Relational algebra by way of adjunctions
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Relational Algebra by Way of Adjunctions

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(joint work with Fritz Henglein, Ralf Hinze, Nicolas Wu)
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1. Summary

- bulk types (sets, bags, lists) are *monads*
- monads have nice *mathematical foundations via adjunctions*
- monads support *comprehensions*
- comprehension syntax provides a *query notation*

\[
\begin{align*}
\left[ (\text{customer.name}, \text{invoice.amount}) \\
| \text{customer} \leftarrow \text{customers}, \\
\text{invoice} \leftarrow \text{invoices}, \\
\text{customer.cid} = \text{invoice.customer}, \\
\text{invoice.due} \leq \text{today} \right]
\end{align*}
\]

- monad structure explains *selection, projection*
- less obvious how to explain *join*
2. Galois connections

Relating monotonic functions between two ordered sets:

\[(A, \leq) \downarrow (B, \sqsubseteq) \quad \text{means} \quad f(b) \leq a \iff b \sqsubseteq g(a)\]

For example,

\[(\mathbb{R}, \leq_{\mathbb{R}}) \downarrow (\mathbb{Z}, \leq_{\mathbb{Z}}) \quad \text{inj} \]

\[(\mathbb{Z}, \leq) \downarrow (\mathbb{Z}, \leq) \quad \times k \]

“Change of coordinates” can sometimes simplify reasoning; e.g., rhs gives $n \times k \leq m \iff n \leq m \div k$, and multiplication is easier to reason about than rounding division.
3. Category theory from ordered sets

A category $\mathbf{C}$ consists of

- a set $|\mathbf{C}|$ of objects,
- a set $\mathbf{C}(X, Y)$ of arrows $X \to Y$ for each $X, Y : |\mathbf{C}|$,
- identity arrows $\text{id}_X : X \to X$ for each $X$,
- composition $f \cdot g : X \to Z$ of compatible arrows $g : X \to Y$ and $f : Y \to Z$,

such that composition is associative, with identities as units.

Think of a directed graph, with vertices as objects and paths as arrows.

An ordered set $(A, \leq)$ is a degenerate category, with objects $A$ and a unique arrow $a \to b$ iff $a \leq b$.

\[ \cdots \longrightarrow -2 \longrightarrow -1 \longrightarrow 0 \longrightarrow 1 \longrightarrow 2 \longrightarrow \cdots \]

Many categorical concepts are generalisations from ordered sets.

*proviso...
4. Concrete categories

Ordered sets are a *concrete category*: roughly,

- the objects are *sets with additional structure*
- the arrows are *structure-preserving mappings*

Many useful categories are of this form.

For example, the category $\mathbf{CMon}$ has commutative monoids $(M, \otimes, \epsilon)$ as objects, and homomorphisms $h : (M, \otimes, \epsilon) \to (M', \oplus, \epsilon')$ as arrows:

\[
\begin{align*}
h (m \otimes n) &= h m \oplus h n \\
h \epsilon &= \epsilon'
\end{align*}
\]

Trivially, category $\mathbf{Set}$ has sets as objects, and total functions as arrows.
5. Functors

Categories are themselves structured objects...

A functor $F : C \to D$ is an operation on both objects and arrows, preserving the structure: $F f : F X \to F Y$ when $f : X \to Y$, and

$$F \text{id}_X = \text{id}_{F X}$$
$$F (f \cdot g) = F f \cdot F g$$

For example, forgetful functor $U : \text{CMon} \to \text{Set}$:

$$U (M, \otimes, \epsilon) = M$$
$$U (h : (M, \otimes, \epsilon) \to (M', \oplus, \epsilon')) = h : M \to M'$$

Conversely, $\text{Free} : \text{Set} \to \text{CMon}$ generates the free commutative monoid (ie bags) on a set of elements:

$$\text{Free } A = (\text{Bag } A, \cup, \emptyset)$$
$$\text{Free } (f : A \to B) = \text{map } f : \text{Bag } A \to \text{Bag } B$$
6. Adjunctions

**Adjunctions** are the categorical generalisation of Galois connections. Given categories \( C, D \), and functors \( L : D \to C \) and \( R : C \to D \), adjunction

\[
\begin{array}{c}
\rightarrow \quad \downarrow \quad \rightarrow \\
L \quad \downarrow \quad D \quad \downarrow \quad R \\
\leftarrow \quad C \quad \downarrow \quad \leftarrow \\
\end{array}
\]

means \( [-] : C(LX, Y) \simeq D(X, RY) : [-] \)

A familiar example is given by **currying**:

\[
\begin{array}{c}
\rightarrow \quad \downarrow \quad \rightarrow \\
- \times P \quad \downarrow \quad \leftarrow \\
\leftarrow \quad \text{Set} \quad \downarrow \quad \leftarrow \\
(\cdot)^P \quad \downarrow \quad \text{Set} \\
\end{array}
\]

with \( \text{curry} : \text{Set}(X \times P, Y) \simeq \text{Set}(X, Y^P) : \text{curry}^\circ \)

hence definitions and properties of \( \text{apply} = \text{uncurry \ id}_{Y^P} : Y^P \times P \to Y \)
7. Products and coproducts

\[
\begin{array}{ccc}
\text{Set} & \downarrow & \text{Set}^2 \\
\Delta & & \Delta \\
\downarrow & & \downarrow \\
\Delta & & \times \\
\end{array}
\]

with

\[\text{fork} : \text{Set}^2(\Delta A, (B, C)) \cong \text{Set}(A, B \times C) : \text{fork}^\circ\]
\[\text{junc}^\circ : \text{Set}(A + B, C) \cong \text{Set}^2((A, B), \Delta C) : \text{junc}\]

hence

\[\text{dup} = \text{fork} \ id_{A,A} : \text{Set}(A, A \times A)\]
\[(\text{fst}, \text{snd}) = \text{fork}^\circ \ id_{B \times C} : \text{Set}^2(\Delta (B, C), (B, C))\]

give tupling and projection. Dually for sums and injections, and generally for any arity—even zero.
8. Free commutative monoids

Adjunctions often capture embedding/projection pairs:

\[
\text{CMon} \quad \perp \quad \text{Set}
\]

\[
\text{with } [-] : \text{CMon}(\text{Free } A, (M, \otimes, \epsilon)) \cong \text{Set}(A, U (M, \otimes, \epsilon)) : [-]
\]

Unit and counit:

\[
single A = [id_{\text{Free } A}] : A \to U (\text{Free } A)
\]

\[
\text{reduce } M = [id_M] : \text{Free} (U M) \to M \quad \text{-- for } M = (M, \otimes, \epsilon)
\]

whence, for \( h : \text{Free } A \to M \) and \( f : A \to U M = M \),

\[
h = \text{reduce } M \cdot \text{Free } f \iff U h \cdot \text{single } A = f
\]

ie 1-to-1 correspondence between homomorphisms from the free commutative monoid (bags) and their behaviour on singletons.
9. Aggregation

Aggregations are bag homomorphisms:

<table>
<thead>
<tr>
<th>aggregation</th>
<th>monoid</th>
<th>action on singletons</th>
</tr>
</thead>
<tbody>
<tr>
<td>count</td>
<td>((\mathbb{N}, 0, +))</td>
<td>({a} \rightarrow 1)</td>
</tr>
<tr>
<td>sum</td>
<td>((\mathbb{R}, 0, +))</td>
<td>({a} \rightarrow a)</td>
</tr>
<tr>
<td>max</td>
<td>((\mathbb{Z}, \text{minBound}, \text{max}))</td>
<td>({a} \rightarrow a)</td>
</tr>
<tr>
<td>min</td>
<td>((\mathbb{Z}, \text{maxBound}, \text{min}))</td>
<td>({a} \rightarrow a)</td>
</tr>
<tr>
<td>all</td>
<td>((\mathbb{B}, \text{True}, \land))</td>
<td>({a} \rightarrow a)</td>
</tr>
<tr>
<td>any</td>
<td>((\mathbb{B}, \text{False}, \lor))</td>
<td>({a} \rightarrow a)</td>
</tr>
</tbody>
</table>

Selection is a homomorphism, to bags, using action

\[
guard : (A \rightarrow \mathbb{B}) \rightarrow \text{Bag } A \rightarrow \text{Bag } A
\]

\[
guard p a = \text{if } p a \text{ then } \{a\} \text{ else } \emptyset
\]

Laws about selections follow from laws of homomorphisms
(and of coproducts, since \(\mathbb{B} = 1 + 1\)).
10. Monads

Bags form a *monad* \((\text{Bag}, \text{union}, \text{single})\) with

\[
\begin{align*}
\text{Bag} & = U \cdot \text{Free} \\
\text{union} & : \text{Bag} (\text{Bag} A) \to \text{Bag} A \\
\text{single} & : A \to \text{Bag} A
\end{align*}
\]

which justifies the use of comprehension notation \(\{ f \ a \ b \mid a \leftarrow x, b \leftarrow g a \}\).

In fact, for any adjunction \(L \dashv R\) between \(C\) and \(D\), we get a monad \((T, \mu, \eta)\) on \(D\), where

\[
\begin{align*}
T & = R \cdot L \\
\mu A & = R \left[ id_A \right] L : T (T A) \to T A \\
\eta A & = \left[ id_A \right] : A \to T A
\end{align*}
\]
11. Maps

Database indexes are essentially maps $\text{Map } K V = V^K$. Maps $(-)^K$ from $K$ form a monad (the $\text{Reader}$ monad in Haskell), so arise from an adjunction.

The *laws of exponents* arise from this adjunction, and from those for products and coproducts:

- $\text{Map } 0 V \simeq 1$
- $\text{Map } 1 V \simeq V$
- $\text{Map } (K_1 + K_2) V \simeq \text{Map } K_1 V \times \text{Map } K_2 V$
- $\text{Map } (K_1 \times K_2) V \simeq \text{Map } K_1 (\text{Map } K_2 V)$
- $\text{Map } K 1 \simeq 1$
- $\text{Map } K (V_1 \times V_2) \simeq \text{Map } K V_1 \times \text{Map } K V_2 : \text{merge}$
12. Indexing

Relations are in 1-to-1 correspondence with set-valued functions:

\[
\text{Rel} \downarrow \downarrow \downarrow \text{Set}
\]

where \( J \) embeds, and \( E R : A \rightarrow \text{Set} B \) for \( R : A \sim B \).

Moreover, the correspondence remains valid for bags:

\[
\text{index} : \text{Bag} (K \times V) \simeq \text{Map} K (\text{Bag} V)
\]

Together, \( \text{index} \) and \( \text{merge} \) give efficient relational joins:

\[
x f \bowtie_g y = \text{flatten} (\text{Map} K \ cp (\text{merge} (\text{groupBy} f x, \text{groupBy} g y)))
\]

\[
\text{groupBy} : (V \rightarrow K) \rightarrow \text{Bag} V \rightarrow \text{Map} K (\text{Bag} V)
\]

\[
\text{flatten} : \text{Map} K (\text{Bag} V) \rightarrow \text{Bag} V
\]
13. Pointed sets and finite maps

Model *finite maps* \( \text{Map}_* \) not as partial functions, but *total* functions to a *pointed* codomain \((A, a)\), i.e. a set \( A \) with a distinguished element \( a : A \).

Pointed sets and point-preserving functions form a category \( \text{Set}_* \).

There is an adjunction to \( \text{Set} \), via

\[
\begin{array}{c}
\text{Set}_* \\
\downarrow \\
\text{Set}
\end{array}
\xleftarrow{\text{Maybe}} \xrightarrow{\text{U}}
\]  

where \( \text{Maybe} A \simeq 1 + A \) adds a point, and \( U (A, a) = A \) discards it.

In particular, \((\text{Bag} A, \emptyset)\) is a pointed set. Moreover, \( \text{Bag} f \) is point-preserving, so we get a functor \( \text{Bag}_* : \text{Set} \to \text{Set}_* \).

Indexing remains an isomorphism:

\[
\text{index} : \text{Bag}_* (K \times V) \simeq \text{Map}_* K (\text{Bag}_* V)
\]
14. Graded monads

A catch: finite maps aren’t a monad, because

\[ \eta a = \lambda k \rightarrow a : A \rightarrow \text{Map } K A \]

in general yields an infinite map.

However, finite maps are a graded monad*:

\[
\mu X : T_m (T_n X) \rightarrow T_{m \otimes n} X \\
\eta X : X \rightarrow T_\epsilon X
\]

satisfying the usual laws. These too arise from adjunctions*.

We use the monoid \((\mathbb{K}, \times, 1)\) of finite key types under product.
15. Conclusions

- *monad comprehensions* for database queries
- structure arising from *adjunctions*
- equivalences from *universal properties*
- fitting in *relational joins*, via indexing
- to do: calculating *query optimisations*

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