterior face of $\mathcal{G}$. The shortest paths in $R$ will be polygonal paths with corners at convex corners of the lakes. Outside the circles $D_v$, the paths along which the man will run are exactly the paths in $\mathcal{G}$. Inside a circle $D_v$, we need to take special care to ensure that the man can always run along an optimal path.

We now explain the construction of the lakes $L_f$ corresponding to faces $f$ of $\mathcal{G}$. Consider a vertex $v$ of $\mathcal{G}$ and the faces $f, g, h$ on which $v$ is a vertex. We first describe how the boundaries of $L_f, L_g, L_h$ look in the circle $D_v$ of radius 1/16 centered at $v$. See Figure 4. Let $a, b, c$ be the points where the edges incident to $v$ enter $D_v$. Suppose that the arc on $D_v$ from $a$ to $b$ is in the face $f$, the arc from $b$ to $c$ is in $g$, and the arc from $c$ to $a$ is in $h$. We now create three polygonal curves $P_{vf}, P_{vg}, P_{vh}$ inside $D_v$ so that the shortest path between any two of $a, b, c$ contained in $D_v$ and not crossing any of $P_{vf}, P_{vg}, P_{vh}$ has length $1/8$. The curve $P_{vf}$ starts at a point $r_{vf}$ on $D_v$ and ends at a point $s_{vf}$ on $D_v$, and the endpoints $r_{vf}, s_{vf}$ are inside $f$, and similarly for the faces $g, h$. These properties are easy to obtain by a construction as shown in Figure 4. The curves $P_{vf}, P_{vg}, P_{vh}$ will be part of the boundary of the lakes $L_f, L_g, L_h$, respectively.

\[ Figure 5 \] The edge $e_{uv}$ of $\mathcal{G}$ is red and is one of the edges bounding the face $f$, which is above $e_{uv}$. The polygonal curve $Q_{uv}$, which is on the boundary of the lake $L_f$, is blue.
We now explain how to construct the rest of the boundary of each lake \( L_f \). Consider a face \( f \) of \( \mathcal{G} \) and assume that the vertices on \( f \) are \( uvxyz \) in that order on \( f \). The curves \( P_{u,v}, P_{v,f}, P_{f,x}, P_{x,y} \) appear on the boundary of \( L_f \) in that order. In the following, we describe how to connect the end \( s_{af} \) of \( P_{af} \) with the start \( r_{vf} \) of \( P_{vf} \) – the other curves are connected in a completely analogous way. See Figure 5. Let \( e_{uv} \) be the edge of \( \mathcal{G} \) between \( u \) and \( v \), thus, \( e_{uv} \) is a polygonal curve. Let a corner of \( e_{uv} \) be a common point of two neighbouring segments of \( e_{uv} \). We make a polygonal curve \( Q_{uv} \) corresponding to \( e_{uv} \). \( Q_{uv} \) starts at \( s_{af} \) and ends at \( r_{vf} \) so that it connects \( P_{af} \) and \( P_{vf} \). \( Q_{uv} \) stays near \( e_{uv} \) inside \( f \) and touches \( e_{uv} \) at the corners of \( e_{uv} \) which are convex corners of \( f \). To summarize, \( Q_{uv} \) has the following properties:

1. \( Q_{uv} \) starts at \( s_{af} \) and ends at \( r_{vf} \).
2. \( Q_{uv} \) is completely contained in \( f \).
3. \( Q_{uv} \) is, except for the endpoints \( s_{af}, r_{vf} \), outside the circles \( D_u \) and \( D_v \).
4. \( Q_{uv} \) and \( Q_{uv'} \) are completely disjoint for any ordered pair \( (u', v') \neq (u, v) \) so that \( \{u', v'\} \) is an edge of \( \mathcal{G} \), and \( e_{uv} \) touches \( e_{uv} \) at a point \( p \) if and only if \( p \) is a corner of \( e_{uv} \) which is a convex corner of \( f \).

Observe that \( Q_{uv} \) (note: not \( Q_{uv}! \)) touches \( e_{uv} \) at the corners which are concave corners of \( f \), since those are convex corners of the neighbouring face on the other side of \( e_{uv} \).

**Theorem 6.** There exists a polygonal region \( R \) in the plane with holes where the exterior and interior boundaries are all pairwise disjoint and such that the man has a winning strategy against two lions.

**Proof.** \( R \) is the region that we get by removing from \( \mathbb{R}^2 \) the interior of each of the lakes \( L_f \). Thus, the boundary of each lake is included in \( R \), so that \( R \) is a closed set. \( R \) is also bounded because we remove the interior of the unbounded lake corresponding to the exterior face of \( \mathcal{G} \). Note that any point on an edge \( e_{uv} \) of \( \mathcal{G} \) which is outside the circles \( D_u \) and \( D_v \) is a point in \( R \). Since the quarters of \( e_{uv} \) are outside the circles \( D_u \) and \( D_v \), it follows that they are also points in \( R \). Furthermore, our construction ensures that the distance in \( R \) between any two quarters is the same as in \( \mathcal{G} \). Let \( \mathcal{G}' \) be the points in \( R \) which are on some shortest path between two quarters in \( R \). Thus, \( \mathcal{G}' \) are the points that the man can possibly visit when running along shortest paths in \( R \) from quarter to quarter.

Let \( l_1 \) and \( l_2 \) be two lions in \( R \). We define projections \( l'_1 \) and \( l'_2 \) of the lions \( l_1 \) and \( l_2 \) to be the closest points in \( \mathcal{G}' \) (with respect to distances in \( R \)). We now define \( l''_1 \) and \( l''_2 \) to be projections of \( l'_1 \) and \( l'_2 \) in \( \mathcal{G} \) in the following way. Outside the circles \( D_u \), \( \mathcal{G} \) and \( \mathcal{G}' \) coincide, and here we simply define \( l''_1 := l'_1 \). Suppose now that \( l'_1 \) is inside a circle \( D_v \) for some vertex \( v \) of \( \mathcal{G} \). See Figure 6. Suppose that the three edges incident to \( v \) enter \( D_v \) at the points \( a, b, c \). The projection \( l''_1 \) is a point on one of the shortest paths between a pair of the points \( a, b, c \). Recall that these shortest paths all have length 1/8. Assume without loss of generality that \( l''_1 \) is on the path from \( a \) to \( c \). Let \( d \) be the distance from \( a \) to \( l''_1 \) in \( R \), so that \( 0 \leq d \leq 1/8 \). If \( d = 1/16 \), we define \( l''_1 := v \). Otherwise, if \( d < 1/16 \), we let \( l''_2 \) be the point on the segment \( av \) in \( \mathcal{G} \) with distance \( d \) to \( a \), i.e., \( l''_1 \in av \) so that \( ||al''_1|| = d \). Similarly, if \( d > 1/16 \), we let \( l''_1 \) be the point on \( bv \) with distance \( 1/8 - d \) to \( b \).

We now prove that \( l''_1 \) moves at at most unit speed in \( \mathcal{G} \). It will then follow from Lemma 2 that the man has a winning strategy.

\( \mathcal{G}' \) subdivides \( R \) into some regions \( R'_1, \ldots, R'_k \), which are the connected components of \( R \setminus \mathcal{G}' \). Let \( R_i = \overline{R'_i} \) be the closure of \( R'_i \). Now, \( R = \bigcup_{i=1}^k R_i \). Inside each circle \( D_v \), there is a triangular region bounded by three segments from \( \mathcal{G}' \). All other regions are bounded by
a polygonal curve $C \subset \partial L_f$ on the boundary of some lake $L_f$ and a concave chain $H \subset G'$. Call such a region normal. If the lion $l_i$ is in a normal region $R_j$ with boundary $\partial R_j = C \cup H$ as described before, the projection $l'_i$ is on $H$. It then follows from the concavity of $H$ that $l'_i$, and thus also $l''_i$, moves continuously and with at most unit speed.

However, when $l_i$ is inside a triangular region in $D_v$, the projection $l'_i$ might jump from one segment of the triangle to another. Suppose that the three edges incident to $v$ enter $D_v$ at the points $a, b, c$ as in Figure 6. Let $a'$ be the point where the shortest paths from $a$ to $b$ and $c$ separate and define $b'$ and $c'$ similarly. Thus, the points $a'b'c'$ are the corners of the triangular region. Suppose that $l'_i$ jumps from $a'b'$ to $a'c'$. Then, the distance from $l_i$ to $a'b'$ and $a'c'$ is the same and the distance from $a$ to $l'_i$ before and after the jump is at most $1/16$, since otherwise, $l_i$ would be closer to the segment $b'c'$ than to $a'b'$ and $a'c'$. It follows that $l'_i$ jumps from one point to another which have the same projection $l''_i$. Thus, $l''_i$ moves continuously and with at most unit speed.

The man now employs the strategy from Lemma 2 in the following way. He imagines that he is playing in the dodecahedron $G$ against the lions $l''_1$ and $l''_2$. Assume therefore that Invariant 1 holds initially. The strategy tells the man to which neighbouring quarter to run. That quarter also exists in $G'$, and has the same distance, so the man runs to that quarter in $G'$. Since $l''_1$ and $l''_2$ run with at most unit speed, the man can escape them forever. When the man is outside the circles $D_v$, it is a necessary condition for the lions to catch the man that $l''_1$ or $l''_2$ coincide with the man, so we conclude that they cannot catch him outside the circles. When the man is inside a circle $D_v$, we know from Lemma 2 that $l''_1$ and $l''_2$ are at least 1 away from the man. Therefore, $l_1$ and $l_2$ must be outside $D_v$, and hence they cannot catch him in that case either. Thus, the man survives forever in $R$. ◀

3 The Fast Man Surviving any Number of Lions in the Plane

Finally, we consider the case where the man is just slightly faster than the lions in the unbounded plane without obstacles. In this case, the man is able to escape arbitrarily many lions. The full proofs of some of the claims below can be found in [1].

Theorem 7. In the plane $\mathbb{R}^2$, for any $\varepsilon > 0$, a man able to run at speed $1 + \varepsilon$ has a locally finite strategy to escape the convex hull of any number $n \in \mathbb{N}$ of unit-speed lions, provided
that the man does not start at the same point as a lion. Thus, the man has a locally finite winning strategy.

In fact, we prove that the man is able to keep some minimum distance \( d_{\varepsilon,n} \) to any lion, where \( d_{\varepsilon,n} \) only depends on \( \varepsilon, n \), and the initial distances to the lions. Thus, if the \( n \) lions and man were disks with radius \( \frac{1}{2}d_{\varepsilon,n} \), the man is still able to escape.

We proceed by induction on the number \( n \) of lions. We define strategies \( M_i \) for the man to keep distance \( c_j \) to the first \( j \) lions. The \( j \)th strategy yields a curve consisting of line segments all of the same length.

Inductively, the man can keep a safety distance \( c_{n-1} \) to the \( n-1 \) first lions by running at speed \( 1 + \varepsilon_{n-1} \), where \( \varepsilon_1 < \varepsilon_2 < \ldots < \varepsilon_n < \varepsilon \). The bends of the curve defined by strategy \( M_{n-1} \) are milestones that he runs towards when avoiding \( n \) lions. If the \( n \)th lion \( \ell_n \) is in the way, the man makes an avoidance move, keeping a much smaller safety distance \( c_n \) to \( \ell_n \) and running slightly faster at speed \( \varepsilon_n \) (see Figure 9). Intuitively, when performing avoidance moves, the man runs counter-clockwise around a fixed-diameter circle centered at the lion.

After a limited number of avoidance moves, the man can make an escape move, where he simply runs towards the milestone defined by the strategy \( M_{n-1} \).

By choosing \( c_n \) sufficiently small, we can make sure that the detour caused by the \( n \)th lion is so small that it can only annoy the man once for each of the segments of the strategy \( M_{n-1} \), and thus that he is ensured to have distance at least \( c_{n-1}/2 \) to the position defined by \( M_{n-1} \) and hence not in danger of the \( (n-1) \)st lions.

\textbf{Theorem 8.} A man able to run at speed \( 1 + \varepsilon \) for any \( \varepsilon > 0 \) has a locally finite strategy to escape the convex hull of any number \( n \in \mathbb{N} \) of unit-speed lions, provided that the man does not start at the same point as a lion. Thus, the man has a locally finite winning strategy.

\textbf{Proof.} We assume without loss of generality that \( \varepsilon < 1 \). Let \( l_1, \ldots, l_n \) be \( n \) arbitrary lion paths and let the man start at position \( m_0 \) such that \( m_0 \not\in l_i(0) \) for all \( i \). We show that the man has a strategy \( M_n \) with the following properties:

1. The man is always running at speed \( 1 + \varepsilon_n \), where \( \varepsilon_n := (1 - 2^{-n}) \cdot \varepsilon \).
2. The path defined by \( M_n(l_1, \ldots, l_n) \) is a polygonal path with corners \( m_0 m_1 \ldots \) and each segment \( m_i m_{i+1} \) has the same length \( \Delta_n \cdot (1 + \varepsilon_n) \). Thus, the time it takes the man to run from \( m_i \) to \( m_{i+1} \) is \( \Delta_n \).
3. Let \( t_i := i \cdot \Delta_n \) be the time where the man leaves \( m_i \) in order to run to \( m_{i+1} \). The point \( m_{i+1} \) can be determined from the positions of the lions at time \( t_i \).
4. There exists a safety distance \( c_n > 0 \) such that for any \( i = 1, \ldots, n \), any \( t \in [t_i, t_{i+1}] \), and any point \( x \in m_im_{i+1} \), it holds that \( \text{dist}(x, \{l_1(t), \ldots, l_n(t)\}) \geq c_n \).
5. There is a corner \( m_i = M_n(t_i) \) such that for all \( t \geq t_i \),

\[ M_n(l_1, \ldots, l_n)(t) \not\in \text{CH}\{l_1(t), \ldots, l_n(t)\}. \]

Clearly, it follows from the properties that \( M_n \) is a winning strategy for the man fulfilling the requirements in the theorem. We prove the statement by induction on \( n \). If there is only one lion, the man will run on the same ray all the time with constant speed \( 1 + \varepsilon_1 = 1 + \varepsilon/2 \). The man chooses the direction of the ray to be \( m_0 = l_1(0) \). This strategy obviously satisfies the stated properties. Assume now that a strategy \( M_{n-1} \) with the stated properties exists for \( n-1 \geq 1 \) lions and consider a situation with \( n \) lions running along paths \( l_1, \ldots, l_n \).

Assume without loss of generality that the lions are numbered according to their (increasing) distance to the man at time 0, i.e., \( \|m_0 l_1(0)\| \leq \|m_0 l_2(0)\| \leq \cdots \leq \|m_0 l_n(0)\| \). For any \( i \in \{1, \ldots, n\} \), let \( M_i \) be shorthand for \( M_i(l_1, \ldots, l_i) \) and \( M \) shorthand for \( M_n \).
At any time $t$, let the succeeding corner on the strategy $M_{n-1}$ be

$$g(t) := M_{n-1} (\lceil t/\Delta_{n-1} \rceil \cdot \Delta_{n-1}).$$

By property 3, the man can always compute the point $g(t)$.

We first describe the intuition behind the man’s strategy without specifying all details, and later give a precise description. In the situation with $n$ lions, the man attempts to run according to the strategy for the $n-1$ first lions, i.e., the strategy $M_{n-1}$. Thus, at any time $t$, the man’s goal is to run towards the point $g(t)$. However, the lion $l_n$ might prevent him from doing so. Compared to the case with $n-1$ lions, the man has increased his speed by $1 + \varepsilon_n - (1 + \varepsilon_{n-1}) = 2^{-n}\varepsilon$, so he has time to take detours while still following the strategy $M_{n-1}$ approximately.

Assume that we have defined the man’s strategy up to time $t$. If he is close to the $n$’th lion, i.e., the distance $\|m(t) - l_n(t)\|$ is close to $r$, for some small constant $r > 0$ to be specified later, he runs counterclockwise around the lion, maintaining approximately distance $r$ to the lion. He does so until he gets to a point where running directly towards $g(t)$ will not decrease his distance to the lion. He then escapes from the lion, running directly towards $g(t)$. Doing so, he can be sure that the lion cannot disturb him anymore until he reaches $g(t)$ or $g(t)$ has changed.

We choose $r$ so small that when the man is running around the lion, we are in one of the following cases:

- The lion is so close to $g(t)$ that the man is within the safety distance $c_{n-1}$ from $g(t)$, and thus in no danger of the lions $l_1, \ldots, l_{n-1}$.
- After running around the lion in a period of time no longer than $12\pi r / \varepsilon_n$, the man escapes by running directly towards $g(t)$ without decreasing the distance to the lion. By choosing $r$ sufficiently small, we can therefore limit the duration, and hence the length, of the detour that the lion can force the man to run, so that the man is ensured to be within the safety distance from the lions $l_1, \ldots, l_{n-1}$ during the detour.

We now describe the details that make this idea work. We define

$$r := \min \left\{ \frac{\Delta_{n-1}\varepsilon_n(\varepsilon_n - \varepsilon_{n-1})}{2 + 2\varepsilon_n + 18\pi(1 + \varepsilon_n)} , \frac{\varepsilon_n c_{n-1}}{4 + 4\varepsilon_n + 24\pi(1 + \varepsilon_n)} \right\},$$

$$\rho := 2r / \varepsilon_n,$$

$$\theta := \arccos \frac{1}{1 + \varepsilon_n},$$

$$\varphi \in (0, \pi/2]$$

so that

$$\tan \theta = \frac{\rho \sin \varphi}{\rho \cos \varphi - 2r},$$

and

$$\Delta_n > 0 \text{ so that } 2 \arcsin \frac{(1 + \varepsilon_n)\Delta_n}{2(r - \Delta_n)} + \frac{\Delta_n}{\rho} \leq \varphi, \Delta_n < \frac{r}{3 + \varepsilon_n}, \text{ and } \Delta_{n-1}/\Delta_n \in \mathbb{N}.$$
Suppose now that \( \|m(t_i)l_n(t_i)\| < r + \Delta_n \) and consider the case \( m(t_i) \neq g(t_i) \). Let \( b \) be the point at distance \((1 + \varepsilon_n)\Delta_n\) from \( m(t_i) \) in the direction towards \( g(t_i) \). If there exist two parallel lines \( W_0 \) and \( W_1 \) such that \( m(t_i) \in W_0, b \in W_1, \text{dist}(l_n(t_i), W_0) \geq r - \Delta_n \), and \( \text{dist}(l_n(t_i), W_1) \geq \text{dist}(l_n(t_i), W_0) + \Delta_n \), then the man runs to \( b \).

In the remaining cases, the circles \( C(m(t_i), \Delta_n(1 + \varepsilon_n)) \) and \( C(l_n(t_i), r) \) intersect at two points \( p \) and \( q \) such that the arc on \( C(l_n(t_i), r) \) from \( p \) counterclockwise to \( q \) is in the interior of \( C(m(t_i), \Delta_n(1 + \varepsilon_n)) \). The man then runs towards the point \( q \).

A move defined by case A, B, or C is called a free move, an escape move, or an avoidance move, respectively. Let move \( i \) be the move that the man does during the interval \([t_i, t_{i+1})\).

\[ \text{Claim 9.} \text{ At any time of choice } t_i, \text{ it holds that} \]
\[ \|m(t_i)l_n(t_i)\| \geq r - \Delta_n \]
and if the preceding move was an avoidance move, it also holds that
\[ \|m(t_i)l_n(t_i)\| \leq r + \Delta_n. \]

Furthermore, at an arbitrary point in time \( t \in [t_{i-1}, t_i] \) and any point \( m' \in m([t_{i-1}, t_i]) \) it holds that
\[ 0 < r - (3 + \varepsilon_n)\Delta_n \leq \|m'l_n(t)\| \]
and if move \( i - 1 \) is an avoidance move then additionally
\[ \|m'l_n(t)\| \leq r + (3 + \varepsilon_n)\Delta_n. \]

\[ \text{Proof.} \text{ (Sketch) The proof is by induction on } i. \text{ The induction step for the first inequality follows easily from the speed of the man and the speed of the lion } l_n \text{ and (in case of an escape or avoidance move) from considering the distance } \|m(t_i)l_n(t_{i-1})\|. \text{ For the second inequality, note that } \|m(t_i)l_n(t_{i-1})\| = r. \text{ The third inequality follows by considering the combined speed of the man and the lion } l_n \text{ and by observing that } \|m(t_{i-1})l_n(t_{i-1})\| \geq r - \Delta_n. \text{ A similar argument using the inequality } \|m(t_{i-1})l_n(t_{i-1})\| \leq r + \Delta_n \text{ shows the fourth inequality.} \]
When the man does an escape move, he will not do an avoidance move before he reaches \( w \) or time has passed: (i) which proves that the man runs from parallel lines defined in case B for move \( w \). Finally, since \( W \parallel\) moves to \( g \) and passing through \( w \). We claim that for any \( j \in \{0, \ldots, k-1\} \), the man moves from \( w_j \) to \( w_{j+1} \) during move \( i+j \) using either an escape move or a free move. We prove this by induction on \( j \). It holds for \( j = 0 \) by assumption, so assume it holds that \( m(t_{i+j}) = w_j \) and that move \( i+j-1 \) was an escape move or a free move. Since the distance between consecutive lines \( W_j \) and \( W_{j+1} \) is at least \( \Delta_n \), it follows that dist \( (l_n(t_j), W_j) \geq r + (j-1)\Delta_n \) and hence that dist \( (l_n(t_{i+j}), W_j) \geq r - \Delta_n \). Now, if \( \|m(t_{i+j})l_n(t_{i+j})\| < r + \Delta_n(1 + \varepsilon_n) \), then the lines \( W_j \) and \( W_{j+1} \) are a witness that move \( i+j \) is an escape move so that the man moves to \( w_{j+1} \). Otherwise, move \( i+j \) is a free move, in which case the man moves to \( w_{j+1} \). Finally, since \( g(t) \) moves or the man reaches \( g \) during move \( i+k \), the statement holds. 

Define \( \rho' := \rho + r + (3 + \varepsilon_n)\Delta_n \) and \( \tau := 6\pi r/\varepsilon_n \).

\textbf{Claim 11.} If move \( i \) is an avoidance move, one of the following events occurs before \( \tau \) time has passed: (i) \( g(t) \) moves, (ii) \( \|m(t)g(t)\| < \rho' \), or (iii) the man makes an escape move.

\textbf{Proof Sketch.} If the first two events do not occur, it follows from Claim 10 that the man keeps doing avoidance moves during this time. Let \( \xi(t) \) resp. \( \eta(t) \) denote the angle of the vector \( l_n(t)m(t) \) resp. \( l_n(t)g(t) \). A key observation is that if the difference in these angles is small, the man makes an escape move since then the lion and the goal \( g \) are roughly on opposite sides of the man. Showing that this difference eventually becomes small involves showing that \( \eta \) increases by at least \( 2\pi \) more than \( \xi \) after \( \tau \) time so that at some point in time \( t \in [t_i, t_i + \tau] \), vectors \( l_n(t)m(t) \) and \( l_n(t)g(t) \) have the same orientation. By Claim 9,
the lion \( l_n \) never gets closer than \( \rho \) to \( g(t_i) \) which implies that the change in \( \eta \) is small in any time interval \([t_j, t_{j+1}]\). Since the man keeps a minimum distance to the lion, it similarly follows that the change in \( \xi \) is small in \([t_j, t_{j+1}]\). Picking \( j \) to be the maximum such that \( t_j \leq t \) gives \( t - t_j \leq \Delta_n \) which implies that the difference in the two angles is small at time \( t_j \) at which point the man makes an escape move. Since \( t_j \leq t + \tau \), the lemma follows. ▶

For \( i \in \mathbb{N}_0 \), define the canonical interval \( I_i \) as \( I_i := ([\Delta_{n-1}, (i+1)\Delta_{n-1}]) \), i.e., \( I_i \) is the interval of time where the man would run from the \( i \)'th to the \( (i+1)' \)st corner on the polygonal line defined by the strategy \( M_{n-1} \). We say that \( I_i \) ends at time \( t = (i+1)\Delta_{n-1} \). Note that if \( t \in I_i \), then \( g(t) = M_{n-1}((i+1)\Delta_{n-1}) \) and \( g(t) \) moves when \( I_i \) ends.

As a consequence of Claim 10 and Claim 11, we get the following.

**Claim 12.** If \( t \in I_i \) and \( \|m(t)g(t)\| \leq \rho' \), then for every \( t' > t, t' \in I_i \), we have
\[
\|m(t')g(t')\| \leq \rho' + (1 + \varepsilon_n)\tau.
\]

**Claim 13.** For any \( i \in \mathbb{N}_0 \) and at any time during the canonical interval \( I_i \), the man is at distance at most \( \rho' + 2(1 + \varepsilon_n)\tau \) away from the segment \( M_{n-1}(I_i) \) and when \( I_i \) ends, the man is within distance \( \rho' + (1 + \varepsilon_n)\tau \) from the endpoint \( M_{n-1}((i+1)\Delta_{n-1}) \) of the segment.

**Proof.** We prove the claim by induction on \( i \). To easily handle the base-case, we introduce an auxiliary canonical interval \( I_{-1} = [-\Delta_{n-1}, 0] \) and assume that the lions and the man are standing at their initial positions during all of \( I_{-1} \). The statement clearly holds for \( i = -1 \).

Assume inductively that the statement holds for \( I_{i-1} \) and consider the interval \( I_i \). Let \( g := M_{n-1}((i+1)\Delta_{n-1}) \). The additional distance that the man runs during \( I_i \) when his speed is \( 1 + \varepsilon_n \) as compared to the speed \( 1 + \varepsilon_{n-1} \) is \( \Delta_{n-1}(\varepsilon_n - \varepsilon_{n-1}) \). It follows from the definition of \( r \) that
\[
\Delta_{n-1}(\varepsilon_n - \varepsilon_{n-1}) \geq \rho + 2\tau + 3(1 + \varepsilon_n)\tau > \rho' + 3(1 + \varepsilon_n)\tau.
\]

By the induction hypothesis, the man is within a distance of \( \rho' + (1 + \varepsilon_n)\tau \) from \( M_{n-1}(i\Delta_{n-1}) \) at time \( i\Delta_{n-1} \). Thus, his distance to \( g \) at the beginning of interval \( I_i \) is at most \( \Delta_{n-1}(1 + \varepsilon_{n-1}) + \rho' + (1 + \varepsilon_n)\tau \), where \( \Delta_{n-1}(1 + \varepsilon_{n-1}) \) is the length of the interval \( M_{n-1}(I_i) \). If the man does not do any avoidance moves during \( I_i \), he runs straight to \( g \), so it follows that he reaches \( g \) at time \( (i+1)\Delta_{n-1} - 2\tau \) at the latest. Therefore, the statement is clearly true in this case.

Otherwise, let \( t \in I_i \) be the first time of choice at which he does an avoidance move during \( I_i \). If he is at distance at most \( \rho' \) from \( g \) at time \( t \), the statement follows from Claim 12. Therefore, assume that the distance is more than \( \rho' \). Then, we must have that \( t < (i+1)\Delta_{n-1} - 2\tau \), since, if \( t \) was larger, he would already have reached \( g \) by the above discussion. Hence, Claim 11 gives that at some time \( t' \leq t + \tau \), either
1. the man gets within a distance of \( \rho' \) from \( g \), or
2. he does an escape move.

We first prove that in the interval \([t, t']\), the distance from the man to the segment \( M_{n-1}(I_i) \) is at most \( \rho' + 2(1 + \varepsilon_n)\tau \). To this end, note that his distance to the segment at time \( t \) is at most \( \rho' + (1 + \varepsilon_n)\tau \). Thus, since \( t' \leq t + \tau \), his distance at time \( t' \) can be at most \( \rho' + 2(1 + \varepsilon_n)\tau \).

It remains to be proven that the man stays within distance \( \rho' + 2(1 + \varepsilon_n)\tau \) from \( M_{n-1}(I_i) \) after time \( t' \) and that he is at distance at most \( \rho' + (1 + \varepsilon_n)\tau \) from \( g \) at time \( (i+1)\Delta_{n-1} \). If we are in case 1, the statement follows from Claim 12, so assume case 2.
By Claim 10, the man will not do an avoidance move again after time \( t' \) until he reaches \( g \) or \( I_i \) ends. While he is running directly towards \( g \), his distance to the segment \( M_{n-1}(I_i) \) is decreasing, so it follows that the distance is always at most \( \rho' + 2(1 + \varepsilon_n)\tau \), as claimed. Since he was doing avoidance moves in a period of length at most \( \tau \) before the escape move at time \( t' \), he can completely compensate for the delay caused by the avoidance moves in the same amount of time by running directly towards \( g \). The total delay is therefore at most \( 2\tau \).

Since he would reach \( g \) at time \((i + 1)\Delta_n - 2\tau\) at the latest if he did not do any avoidance moves, it follows that he reaches \( g \) at time \((i + 1)\Delta_n - 2\tau\) or earlier. The statement then follows from Claim 12.

We are now ready to finish our proof of Theorem 8. Claim 13 implies that during interval \( I_i \) for any \( i \), the distance from the man to any of the lions \( l_1, \ldots, l_{n-1} \) is at most

\[
\rho' + 2(1 + \varepsilon_n)\tau < \rho + 2r + 2(1 + \varepsilon_n)\tau \leq c_n - 1/2.
\]

Thus, these lions will not catch the man. By Claim 9, the distance to \( l_n \) is always at least \( r - (3 + \varepsilon_n)\Delta_n \). Therefore, we now define \( c_n := \min\{c_{n-1}/2, r - (3 + \varepsilon_n)\Delta_n\} \), and it holds that in the time interval \( I := [i\Delta_n, (i + 1)\Delta_n] \), the distance from any point on the segment \( m(I) \) to any lion is at least \( c_n \).

Claim 13 implies that at any time \( t \) and for any \( i \in \{2, \ldots, n\} \), the distance \( \|M_{i-1}(t)M_i(t)\| \) is bounded by some constant. It follows that \( \|M_1(t)M_n(t)\| \leq d_n \) for some constant \( d_n > 0 \). Since \( M_1(t) \) traverses a ray with constant speed \( 1 + \varepsilon/2 > 1 \), the man eventually escapes the convex hull of the lions and that the distance diverges to \( \infty \) as \( t \to \infty \). This proves the theorem.

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**References**

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