Relational algebra by way of adjunctions
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Relational Algebra by Way of Adjunctions

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1. Summary

- bulk types (sets, bags, lists) are monads
- monads have nice mathematical foundations via adjunctions
- monads support comprehensions
- comprehension syntax provides a query notation

\[
\begin{array}{l}
\left[ (\text{customer.name, invoice.amount}) \\
| \text{customer} \leftarrow \text{customers}, \\
\text{invoice} \leftarrow \text{invoices}, \\
\text{customer.cid} = \text{invoice.customer}, \\
\text{invoice.due} \leq \text{today} \right]
\end{array}
\]

- monad structure explains selection, projection
- less obvious how to explain join
2. Galois connections

Relating monotonic functions between two ordered sets:

\[(A, \leq) \perp (B, \sqsubseteq)\]

means \(f b \leq a \iff b \sqsubseteq g a\)

For example,

\[(\mathbb{R}, \leq_{\mathbb{R}}) \perp (\mathbb{Z}, \leq_{\mathbb{Z}})\]

\[(\mathbb{Z}, \leq) \perp (\mathbb{Z}, \leq)\]

“Change of coordinates” can sometimes simplify reasoning; eg rhs gives
\(n \times k \leq m \iff n \leq m \div k\), and multiplication is easier to reason about than rounding division.
3. Category theory from ordered sets

A category $\mathbf{C}$ consists of

- a set* $|\mathbf{C}|$ of objects,
- a set* $\mathbf{C}(X, Y)$ of arrows $X \to Y$ for each $X, Y : |\mathbf{C}|$,
- identity arrows $\text{id}_X : X \to X$ for each $X$
- composition $f \cdot g : X \to Z$ of compatible arrows $g : X \to Y$ and $f : Y \to Z$,
- such that composition is associative, with identities as units.

Think of a directed graph, with vertices as objects and paths as arrows.

An ordered set $(A, \leq)$ is a degenerate category, with objects $A$ and a unique arrow $a \to b$ iff $a \leq b$.

\[ \ldots \rightarrow -2 \rightarrow -1 \rightarrow 0 \rightarrow 1 \rightarrow 2 \rightarrow \ldots \]

Many categorical concepts are generalisations from ordered sets.

*proviso...
4. Concrete categories

Ordered sets are a *concrete category*: roughly,

- the objects are *sets with additional structure*
- the arrows are *structure-preserving mappings*

Many useful categories are of this form.

For example, the category \( \textbf{CMon} \) has commutative monoids \((M, \otimes, \epsilon)\) as objects, and homomorphisms \( h: (M, \otimes, \epsilon) \rightarrow (M', \oplus, \epsilon') \) as arrows:

\[
\begin{align*}
  h (m \otimes n) &= h m \oplus h n \\
  h \epsilon &= \epsilon'
\end{align*}
\]

Trivially, category \( \textbf{Set} \) has sets as objects, and total functions as arrows.
5. Functors

Categories are themselves structured objects...

A functor $F : C \to D$ is an operation on both objects and arrows, preserving the structure: $F f : F X \to F Y$ when $f : X \to Y$, and

$$F \ id_X = id_{F \ X}$$
$$F (f \cdot g) = F f \cdot F g$$

For example, forgetful functor $U : CMon \to Set$:

$$U (M, \otimes, \epsilon) = M$$
$$U (h : (M, \otimes, \epsilon) \to (M', \oplus, \epsilon')) = h : M \to M'$$

Conversely, Free : Set $\to$ CMon generates the free commutative monoid (ie bags) on a set of elements:

$$\text{Free } A = (\text{Bag } A, \uplus, \emptyset)$$
$$\text{Free } (f : A \to B) = \text{map } f : \text{Bag } A \to \text{Bag } B$$
6. Adjunctions

Adjunctions are the categorical generalisation of Galois connections. Given categories $\mathbf{C}, \mathbf{D}$, and functors $L : \mathbf{D} \to \mathbf{C}$ and $R : \mathbf{C} \to \mathbf{D}$, adjunction

\[
\begin{array}{ccc}
\mathbf{C} & \perp & \mathbf{D} \\
\downarrow^L & & \downarrow^R \\
\mathbf{C}(L X, Y) & \simeq & \mathbf{D}(X, R Y)
\end{array}
\]

A familiar example is given by currying:

\[
\begin{array}{ccc}
\mathbf{Set} & \perp & \mathbf{Set} \\
\downarrow^{(-)P} & & \downarrow^{(-)P} \\
\mathbf{Set}(X \times P, Y) & \simeq & \mathbf{Set}(X, Y^P)
\end{array}
\]

with $\text{curry} : \mathbf{Set}(X \times P, Y) \simeq \mathbf{Set}(X, Y^P) : \text{curry}^\circ$

hence definitions and properties of $\text{apply} = \text{uncurry}\ id_{Y^P} : Y^P \times P \to Y$
7. Products and coproducts

\[ \text{Set} \xleftarrow{\Delta} \text{Set}^2 \xrightarrow{\Delta} \text{Set} \]

\[ \xleftarrow{\times} \\xrightarrow{\perp} \]

with

\[ \text{fork} : \text{Set}^2(\Delta A, (B, C)) \simeq \text{Set}(A, B \times C) \quad : \text{fork}^\circ \]

\[ \text{junc}^\circ : \text{Set}(A + B, C) \simeq \text{Set}^2((A, B), \Delta C) : \text{junc} \]

hence

\[ \text{dup} = \text{fork } \text{id}_{A,A} : \text{Set}(A, A \times A) \]

\[ (\text{fst}, \text{snd}) = \text{fork}^\circ \text{id}_{B \times C} : \text{Set}^2(\Delta(B, C), (B, C)) \]

give tupling and projection. Dually for sums and injections, and generally for any arity—even zero.
8. Free commutative monoids

Adjunctions often capture embedding/projection pairs:

\[ \text{CMon} \perp \text{Set} \]

\[ \text{with } [-] : \text{CMon}(\text{Free } A, (M, \otimes, \epsilon)) \]
\[ \cong \text{Set}(A, U (M, \otimes, \epsilon)) : [-] \]

Unit and counit:

\[ \text{single } A = [id_{\text{Free } A}] : A \rightarrow U (\text{Free } A) \]
\[ \text{reduce } M = [id_M] : \text{Free } (U M) \rightarrow M \quad \text{-- for } M = (M, \otimes, \epsilon) \]

whence, for \( h : \text{Free } A \rightarrow M \) and \( f : A \rightarrow U M = M \),

\[ h = \text{reduce } M \cdot \text{Free } f \iff U h \cdot \text{single } A = f \]

ie 1-to-1 correspondence between homomorphisms from the free commutative monoid (bags) and their behaviour on singletons.
## 9. Aggregation

Aggregations are bag homomorphisms:

<table>
<thead>
<tr>
<th>aggregation</th>
<th>monoid</th>
<th>action on singletons</th>
</tr>
</thead>
<tbody>
<tr>
<td>count</td>
<td>$\langle \mathbb{N}, 0, + \rangle$</td>
<td>$\lhd a \rhd \mapsto 1$</td>
</tr>
<tr>
<td>sum</td>
<td>$\langle \mathbb{R}, 0, + \rangle$</td>
<td>$\lhd a \rhd \mapsto a$</td>
</tr>
<tr>
<td>max</td>
<td>$\langle \mathbb{Z}, \text{minBound}, \text{max} \rangle$</td>
<td>$\lhd a \rhd \mapsto a$</td>
</tr>
<tr>
<td>min</td>
<td>$\langle \mathbb{Z}, \text{maxBound}, \text{min} \rangle$</td>
<td>$\lhd a \rhd \mapsto a$</td>
</tr>
<tr>
<td>all</td>
<td>$\langle \mathbb{B}, \text{True}, \wedge \rangle$</td>
<td>$\lhd a \rhd \mapsto a$</td>
</tr>
<tr>
<td>any</td>
<td>$\langle \mathbb{B}, \text{False}, \lor \rangle$</td>
<td>$\lhd a \rhd \mapsto a$</td>
</tr>
</tbody>
</table>

Selection is a homomorphism, to bags, using action

$\text{guard} : (A \rightarrow \mathbb{B}) \rightarrow \text{Bag } A \rightarrow \text{Bag } A$

$\text{guard } p \ a = \text{if } p \ a \text{ then } \lhd a \rhd \text{ else } \emptyset$

Laws about selections follow from laws of homomorphisms (and of coproducts, since $\mathbb{B} = 1 + 1$).
10. Monads

Bags form a *monad* \((\text{Bag}, \text{union}, \text{single})\) with

\[
\text{Bag} = U \cdot \text{Free} \\
\text{union} : \text{Bag} (\text{Bag} A) \to \text{Bag} A \\
\text{single} : A \to \text{Bag} A
\]

which justifies the use of comprehension notation \(\{ f \ a \ b \mid a \leftarrow x, b \leftarrow g \ a \}^\text{+}\).

In fact, for any adjunction \(L \dashv R\) between \(C\) and \(D\), we get a monad \((T, \mu, \eta)\) on \(D\), where

\[
T = R \cdot L \\
\mu A = R \left[ id_A \right] L : T (T A) \to T A \\
\eta A = [id_A] : A \to T A
\]
11. Maps

Database indexes are essentially maps $\text{Map } K \ V = V^K$. Maps $(-)^K$ from $K$ form a monad (the Reader monad in Haskell), so arise from an adjunction.

The laws of exponents arise from this adjunction, and from those for products and coproducts:

- $\text{Map } 0 \ V \cong 1$
- $\text{Map } 1 \ V \cong V$
- $\text{Map } (K_1 + K_2) \ V \cong \text{Map } K_1 \ V \times \text{Map } K_2 \ V$
- $\text{Map } (K_1 \times K_2) \ V \cong \text{Map } K_1 (\text{Map } K_2 \ V)$
- $\text{Map } K \ 1 \cong 1$
- $\text{Map } K (V_1 \times V_2) \cong \text{Map } K \ V_1 \times \text{Map } K \ V_2 : \text{merge}$
12. Indexing

Relations are in 1-to-1 correspondence with set-valued functions:

\[ \text{Rel} \xrightarrow{\perp} \text{Set} \]

where \( J \) embeds, and \( E R : A \rightarrow \text{Set} B \) for \( R : A \sim B \).

Moreover, the correspondence remains valid for bags:

\[ \text{index} : \text{Bag} (K \times V) \simeq \text{Map} K (\text{Bag} V) \]

Together, \textit{index} and \textit{merge} give efficient relational joins:

\[ x f \bowtie g y = \text{flatten} (\text{Map} K \ cp (\text{merge} (\text{groupBy} f x, \text{groupBy} g y))) \]

\[ \text{groupBy} : (V \rightarrow K) \rightarrow \text{Bag} V \rightarrow \text{Map} K (\text{Bag} V) \]

\[ \text{flatten} : \text{Map} K (\text{Bag} V) \rightarrow \text{Bag} V \]
13. Pointed sets and finite maps

Model *finite maps* $\text{Map}_*$ not as partial functions, but *total* functions to a *pointed* codomain $(A, a)$, i.e. a set $A$ with a distinguished element $a : A$.

Pointed sets and point-preserving functions form a category $\text{Set}_*$. There is an adjunction to $\text{Set}$, via

\[
\begin{array}{ccc}
\text{Set}_* & \xrightarrow{\bot} & \text{Set} \\
\downarrow & & \downarrow \\
\text{Maybe} & \xleftarrow{\bot} & \text{Set} \\
\end{array}
\]

where $\text{Maybe } A \simeq 1 + A$ adds a point, and $U (A, a) = A$ discards it.

In particular, $(\text{Bag } A, \emptyset)$ is a pointed set. Moreover, $\text{Bag } f$ is point-preserving, so we get a functor $\text{Bag}_* : \text{Set} \to \text{Set}_*$.

Indexing remains an isomorphism:

\[
\text{index} : \text{Bag}_* (K \times V) \simeq \text{Map}_* K (\text{Bag}_* V)
\]
14. Graded monads

A catch: finite maps aren’t a monad, because

\[ \eta a = \lambda k \rightarrow a : A \rightarrow \text{Map } K A \]

in general yields an infinite map.

However, finite maps are a graded monad*: for monoid \((M, \otimes, \epsilon)\),

\[
\begin{align*}
\mu X &: T_m (T_n X) \rightarrow T_{m \otimes n} X \\
\eta X &: X \rightarrow T_\epsilon X
\end{align*}
\]

satisfying the usual laws. These too arise from adjunctions*.

We use the monoid \((\mathbb{K}, \times, 1)\) of finite key types under product.
15. Conclusions

- *monad comprehensions* for database queries
- structure arising from *adjunctions*
- equivalences from *universal properties*
- fitting in *relational joins*, via indexing
- to do: calculating *query optimisations*

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