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Published in:
24th Annual European Symposium on Algorithms (ESA 2016)

DOI:
10.4230/LIPIcs.ESA.2016.4

Publication date:
2016

Document Version
Publisher's PDF, also known as Version of record

Citation for published version (APA):
Outer Common Tangents and Nesting of Convex Hulls in Linear Time and Constant Workspace

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Abstract
We describe the first algorithm to compute the outer common tangents of two disjoint simple polygons using linear time and only constant workspace. A tangent of a polygon is a line touching the polygon such that all of the polygon lies on the same side of the line. An outer common tangent of two polygons is a tangent of both polygons such that the polygons lie on the same side of the tangent. Each polygon is given as a read-only array of its corners in cyclic order. The algorithm detects if an outer common tangent does not exist, which is the case if and only if the convex hull of one of the polygons is contained in the convex hull of the other. Otherwise, two corners defining an outer common tangent are returned.

1998 ACM Subject Classification I.3.5 Computational Geometry and Object Modeling

Keywords and phrases simple polygon, common tangent, optimal algorithm, constant workspace

Digital Object Identifier 10.4230/LIPIcs.ESA.2016.4

1 Introduction

A tangent of a polygon is a line touching the polygon such that all of the polygon lies on the same side of the line. An outer common tangent of two polygons is a tangent of both polygons such that the polygons lie on the same side of the tangent. Two disjoint polygons have exactly two outer common tangents unless their convex hulls are nested. If they are properly nested, there is no outer common tangent. In this paper, we study the problem of computing the outer common tangents of two disjoint simple polygons, each given as a read-only array of its corners in cyclic order. We give an algorithm computing the outer common tangents in linear time using only a constant number of variables each storing a boolean value or an index of a corner in the array. We are therefore working in the constant workspace model of computation.

The constant workspace model is a restricted version of the RAM model in which the input is read-only, the output is write-only, and only \(O(\log n)\) additional bits of workspace (with both read and write access) are available, where \(n\) is the size of the input. Clearly, \(\Omega(\log n)\) bits in the workspace are necessary to solve any interesting computational problem, because that many bits are required to store an index of or a pointer to an entry in the input. Since blocks of \(\Theta(\log n)\) bits are considered to form words in the memory, algorithms in the constant workspace model use \(O(1)\) words of memory, which explains the name of the model.

* Research partially supported by Mikkel Thorup’s Advanced Grant from the Danish Council for Independent Research under the Sapere Aude research career programme.
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The practical relevance of studying problems in the constant workspace model is increasing, as there are many current and emerging memory technologies where writing can be much more expensive than reading in terms of time and energy [8].

The constant workspace model was first studied explicitly for geometric problems by Asano et al. [4]. Recently, there has been growing interest in algorithms for geometric problems using constant or restricted workspace, see for instance [1, 3, 5, 6, 9, 11, 14].

The problem of computing common tangents of two polygons has received most attention in the case that the polygons are convex. For instance, computing the outer common tangents of disjoint convex polygons is used as a subroutine in the classical divide-and-conquer algorithm for the convex hull of a set of \(n\) points in the plane due to Preparata and Hong [17]. They give a naive linear-time algorithm for outer common tangents, as it suffices for an \(O(n \log n)\)-time convex hull algorithm. The problem is also considered in various dynamic convex hull algorithms [7, 12, 16]. Overmars and van Leeuwen [16] give an \(O(\log n)\)-time algorithm for computing an outer common tangent of two disjoint convex polygons when a separating line is known, where each polygon has at most \(n\) corners. Kirkpatrick and Snoeyink [13] give an \(O(\log n)\)-time algorithm for the same problem but without using a separating line. Guibas et al. [10] give a lower bound of \(\Omega(\log^2 n)\) on the time required to compute an outer common tangent of two intersecting convex polygons, even if they are known to intersect in at most two points. They also describe an algorithm achieving that bound. Toussaint [18] considers the problem of computing separating common tangents of convex polygons and notes that the problem occurs in problems related to visibility, collision avoidance, range fitting, etc. He gives a linear-time algorithm. Guibas et al. [10] give an \(O(\log n)\)-time algorithm for the same problem. All the above-mentioned algorithms with sublinear running times make essential use of the convexity of the polygons. If the polygons are not convex, a linear-time algorithm can be used to compute the convex hulls before computing the tangents [15]. However, if the polygons are given in read-only memory, \(\Omega(n)\) extra bits are required to store the convex hulls, so this approach does not work in the constant workspace model.

Abrahamsen [2] gives a linear-time constant-workspace algorithm to compute the outer common tangents of two simple polygons the convex hulls of which are disjoint. In this paper, we show that the same is possible as long as the polygons (but not necessarily their convex hulls) are disjoint. The algorithm is only slightly different from the one in [2], but its proof of correctness requires much more effort. In particular, the proof relies on an intricate continuous analysis of the algorithm. Before, it was not even clear whether to expect existence of a linear-time constant-workspace algorithm that does not require the convex hulls to be disjoint, because it happens quite often that a computational problem exhibits different behavior for disjoint polygons and for polygons that are not disjoint. For instance, as it has been mentioned above, the outer common tangents of two disjoint convex polygons can be computed in time \(O(\log n)\), while doing the same for two convex polygons that intersect in two points requires time \(\Omega(\log^2 n)\).

A separating common tangent of two polygons is a tangent of both polygons such that the polygons lie on the opposite sides of the tangent. Two disjoint polygons have exactly two separating common tangents provided that their convex hulls are disjoint. If they intersect properly, there is no separating common tangent. Abrahamsen [2] describes a linear-time constant-workspace algorithm that computes the separating common tangents of two simple polygons. In particular, it detects whether the convex hulls of two simple polygons are disjoint. Our current algorithm can decide whether the convex hulls two simple polygons are nested, which happens when it is unable to find an outer common tangent. To the best of our knowledge, this was not known to be possible in linear time and constant workspace prior to
this work. Our algorithm and the algorithm from [2] together enable us to determine, for
two disjoint simple polygons in general position, the full relation between their convex hulls
(whether they are nested, overlapping, or disjoint) in linear time and constant workspace.

It remains open whether an outer common tangent of two polygons that are not disjoint
can be found in linear time using constant workspace.

2 Terminology and Notation

For any two points $a$ and $b$ in the plane, the closed line segment with endpoints $a$ and $b$
is denoted by $ab$. When $a \neq b$, the straight line containing $a$ and $b$ that is infinite in both
directions is denoted by $L(a, b)$, and the ray starting at $a$ and going through $b$ is denoted by
$R(a, b)$. For three points $a$, $b$, and $c$, consider the line $L(a, b)$ as oriented from $a$ towards $b$,
and define $T(a, b, c)$ to be 1 if $c$ lies to the left of $L(a, b)$, 0 if $a$, $b$, $c$ are collinear, and $-1$ if
$c$ lies to the right of $L(a, b)$. Let LHP$(a, b)$ denote the closed half-plane lying to the left of
$L(a, b)$ and RHP$(a, b)$ denote the closed half-plane lying to the right of $L(a, b)$.

A simple polygon, or just a polygon, with corners $x_0, \ldots, x_{n-1}$ is a closed polygonal curve
in the plane composed of $n$ edges $x_0x_1, \ldots, x_{n-2}x_{n-1}, x_{n-1}x_0$ such that the segments have
no common points other than the common endpoints of pairs of consecutive edges. The
region of the plane bounded by a polygon $P$ (including $P$ itself) is a polygonal region.

Assume for the rest of this paper that $P_0$ and $P_1$ are two disjoint simple polygons with
$n_0$ and $n_1$ corners, respectively. (We allow one of $P_0$ and $P_1$ to be contained in the “interior
region” of the other – in that case our algorithm will report that the convex hulls are
nested and no outer common tangent exists.) Assume that $P_k$ is defined by a read-only
array of its corners $p_k[0], p_k[1], \ldots, p_k[n_k - 1]$ for $k \in \{0, 1\}$. Assume further, without loss
of generality, that the corners of $P_0$ are given in counterclockwise order and the corners
of $P_1$ are given in clockwise order. (The orientation of a polygon can be easily tested in
linear time using constant workspace, and the algorithm can choose to traverse the polygon
forwards or backwards, accordingly.) Finally, assume that the corners are in general position
in the sense that $P_0$ and $P_1$ have no corners in common and the combined set of corners
\{p_0[0], \ldots, p_0[n_0 - 1], p_1[0], \ldots, p_1[n_1 - 1]\} contains no triple of collinear points.

Indices of the corners of $P_k$ are considered modulo $n_k$, so that $p_k[i]$ and $p_k[j]$ denote the
same corner when $i \equiv j \pmod{n_k}$. For $a, b \in P_k$, the chain $P_k[a, b]$ is the portion of $P_k$ from
$a$ to $b$ in the order assigned to $P_k$ (counterclockwise for $P_0$, clockwise for $P_1$). If $i$ and $j$ are
indices of corners on $P_k$, we write $P_k[i, j]$ to denote $P_k[p_k[i], p_k[j]]$.

A tangent of $P_k$ is a line $\ell$ such that $\ell$ and $P_k$ are not disjoint and $P_k$ is contained in one
of the closed half-planes determined by $\ell$. The line $\ell$ is a common tangent of $P_0$ and $P_1$ if it
is a tangent of both $P_0$ and $P_1$. A common tangent is an outer common tangent if $P_0$ and
$P_1$ are on the same side of the tangent, otherwise the common tangent is separating.

For a simple polygon $P$, let $H(P)$ denote the convex hull of $P$. The following lemma
asserts well-known properties of common tangents of polygons. See Figures 1–3.

Lemma 1. A line is a tangent of a polygon $P$ if and only if it is a tangent of $H(P)$.
Under our general position assumptions, the following holds. If one of $H(P_0)$ and $H(P_1)$
is completely contained in the other, there are no outer common tangents of $P_0$ and $P_1$.
Otherwise, there are two or more, and there are exactly two if $P_0$ and $P_1$ are disjoint.
If $H(P_0)$ and $H(P_1)$ are not disjoint, there are no separating common tangents of $P_0$ and $P_1$.
Otherwise, there are exactly two.
Algorithm 1: OuterCommonTangent($P_0, P_1$)

1: $s_0 \leftarrow 0$; $v_0 \leftarrow 0$; $b_0 \leftarrow \text{false}$; $s_1 \leftarrow 0$; $v_1 \leftarrow 0$; $b_1 \leftarrow \text{false}$; $u \leftarrow 0$
2: while $s_0 < 2n_0$ and $s_1 < 2n_1$ and $(v_0 < s_0 + n_0$ or $v_1 < s_1 + n_1)$
3:     $v_u \leftarrow v_u + 1$
4:     if $T(p_0[s_0], p_1[s_1], p_u[v_u]) = 1$
5:         if $p_{1-u}[s_{1-u}] \in \Delta(p_u[s_u], p_u[v_u-1], p_u[v_u])$
6:             $b_u \leftarrow \text{true}$
7:         if not $b_u$
8:             $s_u \leftarrow v_u$; $v_{1-u} \leftarrow s_{1-u}$; $b_{1-u} \leftarrow \text{false}$
9:     $u \leftarrow 1 - u$
10: if $s_0 \geq 2n_0$ or $s_1 \geq 2n_1$ or $b_0$ or $b_1$
11: return nested
12: return $(s_0, s_1)$

3: Algorithm

Let the outer common tangents of $P_0$ and $P_1$ be defined by pairs of corners $(\ell_0, \ell_1)$ and $(r_0, r_1)$ so that $\ell_0, r_0 \in P_0$, $\ell_1, r_1 \in P_1$, and $P_0, P_1 \subset \text{LHP}(\ell_0, \ell_1) \cap \text{RHP}(r_0, r_1)$. Algorithm 1 returns a pair of indices $(s_0, s_1)$ such that $(r_0, r_1) = (p_0[s_0], p_1[s_1])$ or, if the convex hulls of $P_0$ and $P_1$ are nested so that the tangents do not exist, the algorithm reports that by returning nested. Finding $(\ell_0, \ell_1)$ requires running Algorithm 1 with the roles of $P_0$ and $P_1$ interchanged and with the orders of the corners of $P_0$ and $P_1$ reversed – each array reference $p_k[i]$ is translated to $p_{1-k}[i]$ for $k \in \{0, 1\}$, and the returned result is $(s_1, s_0)$ such that $(\ell_0, \ell_1) = (p_0[s_0], p_1[s_1])$.

The algorithm maintains a pair of indices $(s_0, s_1)$ which determines the tangent candidate $L(p_0[s_0], p_1[s_1])$. Starting from $(s_0, s_1) = (0, 0)$ and advancing the indices $s_0$, $s_1$ appropriately, the algorithm attempts to reach a situation that $(p_0[s_0], p_1[s_1]) = (r_0, r_1)$, that is, $P_0, P_1 \subset \text{RHP}(p_0[s_0], p_1[s_1])$. At the start and after each update to $(s_0, s_1)$, the algorithm traverses $P_0$ and $P_1$ in parallel with indices $(v_0, v_1)$, starting from $(v_0, v_1) = (s_0, s_1)$ and advancing $v_0$ and $v_1$ alternately. The variable $u \in \{0, 1\}$ determines the polygon $P_u$ in which we advance the traversal in a given iteration. If the test in line 4 happens to be positive, then the corner $p_u[v_u]$ lies on the “wrong side” of the tangent candidate, witnessing $P_u \not\subset \text{RHP}(p_0[s_0], p_1[s_1])$. 

\begin{figure}[h]
\centering
\includegraphics[width=\textwidth]{figure1.png}
\caption{The convex hulls are disjoint – separating and outer common tangents exist.}
\end{figure}

\begin{figure}[h]
\centering
\includegraphics[width=\textwidth]{figure2.png}
\caption{The convex hulls overlap – only outer common tangents exist.}
\end{figure}

\begin{figure}[h]
\centering
\includegraphics[width=\textwidth]{figure3.png}
\caption{The convex hulls are nested – no common tangents exist.}
\end{figure}
Algorithm 1 ends with which was on the correct side of the tangent candidate before the update to $s_0$. In particular, Lemma 2 would no longer be true. On the other hand, if making updates to $s_0$ in line 8 regardless of the current values of $b_0$ and $b_1$, the algorithm could never end making updates to $s_0$ and $s_1$ even if the outer common tangents exist (see [2] for an example of such a behavior). In particular, Lemma 2 would no longer be true. On the other hand, if the convex hulls of $P_0$ and $P_1$ are nested, then one of the following happens:

\[ P_0, P_1 \subset \text{RHP}(p_0[s_0], p_1[s_1]). \]
the algorithm never ends making updates to $s_0$ and $s_1$,
- one of $b_0$, $b_1$, say $b_k$, is true and the algorithm has traversed $P_{1-k}$ entirely with the index $v_{1-k}$ after last update to $s_{1-k}$ without detecting any corner on the wrong side of the tangent candidate.

In both cases, taking the conditions $s_0 < 2n_0$ and $s_1 < 2n_1$ in line 2 back into account, the algorithm reports that the convex hulls of $P_0$ and $P_1$ are nested in line 11.

\textbf{Lemma 3.} If the outer common tangents of $P_0$ and $P_1$ exist, then the loop in line 2 of Algorithm 1 ends with $b_0 = b_1 = \text{false}$.

\textbf{Proof.} We prove a slightly stronger statement, namely, that at most one of $b_0$ and $b_1$ can be true at a time, and if one of $b_0$ and $b_1$ is true, then it will be cleared subsequently. Hence, the algorithm cannot terminate with $b_0 = \text{true}$ or $b_1 = \text{true}$.

Consider an iteration $i$ of the loop in line 2 which leads to changing the value of $b_0$ from false to true in line 6. By induction, we can assume that $b_1 = \text{false}$. Since the test in line 5 is positive, the edge $P_0[v_0 - 1, v_0]$ intersects $L(p_0[s_0], p_1[s_1])$ at a point $x$ such that $p_1[s_1]$ lies on the segment $p_0[s_0]x$. Moreover, $P_0[p_0[s_0], x] \subset RHP(p_0[s_0], p_1[s_1])$, otherwise $b_0$ would be set before. Let $y$ be the first corner of $P_1$ after $p_1[s_1]$ such that $y \notin RHP(p_0[s_0], p_1[s_1])$. Such a corner exists, otherwise $P_1$ would be contained in the convex hull of $P_0$. It follows that the test in line 4 will be positive in the first iteration $j$ after $i$ in which $u = 1$ and $p_1[v_1] = y$. The edge $P_1[v_1 - 1, v_1]$ intersects $L(p_0[s_0], p_1[s_1])$ at a point on the segment $p_0[s_0]x$, and hence the test in line 5 is negative in iteration $j$. Therefore, $b_0$ is cleared and we again have $b_0 = b_1 = \text{false}$. The same argument shows that $b_1$ will be cleared after being set.

\textbf{Theorem 4.} Algorithm 1 is correct, runs in linear time, and uses constant workspace. Specifically, if the outer common tangents exist, then Algorithm 1 returns a pair of indices $(s_0, s_1)$ such that $(r_0, r_1) = (p_0[s_0], p_1[s_1])$, that is, $P_0, P_1 \subset RHP(p_0[s_0], p_1[s_1])$. Otherwise, the algorithm returns nested.

\textbf{Proof.} First, suppose the algorithm returns $(s_0, s_1)$ in line 12. Consider the final values of $s_0$, $s_1$, $b_0$ and $b_1$. Due to the test in line 10, we have $s_0 < 2n_0$, $s_1 < 2n_1$, and $b_0 = b_1 = \text{false}$, so the loop in line 2 has ended because $v_0 \geq s_0 + n_0$ and $v_1 \geq s_1 + n_1$. After the last update to $(s_0, s_1)$, the test in line 4 has been performed for every $v_0 \in \{s_0 + 1, \ldots, s_0 + n_0\}$ and every $v_1 \in \{s_1 + 1, \ldots, s_1 + n_1\}$ and was negative - otherwise a further update would have been performed in line 8, as $b_0 = b_1 = \text{false}$. This shows that $P_0, P_1 \subset RHP(p_0[s_0], p_1[s_1])$.

Now, suppose that the outer common tangents exist. By Lemma 2 and Lemma 3, the loop in line 2 ends with $s_0 < 2n_0$, $s_1 < 2n_1$, and $b_0 = b_1 = \text{false}$. Hence $(s_0, s_1)$ is returned in line 12. In view of the discussion above, this proves correctness of the algorithm.

It is clear that Algorithm 1 uses constant workspace. For the running time, note that if an update to $(s_0, s_1)$ happens in iteration $i$, the sum $s_0 + s_1$ is increased by at least $\frac{n_0 - n_1}{2}$, where $j$ is the number of the previous iteration in which an update to $(s_0, s_1)$ happened or $j = 0$ if there has been no update before. By induction, we see that there have been at most $2(s_0 + s_1)$ iterations until an update to $(s_0, s_1)$. Suppose first that $s_0 < 2n_0$ and $s_1 < 2n_1$ when the loop in line 2 terminates. There have been at most $4(n_0 + n_1)$ iterations until the final update to $(s_0, s_1)$. Thereafter, at most $2 \max\{n_0, n_1\} \leq 2(n_0 + n_1)$ iterations follow until $v_0 \geq s_0 + n_0$ and $v_1 \geq s_1 + n_1$, when the loop in line 2 terminates. Hence, there are at most $6(n_0 + n_1)$ iterations in total. Now, suppose that $s_0 \geq 2n_0$ or $s_1 \geq 2n_1$ when the loop terminates. By the same argument, the second to last update to $(s_0, s_1)$ happens after at most $4(n_0 + n_1)$ iterations, after which at most $2(n_0 + n_1)$ iterations follow until the last update to $(s_0, s_1)$. The loop is terminated immediately after the last update. Hence, we get
the same bound of $6(n_0 + n_1)$ iterations. Clearly, each iteration takes constant time, so the total running time of the algorithm is linear.

4 Proof of Lemma 2

For our analysis, it will be convenient to imagine the execution of Algorithm 1 in continuous time. By considering various discrete events happening during the continuous execution of the algorithm, we are able to prove the invariant stated in Lemma 2.

4.1 Additional Terminology and Notation

For $U \subseteq \mathbb{R}^2$, let $\mathcal{F}(U)$ denote the set of compact subsets of $U$. By an interval, we mean a bounded interval of real numbers. We allow an interval to be closed or open at each endpoint independently. We shall consider functions defined on an interval $I$ with the following sets (or their subsets) as codomains: $\mathbb{R}$ with the standard metric, $\mathbb{R}^2$ with the Euclidean metric, and $\mathcal{F}(\mathbb{R}^2)$ with the Hausdorff metric, a set $\mathcal{S}$ of functions with the discrete metric, and the power set $2^\mathcal{S}$ of a set $\mathcal{S}$ of functions, again with the discrete metric. The only purpose of these metrics is to have a suitable notion of convergence. We think of the domain $I$ as time. If $f$ is a function with domain $I$ and $I'$ is a subinterval of $I$, then $f \upharpoonright I'$ denotes the restriction of $f$ to $I'$. For a function $f : I \to X$, where $X$ is a (sub)set of one of the codomains above, a point in time $t \in I$ is a discontinuity of $f$ if $f$ is not continuous at $t$. We write $f(\nearrow t^*)$ to denote the limit of $f(t)$ as $t \to t^*$ from below, where $t^* \in I \setminus \{\inf I\}$, $f(\searrow t^*)$ to denote the limit of $f(t)$ as $t \to t^*$ from above, where $t^* \in I \setminus \{\sup I\}$. If the limits $f(\nearrow t^*)$ exist for all $t^* \in I \setminus \{\inf I\}$ and the limits $f(\searrow t^*)$ exist for all $t^* \in I \setminus \{\sup I\}$, then we say that $f$ has one-sided limits. Each of the functions $f$ that we consider has one-sided limits and finitely many discontinuities. Note that $f$ has a discontinuity at a point in time $t \in I$ if and only if $f(\nearrow t) \neq f(t)$ or $f(\searrow t) \neq f(t)$. A function $f : I \to \mathcal{F}(U)$, where $U \subseteq \mathbb{R}^2$, is monotonically decreasing if $f(t) \supseteq f(t')$ for any $t, t' \in I$ such that $t < t'$.

Lemma 5. Let $I$ be an interval and $f : I \to \mathcal{F}(U)$ be a function with one-sided limits and finitely many discontinuities, where $U \subseteq \mathbb{R}^2$. Suppose $f \upharpoonright I'$ is monotonically decreasing for every subinterval $I' \subseteq I$ such that $f \upharpoonright I'$ is continuous on $I'$. Furthermore, suppose that $f(\nearrow t) \supseteq f(t)$ for any $t \in I \setminus \{\inf I\}$ such that $f(\nearrow t) \neq f(t)$, $f(t) \supseteq f(\searrow t)$ for any $t \in I \setminus \{\sup I\}$ such that $f(t) \neq f(\searrow t)$. Then $f$ is monotonically decreasing in the entire domain $I$.

Proof. Let $t_1 < \cdots < t_n$ be the discontinuities of $f$. Let $t, t' \in I$ and $t < t'$. If there is no $i$ with $t \leq t_i \leq t'$, then $f \upharpoonright [t, t']$ is continuous, so it follows from the assumption that $f(t) \supseteq f(t')$. Otherwise, let $i$ be minimum and $j$ be maximum such that $t \leq t_i \leq t_j \leq t'$. If $t < t_i$, then the assumptions yield $f(t) \supseteq f(\nearrow t_i) \supseteq f(t_i)$, $f(t) \supseteq f(\searrow t_{j-1}) \supseteq f(t_{j-1})$ for $k \in \{i, \ldots, j-1\}$, and $f(t_j) \supseteq f(\searrow t_j) \supseteq f(t')$ if $t_j < t'$. Thus $f(t) \supseteq f(t')$.

4.2 Continuous Interpretation of the Algorithm

Let $m$ denote the number of iterations of the loop in line 2 performed by Algorithm 1. For $i \in \{0, \ldots, m\}$ and $k \in \{0, 1\}$, let $v_k(i)$ and $s_k(i)$ denote the values of $v_k$ and $s_k$, respectively, after $i$ iterations of the loop. In particular, $v_k(0) = s_k(0) = 0$. For $x \in \mathbb{R} \setminus \mathbb{Z}$, let $p_k[x]$ denote the interpolated point $([x] - x)p_k([x]) + (x - [x])p_k([x])$ on the edge $P_k[x, [x]]$. 

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We extend the functions $s_0$ and $s_1$ to the real interval $[0,m]$ as follows. We imagine that the $i$th iteration of the loop in line 2 starts at time $i-1$ and ends at time $i$, and during that iteration $v_u$ grows continuously from $v_u(i-1) = v_u(i-1)$ to $v_u(i)$. Thus we define $v_u(t) = v_u(i-1) + t$ for $t \in (i-1, i)$. Suppose that the update in line 8 is to be performed in the $i$th iteration. If $s_u(i-1) = v_u(i-1)$, then all of the edge $P_u[v_u(i-1), v_u(i)]$ is in LHP($p_0[s_0(i-1)], p_1[s_1(i-1)]$). We therefore imagine that the update happens at time $i-1$ and then $s_u(t) = v_u(t)$ and $v_{1-u}(t) = s_{1-u}(i-1)$ for $t \in (i-1, i)$. If $s_u(i-1) < v_u(i-1)$, then the edge $P_u[v_u(i-1), v_u(i)]$ intersects the tangent candidate at a point $p_u[v_u(t^*)]$, where $t^* \in (i-1, i)$. We therefore imagine that the update in line 8 happens at time $t^*$ and then $s_u(t) = v_u(i-1)$ and $v_{1-u}(t) = s_{1-u}(i-1)$ for $t \in (t^*, i)$, and we say that $s_u$ jumps from $s_u(t)$ to $v_u(t^*)$. Finally, in either case, we define $s_{1-u}(t) = s_{1-u}(i-1)$ for $t \in (i-1, i)$. The functions $s_0, s_1 : [0,m] \to \mathbb{R}$ thus defined are nondecreasing, have one-sided limits and finitely many discontinuities, and are left-continuous, that is, $s_0(\succ t) = s_0(t)$ and $s_1(\succ t) = s_1(t)$ for every $t \in (0,m]$. We have also defined functions $v_0, v_1 : [0,m] \to \mathbb{R}$, but we are not going to use them any more.

Observation 6. At any point in time during the execution of the continuous version of Algorithm 1, at most one of $s_0, s_1$ is changing. The tangent candidate $L(p_0[s_0], p_1[s_1])$ either is not moving, or is turning continuously counterclockwise around $p_0[s_0]$ (when $s_1$ is changing), or is turning continuously clockwise around $p_1[s_1]$ (when $s_0$ is changing).

The following is trivial if $s_k(t) = s_k(\prec t)$ and otherwise is a direct consequence of the test in line 5 and of the fact that the update in line 8 is only performed when $b_u = \text{false}$.

Observation 7. If $t \in (0,m]$ and $k \in \{0,1\}$, then $p_k[s_k(\prec t)] \in \mathcal{R}(p_1-k[s_1-k(t)], p_k[s_k(t)])$ and $P_k[s_k(t), s_k(\prec t)] \subset \text{RHP}(p_0[s_0(t)], p_1[s_1(t)])$.

4.3 Auxiliary Structure on the Polygons

In this subsection, we introduce some auxiliary concepts used in the proof of Lemma 2. They are defined in terms of the polygons $P_0, P_1$ only and are independent of the algorithm.

Assume for this entire subsection that the convex hulls of $P_0$ and $P_1$ are not nested. Thus there are two outer common tangents - let them be given by points $\ell_0, r_0 \in P_0$ and $\ell_1, r_1 \in P_1$ such that $P_0, P_1 \subset \text{LHP}(\ell_0, \ell_1) \cap \text{RHP}(r_0, r_1)$. Let $L = \ell_0\ell_1$ and $R = r_0r_1$. Let $E$ be the polygonal region bounded by the chains $P_0[\ell_0, r_0], P_1[\ell_1, r_1]$ and by the segments $L, R$. Since $P_0$ is oriented counterclockwise and $P_1$ clockwise, the interiors of $P_0$ and $P_1$ lie outside $E$.

Lemma 8. Every segment $xy$ such that $xy \cap P_0 = \{x\}$ and $xy \cap P_1 = \{y\}$ is contained in $E$.

Proof. The set $E \setminus (P_0[\ell_0, r_0] \cup P_1[\ell_1, r_1])$ separates $P_0$ and $P_1$ in $\text{LHP}(\ell_0, \ell_1) \cap \text{RHP}(r_0, r_1)$, so it contains a point $z$ in common with the segment $xy$. If $z \in L$ or $z \in R$, then $xy = \ell_0\ell_1$ or $xy = r_0r_1$, respectively, so $xy$ lies in $E$. So suppose $z$ is in the interior of $E$. The segment $xz$ cannot cross the boundary of $E$ at any point other than $x$, and $zy$ at any point other than $y$. This shows that $xy$ lies in $E$. □

See Figure 5. Let $q_0 \in P_0$ and $q_1 \in P_1$ be fixed points such that at least one of $q_0, q_1$ is a corner of the respective polygon $P_0$ or $P_1$. Let $S = \{q_0, q_1\}$. We consider the segment $S$ as oriented from $q_0$ to $q_1$, so that we can speak of the left side of $S$, LHP($q_0, q_1$), and the
right side of $S$. RHP($q_0, q_1$). A door is a subsegment $xy$ of $S$ such that $xy \cap P_k = \{x\}$ and $xy \cap P_{1-k} = \{y\}$ for some $k \in \{0, 1\}$. By Lemma 8, every door is contained in $E$. A fence is a subsegment $xy$ of $S$ such that $xy \subset E$, $xy \cap P_k = \{x, y\}$, and $xy \cap P_{1-k} = \emptyset$ for some $k \in \{0, 1\}$. Exceptionally, when $S$ contains an edge $xy$ of $P_k$, we call the whole edge $xy$ a fence. Since at least one of $q_0, q_1$ is a corner, the latter is possible only when $x = q_k$ or $y = q_k$. Let $D$ be the set of all doors defined by the fixed points $q_0$ and $q_1$. Figure 5 also illustrates the following lemma.

**Lemma 9.** The doors in $D$ can be ordered as $D_1, \ldots, D_d$ so that if $D_i \cap P_0 = \{x_i\}$ and $D_i \cap P_1 = \{y_i\}$ for $i \in \{1, \ldots, d\}$, then

- the order of points along $P_0[\ell_0, r_0]$ is $\ell_0, x_1, \ldots, x_d, r_0$ (with possible coincidences),
- the order of points along $P_1[\ell_1, r_1]$ is $\ell_1, y_1, \ldots, y_d, r_1$ (with possible coincidences).

The doors partition $E$ into polygonal regions $E_0, \ldots, E_d$ such that

- $E_0$ is bounded by $L, P_0[\ell_0, x_1], D_1$ and $P_1[\ell_1, y_1]$ (it is degenerate when $D_1 = L$),
- $E_i$ is bounded by $D_i, P_0[x_i, x_{i+1}], D_{i+1}$ and $P_1[y_i, y_{i+1}]$, for $i \in \{1, \ldots, d-1\}$,
- $E_d$ is bounded by $D_d, P_0[x_d, r_0], R$ and $P_1[y_d, r_1]$ (it is degenerate when $D_d = R$).

**Proof.** Suppose there are doors $xy, x'y' \in D$ such that $x$ is strictly before $x'$ on $P_0[\ell_0, r_0]$ while $y'$ is strictly before $y$ on $P_1[\ell_1, r_1]$. It follows that the clockwise order of the four points along the boundary of $E$ is $x, x', y, y'$ and no two of these points coincide. By Lemma 8, both $xy$ and $x'y'$ lie in $E$, so they must cross at a point different from their endpoints, which is a contradiction. This shows that the order of endpoints of the doors along $P_0[\ell_0, r_0]$ agrees with that along $P_1[\ell_1, r_1]$, which proves the first statement. The second statement is a straightforward corollary to the first.

From now on, we use $D_1, \ldots, D_d$ to denote the doors in their order according to Lemma 9, and we use $E_0, \ldots, E_d$ to denote the regions defined in Lemma 9.
Recall that we consider $S$ as a segment oriented from $q_0$ to $q_1$. Every door inherits that orientation, so that we can speak of the left side and the right side of the door. Taking into account that the regions $E_{i-1}$ and $E_i$ lie on opposite sides of $D_i$, we classify each door $D_i$ as

- a right-door if $E_{i-1}$ lies to the right and $E_i$ lies to the left of $D_i$ (in particular, if $D_i = L$),
- a left-door if $E_{i-1}$ lies to the left and $E_i$ lies to the right of $D_i$ (in particular, if $D_i = R$).

Lemma 10. Consider a chain $P_k[a, b]$, where $k \in \{0, 1\}$. If $P_k[a, b] \cap S = \{a, b\}$ and $P_k[a, b] \subset \text{RHP}(q_0, q_1)$, then all doors contained in the segment $ab$ occur in pairs of a left-door followed by a right-door, consecutive in the order on $D$.

Proof. Consider the polygonal region $F$ bounded by the chain $P_k[a, b]$ and by the segment $ab$. It follows that $F \subset \text{RHP}(q_0, q_1)$. Each of the regions $E_0, \ldots, E_d$ lies either inside or outside $F$, where $E_0$ and $E_d$ lie outside $F$. Each region $E_i$ lying inside $F$ connects the door $D_i$, which is therefore a left-door, and the door $D_{i+1}$, which is therefore a right-door.

So far we were considering $q_0$ and $q_1$ as fixed points. Now, we allow them to change in time. Specifically, let $I$ be a real interval that can be open or closed at each endpoint independently, and consider $q_0$ and $q_1$ as continuous functions $q_0 : I \rightarrow P_0$ and $q_1 : I \rightarrow P_1$. This way $S$ becomes a continuous function $S : I \rightarrow \mathcal{F}(\mathbb{R}^2)$. Furthermore, suppose at least one of $q_0(t), q_1(t)$ is a corner of the respective polygon for every $t \in I$, so that $S(t)$ can contain at most one other corner (by the general position assumption). Let $X(t)$ denote the set of intersection points of $S(t)$ with $P_0 \cup P_1$. In the exceptional case that $S(t)$ contains an edge of $P_0$ or $P_1$, we only include the endpoints of the edge in $X(t)$. The points in $X(t)$ are changing continuously except that an intersection point appears or disappears at a point in time $t \in I$ when $S(t)$ sweeps over a corner whose both incident edges lie on the same side of $S(t)$. Note that since the corners of $P_0$ and $P_1$ are assumed to be in general position and one of $q_0$ and $q_1$ is a corner, at most one point can appear in or disappear from $X(t)$ at any point in time. The doors are changing continuously except when one of the following door events happens as a point appears in or disappears from $X(t)$:

1. a fence splits into two doors,
2. two doors merge into a fence,
3. a door splits into a smaller door and a fence,
4. a door and a fence merge into a larger door.

Specifically, every door $D$ can be represented as a continuous function $D : I_D \rightarrow \mathcal{F}(\mathbb{R}^2)$, where $I_D$ is a subinterval of $I$ (open or closed at each endpoint independently) such that

1. if $t = \inf I_D \in I_D$, then an endpoint of $D(t)$ is in $X(t)$ but not in $X(\leftarrow t)$,
2. if $t = \sup I_D \in I_D$, then an endpoint of $D(t)$ is in $X(t)$ but not in $X(\rightarrow t)$,
3. if $t = \sup I_D \notin I_D$, then an interior point of $D(\leftarrow t)$ is in $X(t)$ but not in $X(\rightarrow t)$,
4. if $t = \inf I_D \notin I_D$, then an interior point of $D(\rightarrow t)$ is in $X(t)$ but not in $X(\leftarrow t)$.

At any point in time $t \in I$, the set of doors $D(t)$ consists of the doors $D$ such that $t \in I_D$ ordered according to Lemma 9. The following observation, a straightforward consequence of Lemma 9, summarizes how $D(t)$ and the order on $D(t)$ are changing in time.

Observation 11. The set $D(t)$ and the order on $D(t)$ are constant in time intervals where no door event happens. A door event at time $t$ makes the following change to $D(t)$:

1. if a fence splits into two doors $D$ and $D'$, then $D$ and $D'$ are added to $D(\leftarrow t)$ as consecutive doors to form $D(t)$,
2. if two doors $D$ and $D'$ merge into a fence, then $D$ and $D'$ are consecutive in $D(t)$ and they are removed from $D(t)$ to form $D(\rightarrow t)$,
3. if a door $D$ splits into a smaller door $D'$ and a fence, then $D$ is replaced by $D'$ in $\mathcal{D}(\setminus t)$ to form $\mathcal{D}(t)$,
4. if a door $D$ and a fence merge into a larger door $D'$, then $D$ is replaced by $D'$ in $\mathcal{D}(t)$ to form $\mathcal{D}(\setminus t)$.

In case of door events 1 and 2, the two doors $D$ and $D'$ are, in their order in $\mathcal{D}(t)$,
- a right-door followed by a left-door if the edges incident to $w$ lie to the right of $S(t)$,
- a left-door followed by a right-door if the edges incident to $w$ lie to the left of $S(t)$,

where $w$ denotes the corner that triggers the event (i.e., the corner that appears in or disappears from $X(t)$ at time $t$). In case of door events 3 and 4, the door $D'$ keeps the left/right-door status of $D$. The left/right-door status of every door $D$ remains constant over the entire time interval $I_D$.

Now, consider $q_0$ and $q_1$ again as fixed points. Recall that $D_1, \ldots, D_d$ denote the doors in their order according to Lemma 9. We define the weight $W(D_i)$ of every door $D_i$ by induction, as follows:

$$W(D_1) = \begin{cases} 1 & \text{if } D_1 \text{ is a right-door,} \\ -1 & \text{if } D_1 \text{ is a left-door,} \end{cases}$$

$$W(D_i) = \begin{cases} W(D_{i-1}) + 1 & \text{if } D_i \text{ is a right-door,} \\ W(D_{i-1}) - 1 & \text{if } D_i \text{ is a left-door,} \end{cases}$$

for $i \in \{2, \ldots, d\}$. See Figure 5. The following is a direct consequence of Observation 11.

**Observation 12.** When $q_0: I \to P_0$, $q_1: I \to P_1$ are continuous functions, every door $D: I_D \to \mathcal{F}(\mathbb{R}^2)$ maintains constant weight over the entire time interval $I_D$. Furthermore, the function $W^*: I \to \mathbb{Z}$ defined so that $W^*(t)$ is the weight of the last door in the order on $\mathcal{D}(t)$ is constant over the entire time interval $I$.

**Lemma 13.** For any fixed points $q_0$, $q_1$, there is at least one door with weight 1.

**Proof.** The statement is obvious if $q_0 = \ell_0$ and $q_1 = \ell_1$, because in that case there is just one door $L$, which is a right-door by definition, so it has weight 1. To prove the lemma in general, let $I = [0, 1]$ and (abusing notation) consider arbitrary continuous functions $q_0: I \to P_0$ and $q_1: I \to P_1$ such that $q_0(0) = \ell_0$, $q_1(0) = \ell_1$, and $q_0(1)$, $q_1(1)$ are the points $q_0$, $q_1$ fixed in the statement of the lemma. By Observation 12, the function $W^*: I \to \mathbb{Z}$ is constant over $I$, so $W^*(1) = W^*(0) = 1$ as observed above. This shows that the last door in the order on $\mathcal{D}(1)$ has weight 1.

For any fixed points $q_0$, $q_1$, let the primary door $D'$ be the first door with weight 1 in the order on $\mathcal{D}$. Such a door always exists due to Lemma 13.

**Observation 14.** The primary door $D'$ is a right-door and is not preceded by a left-door in the order on $\mathcal{D}$.

Let $y_0$ and $y_1$ denote the endpoints of $D'$ so that $y_0 \in P_0$ and $y_1 \in P_1$. Let $Y_0 = P_0[y_0, r_0]$ and $Y_1 = P_1[y_1, r_1]$. Finally, let the primary region $E'$ be defined as the polygonal region determined by $D'$, $Y_0$, $R$ and $Y_1$. See Figure 5.

**Observation 15.** If $D' = D_i$, then $E'$ is the union of $E_{i+1}, \ldots, E_d$. In particular, $E'$ contains the doors $D_{i+1}, \ldots, D_d$. By Observation 14, the region $E'$ meets $D'$ from the left.
4.4 Back to the Algorithm

We recall the functions $s_0, s_1 : [0, m] \to \mathbb{R}$ describing the execution of Algorithm 1 as explained in Section 4.2, and we define functions $q_0 : [0, m] \to P_0$ and $q_1 : [0, m] \to P_1$ as follows:

$$q_0(t) = p_0[s_0(t)], \quad q_1(t) = p_1[s_1(t)] \quad \text{for } t \in [0, m].$$

They have the property that at least one of $q_0(t)$, $q_1(t)$ is a corner at any point in time $t \in [0, m]$. Some other objects that have been defined in Section 4.3 based on fixed points $q_0$, $q_1$ now become functions of time $t \in [0, m]$: the segment $S$, the primary door $D'$, the points $y_0$, $y_1$, the chains $Y_0$, $Y_1$, and the primary region $E'$.

The functions $q_0$ and $q_1$ have finitely many discontinuities – the points of time $t \in [0, m)$ when the respective $s_k$ jumps from $s_k(t)$ to $s_k(\wedge t)$. It is also clear that they have one-sided limits, since the functions $s_0$ and $s_1$ are bounded and piecewise monotone. It follows that the functions $D' : [0, m] \to F(\mathbb{R}^2)$, $y_0 : [0, m] \to P_0$, $y_1 : [0, m] \to P_1$, $Y_0 : [0, m] \to F(P_0)$, $Y_1 : [0, m] \to F(P_1)$, and $E' : [0, m] \to F(\mathbb{R}^2)$ also have one-sided limits and finitely many discontinuities, which arise from discontinuities of $q_0$, $q_1$ and from door events in between.

The following lemma is the heart of the proof of correctness of the algorithm. Informally speaking, it asserts that the primary region $E'$ can only shrink in time, since the primary door $D'$ always sweeps continuously into or jumps into $E'$.

**Lemma 16.** The functions $Y_0$ and $Y_1$ are monotonically decreasing.

**Proof.** First, we let $I$ be an arbitrary subinterval of $[0, m]$ in which $q_0$ and $q_1$ are continuous, and we prove the lemma for functions restricted to $I$: $y_0 \mid I$, $y_1 \mid I$, $Y_0 \mid I$ and $Y_1 \mid I$. Following the convention from Section 4.3, we consider doors as continuous functions $D : I_D \to F(\mathbb{R}^2)$ with $I_D \subseteq I$ and, for $t \in I$, we let $D(t)$ denote the set of doors $D$ such that $t \in I_D$. Accordingly, we redefine $D'(t)$ to denote the function $D : I_D \to F(\mathbb{R}^2)$ that is chosen as the primary door at time $t \in I$.

By Observation 12, every door $D : I_D \to F(\mathbb{R}^2)$ maintains constant weight over the entire time interval $I_D$. By Observation 11, the only possible changes to $D$ and to the order on $D$ over time interval $I$ are that doors are being added to or removed from $D$. Therefore, any change to the choice of the primary door can only occur at a point in time $t \in I$ when a door event happens; moreover, the primary door $D'(t)$ must participate in that event, that is, if $D'(t) = D$, then $t = \inf I_D$ or $t = \sup I_D$.

Consider an interval $I' \subseteq I$ over which the choice of the primary door remains constant, that is, there is a door $D : I_D \to F(\mathbb{R}^2)$ such that $I' \subseteq I_D$ and $D'(t) = D$ for every $t \in I'$. Since $D$ is a continuous function, so are the functions $y_0 \mid I'$, $y_1 \mid I'$, $Y_0 \mid I'$ and $Y_1 \mid I'$. Furthermore, it follows from Observation 6 that the segment $S$ is constant or is sweeping continuously to the left at any point in time $t \in I$. By Observation 15, $D$ can only be moving towards the interior of $E'$ in time interval $I'$. This shows that $E' \cap I'$ and hence $Y_0 \mid I'$ and $Y_1 \mid I'$ are monotonically decreasing functions.

In view of Lemma 5, to complete the proof that $Y_0 \mid I$ and $Y_1 \mid I$ are monotonically decreasing, it remains to prove that

1. A fence splits into two doors. Since $S(t)$ can only be sweeping to the left, both polygon edges incident to the corner triggering that event lie to the left of $S(t)$. Therefore, by
Observation 11, the two doors are a left-door followed by a right-door in the order on $D(t)$. Consequently, by Observation 14, neither of the two doors can be primary.

2. Two doors merge into a fence. If $D'(t)$ is one of the two doors, then the choice of the primary door changes to some door $D \in D(\wedge t) \subseteq D(t)$ that is after $D'(t)$ in the order on $D(t)$. By Lemma 9, the endpoints $y_0(\wedge t)$ and $y_1(\wedge t)$ of $D(t)$ lie on $Y_0(t)$ and $Y_1(t)$, respectively, so $Y_0(\wedge t) \subseteq Y_0(t)$ and $Y_1(\wedge t) \subseteq Y_1(t)$ as required.

A door splits into a smaller door and a fence. It follows from Observation 11 that the door added to $D(t)$ maintains the weight of the door removed from $D(\wedge t)$. Therefore, assuming $D'(t) \neq D'(\wedge t)$, $D'(t)$ is the door added to $D(t)$ and $D'(\wedge t)$ is the one removed from $D(\wedge t)$. Let $w \in P_k$ denote the corner that triggers the event, where $k \in \{0, 1\}$. It follows that $y_{1-k}(t) = y_{1-k}(\wedge t)$, so $Y_{1-k}(t) = Y_{1-k}(\wedge t)$. Since $w = y_k(t)$, we need to prove that $w \in Y_k(\wedge t)$. Let $D = D'(\wedge t)$ and let $t_0$ be a value in $I_D \cap I$ such that $t_0 < t$ and no door event happens in time interval $[t_0, t]$. For every $t' \in [t_0, t]$, let $\varphi(t')$ be the point on $D(t')$ closest to $w$. Since $D$ moves continuously, $\varphi$ is a continuous curve. Since $\varphi(t') \in E'(t')$ and $E'(t') \subseteq E'(t_0)$ as shown before for every $t' \in [t_0, t)$, $\varphi$ must be contained in $E'(t_0)$. Since $w \in D(\wedge t)$, we have $\varphi(\wedge t) = w$. Furthermore, $E'(t_0)$ is a closed set, and hence $w \in E'(t_0)$. The interior of $E'(t_0)$ is disjoint from $P_0$ and $P_1$, so $w$ must be a corner on the chain $Y_k(t_0) = P_k[y_k(t_0), r_k]$. Since $w = y_k(t)$, it follows that $Y_k(t) \subseteq Y_k(t_0)$.

4. A door and a fence merge into a larger door. Again, it follows from Observation 11 that the door added to $D(\wedge t)$ maintains the weight of the door removed from $D(t)$. Therefore, assuming $D'(t) \neq D'(\wedge t)$, $D'(t)$ is the door removed from $D(t)$ and $D'(\wedge t)$ is the one added to $D(\wedge t)$. Let $w \in P_k$ denote the corner that triggers the event, where $k \in \{0, 1\}$. It follows that $y_{1-k}(t) = y_{1-k}(\wedge t)$, so $Y_{1-k}(t) = Y_{1-k}(\wedge t)$. We make an argument similar to the one in the above case to show that $Y_k(\wedge t) \subseteq Y_k(t)$, but using reversed time. Let $D = D'(\wedge t)$ and let $t_0$ be a value in $I_D \cap I$ such that $t_0 > t$ and no door event happens in time interval $[t, t_0]$. For every $t' \in (t, t_0]$, let $\varphi(t')$ be the point on $D(t')$ closest to $w$. Since $D$ moves continuously, $\varphi$ is a continuous curve. For every $t' \in [0, m]$, let $F(t')$ be the polygonal region bounded by $P_0[\ell_0, y_0(t')]$, the primary door at time $t'$, $P_1[\ell_1, y_1(t')]$, and the segment $L = \ell_0 \ell_1$. Thus $F(t')$ is a sort of complementary region to $E'(t')$ in the region $E$. Since $E'$ is monotonically decreasing on $(t, t_0]$ as shown before, $F'$ is monotonically increasing on $(t, t_0]$. Therefore, since $\varphi(t') \in F(t')$, $\varphi$ must be contained in $F(t_0)$. Since $w \in D(\wedge t)$, we have $\varphi(\wedge t) = w$. Furthermore, $F(t_0)$ is a closed set, and hence $w \in F(t_0)$. The interior of $F(t_0)$ is disjoint from $P_0$ and $P_1$, so $w$ must be a corner on the chain $P_k[\ell_0, y_k(t_0)]$. Since $w = y_k(t)$, it follows that $P_k[\ell_0, y_k(t_0)] \subseteq P_k[\ell_0, y_k(t_0)]$. By letting $t_0$ approach $t$ from above, we get $P_k[\ell_0, y_k(t)] \subseteq P_k[\ell_0, y_k(t)]$. Hence, $y_k(\wedge t)$ is on the chain $Y_k(t) = P_k[y_k(t), r_k]$ and therefore $Y_k(\wedge t) \subseteq Y_k(t)$.

Now, we return to the general case of functions $y_0$, $y_1$, $Y_0$ and $Y_1$ defined on the entire interval $[0, m]$. Consider a point in time $t \in [0, m]$ that is a discontinuity of $q_k$, where $k \in \{0, 1\}$. That is, $s_k$ jumps from $s_k(t)$ to $s_k(\wedge t)$ at time $t$. We shall see that the jump of $s_k$ has no effect on the choice of the primary door. By Observation 7, the point $p_k[s_k(\wedge t)] = q_k(\wedge t)$ lies on the ray $R(q_{1-k}(t), q_k(t))$ and the chain $P_k[q_k(t), q_k(\wedge t)]$ belongs to $\text{RHP}(q_0(t), q_1(t))$. Let $D(t)$ and $D(\wedge t)$ denote the sets of doors as defined for the segments $S(t)$ and $S(\wedge t)$, respectively. We shall prove that the primary door with respect to $D(t)$ (i.e., defined for $S(t)$) is the same as with respect to $D(\wedge t)$ (i.e., defined for $S(\wedge t)$).

Suppose $q_0(\wedge t)$ is on the segment $S(t)$. It follows that $S(\wedge t) \subseteq S(t)$, $D(\wedge t) \subseteq D(t)$, and $D(t) \setminus D(\wedge t)$ is the set of doors on $S(t) \setminus S(\wedge t)$. Since $P_k[q_k(t), q_k(\wedge t)] \subseteq \text{RHP}(q_0(t), q_1(t))$, it follows from Lemma 10 that the doors in $D(t) \setminus D(\wedge t)$ occur in pairs of a left-door followed
by a right-door, consecutive in the order on $D(t)$. Therefore, the weights of every door $D \in D(\leq t)$ with respect to the sets of doors $D(\leq t)$ and $D(t)$ are equal. By Observation 14, none of the doors in $D(t) \setminus D(\leq t)$ can be primary with respect to $D(t)$, so the primary door is the same with respect to $D(t)$ as with respect to $D(\leq t)$.

Now, suppose $q_k(\leq t)$ is not on the segment $S(t)$. It follows that $S(t) \subset S(\leq t)$, $D(t) \subset D(\leq t)$, and $D(\leq t) \setminus D(t)$ is the set of doors on $S(\leq t) \setminus S(t)$. An argument analogous to that for $q_k(\leq t) \in S(t)$ above shows that the primary door is the same with respect to $D(\leq t)$ as with respect to $D(t)$.

To conclude, let $t_0 = 0$, $t_1, \ldots, t_{n-1}$ be the discontinuities of $q_0$ or $q_1$ ordered so that $t_1 < \cdots < t_{n-1}$, and $t_n = m$, and consider the closed intervals $I_i = [t_{i-1}, t_i]$ for $i \in \{1, \ldots, n\}$. Fix an index $i$ and consider the restrictions $q_0 \upharpoonright I_i$ and $q_1 \upharpoonright I_i$. Only one of them, say $q_k \upharpoonright I_i$, is not continuous, and the only discontinuity of $q_k \upharpoonright I_i$ is $t_{i-1}$. As we have proved above, if we redefine $q_k(t_{i-1})$ by letting $q_k(t_{i-1}) = q_k(\leq t_{i-1})$, the primary door $D'(t_{i-1})$ does not change, but then $q_k \upharpoonright I_i$ becomes continuous. Therefore, what we have proved for restrictions of $Y_0$ and $Y_1$ to subintervals $I \subseteq [0, m]$ such that $q_0 \upharpoonright I$ and $q_1 \upharpoonright I$ are continuous implies that $Y_0 \upharpoonright I_i$ and $Y_1 \upharpoonright I_i$ are monotonically decreasing, for every $i \in \{1, \ldots, n\}$. The assumptions of Lemma 5 are satisfied for $Y_0$ and $Y_1$, so $Y_0$ and $Y_1$ are monotonically decreasing in the entire domain $[0, m]$.

We are now ready to prove Lemma 2. Using the continuous interpretation of the algorithm, it can be rephrased as follows.

\begin{lemma}
For any $t \in [0, m]$, we have $0 \leq s_0(t) < 2n_0$ and $0 \leq s_1(t) < 2n_1$.
\end{lemma}

\textbf{Proof.} We only present the proof of the bound on $s_0(t)$. That for $s_1(t)$ is analogous. Let $c_0(0)$ be the unique real in the interval $[0, n_0)$ such that $y_0(0) = p_0[c_0(0)]$. Let $c_0$ be the unique real in the interval $[c_0(0), c_0(0) + n_0)$ such that $y_0(t) = p_0[c_0(t)]$. By Lemma 16, for $t \in (0, m]$, there is a unique real $c_0(t) \in [c_0(0), c_0(0) + n_0)$ such that $y_0(t) = p_0[c_0(t)]$, and this defines a nondecreasing function $c_0: [0, m] \to \mathbb{R}$ with one-sided limits and finitely many discontinuities. Obviously, $0 \leq s_0(t)$ and $c_0(t) \leq c_0 < c_0(0) + n_0 < 2n_0$. It remains to prove $s_0(t) \leq c_0(t)$ for $t \in [0, m]$.

We first prove that for every $t \in [0, m]$ with $s_0(t) \leq c_0(t)$, there is $\varepsilon > 0$ such that
\[ s_0(t') \leq c_0(t') \]
for all $t' \in [t, t + \varepsilon)$. Let $t \in [0, m]$ be such that $s_0(t) \leq c_0(t)$. First, suppose $s_0$ is continuous and either constant or strictly increasing on some interval $[t, t + \varepsilon)$ with $\varepsilon > 0$. If $s_0(t) < c_0(t)$, then the statement is clear, so suppose $s_0(t) = c_0(t)$. If $s_0$ is constant on $[t, t + \varepsilon)$, then the statement is clear, as $c_0$ is nondecreasing. If $s_0$ is strictly increasing on $[t, t + \varepsilon)$, we either have $c_0(t') = s_0(t')$ for $t' \in [t, t + \varepsilon)$, for some $\varepsilon' \in (0, \varepsilon]$, or $c_0$ jumps at time $t$ to a higher value, that is, $c_0(t) < c_0(\leq t)$. In both cases, the statement holds.

Now, suppose $s_0$ jumps at time $t$, that is, $s_0(t) < s_0(\leq t)$. By choosing $\varepsilon > 0$ small enough, we can assume that $s_0$ is continuous on the interval $(t, t + \varepsilon)$ and that the points $\{p_0[s_0(t')]: t' \in (t, t + \varepsilon)\}$ are a part of one edge $e$ of $P_0$, which also contains the point $p_0[s_0(\leq t)]$. By Observation 7, we have $P_0[s_0(t), s_0(\leq t)] \subset \text{RHP}(p_0[s_0(t)], p_1[s_1])$. This and the facts that $p_0[c_0(t')] \in S(t')$ and $S(t') \cap \text{RHP}(p_0[s_0(t)], p_1[s_1]) = \{p_1[s_1]\}$ imply $c_0(t') > s_0(t')$, for every $t' \in (t, t + \varepsilon)$. We conclude that for every $t' \in (t, t + \varepsilon)$, either $c_0(t') = s_0(t')$ or $p_0[c_0(t')]$ is on an edge of $P_0$ other than $e$, in which case $c_0(t') > s_0(t')$.

We now return to proving that $s_0(t) \leq c_0(t)$ for every $t \in [0, m]$. Suppose the contrary, and let $t^* = \inf\{t \in [0, m]: s_0(t) > c_0(t)\}$. In view of the discussion above, we must have $s_0(t^*) > c_0(t^*)$. Then $t^* > 0$, because $c_0(0) \geq 0 = s_0(0)$. By the definition of $s_0$, we have $s_0(\leq t^*) = s_0(t^*) > c_0(t^*) \geq c_0(\leq t^*)$. This contradicts the definition of $t^*$.  

\hfill $\blacksquare$
References


